On Unique Recovery of Integer Bounded Signals

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Summary

The thesis considers the problem of unique recovery of finite-valued integer signals using a single linear integer measurement. A signal is an integer ndimensional vector x with absolute entries bounded by a positive integer r, that is $x \in [-r, r]^n$. We assume that the signal x is sufficiently sparse. Specifically, the number of nonzero entries of x is assumed to be bounded by a positive integer l with 2l < n. A single linear integer measurement is represented by an integer $1 \times n$ measurement matrix, or row vector, H. Naturally, it is desirable to construct H with as small absolute entries as possible. We give a constructive proof for the existence of measurement matrices H with maximum absolute entry $\Delta = O(r^{2l-1})$. The capital *O* in this bound contains an implicit constant that depends on *l* and *n* and probably far from being optimal, however the exponent 2l-1 is optimal. The optimality of the exponent is the main advantage of the latter upper bound. Additionally, we show that, in the above setting, a single measurement can be replaced by several measurements with absolute entries sub-linear in Δ . The proofs make use of results on admissible (n-1)-dimensional integer lattices for *m*-sparse *n*cubes that are of independent interest. The main tools include the aggregation of linear Diophantine equations and Siegel's lemma. Additionally, we discuss some probabilistic aspects of unique recovery for finite-valued integer signals.

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Chapter 1

Preliminaries

1.1 Basics and notation

We denote by \mathbb{R} the set of real numbers, by \mathbb{Q} the set of rational numbers and by \mathbb{Z} the set of integers. For a real number $x \in \mathbb{R}$, we denote by $\lceil x \rceil$ the *ceiling function* of x, that is the least integer greater than or equal to x. Similarly, $\lfloor x \rfloor$ stands for the *floor function* of x, the greatest integer less than or equal to x. The *support* of a vector $\mathbf{x} \in \mathbb{R}^n$ is denoted by $\operatorname{supp}(\mathbf{x})$. That is $\operatorname{supp}(\mathbf{x}) = \{i : \mathbf{x}_i \neq 0\}$. Given $K \subset \mathbb{R}^n$, we use the notation $\operatorname{vol}(K)$ for the volume , i.e., the *n*-dimensional Lebesgue measure of the set K.

Norms

A general norm on a vector space can be defined as follows.

Definition 1. A function $\|\cdot\| : \mathbb{R}^n \to \mathbb{R}$ is called a *norm* if it satisfies three properties:

(i) $||\mathbf{x}|| \ge 0$ for any vector $\mathbf{x} \in \mathbb{R}^n$, and $||\mathbf{x}|| = 0$ if and only if $\mathbf{x} = \mathbf{0}$.



Figure 1.1: The unit norm balls for ℓ_p -norms.

(ii) $\|\alpha x\| = |\alpha| \|x\|$ for any vector $x \in \mathbb{R}^n$ and any scalar $\alpha \in \mathbb{R}$.

(iii) $||x + y|| \le ||x|| + ||y||$ for any vectors $x, y \in \mathbb{R}^n$.

The first property is referred to as the positivity of the norm, and the second expresses its homogenity, while the third is known as the triangle inequality.

The most commonly used norms belong to the family of ℓ_p -norms, which are defined for $x \in \mathbb{R}^n$ as

$$\|\mathbf{x}\|_{p} = \left(\sum_{i=1}^{n} |x_{i}|^{p}\right)^{1/p},$$

where $p \ge 1$ is a real number.

Here are some particularly important ℓ_p -norms:

• p = 1: the ℓ_1 -norm, which is also referred to as the Manhattan norm, is

$$\|\mathbf{x}\|_1 = |x_1| + |x_2| + \dots + |x_n|.$$

• p = 2: the ℓ_2 -norm, which is also referred to as the *Euclidean norm*, is

$$\|\mathbf{x}\|_2 = \sqrt{x_1^2 + x_2^2 + \dots + x_n^2}$$

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• As p approaches the infinity, the ℓ_p norm approaches the ℓ_{∞} -norm, which is also referred to as the *maximum norm*,

$$\|\boldsymbol{x}\|_{\infty} = \max_{1 \le i \le n} |x_i|.$$

The unit norm balls $\{x \in \mathbb{R}^2 : ||x||_p \le 1\}$ for the above norms are shown in Figure 1.1.

In this thesis, an important role will be played by the ℓ_0 -"norm", widely used in the theory of compressed sensing.

Definition 2. For a given vector $x \in \mathbb{R}^n$, the ℓ_0 -norm $||x||_0$ is the cardinality of the support of x, that is the number of non-zero entries of x

$$\|\boldsymbol{x}\|_0 = |\operatorname{supp}(\boldsymbol{x})|.$$

In fact ℓ_0 -norm is not a proper norm, as it is not homogeneous. If x is a non-zero vector and $\alpha \neq 1$, then

$$\|\alpha x\|_0 \neq \alpha \|x\|_0$$

ℓ_{∞} -norm of a matrix

The ℓ_{∞} -norm for a vector \mathbf{x} is the largest absolute value of its components. For a matrix $A = (a_{ij}) \in \mathbb{R}^{m \times n}$ we define the ℓ_{∞} -norm $||A||_{\infty}$ in a similar way.

Definition 3. For a matrix $A = (a_{ij}) \in \mathbb{R}^{m \times n}$, the ℓ_{∞} -norm is the maximum absolute value of the entries of A, that is

$$\|A\|_{\infty} = \max_{\substack{1 \le i \le m \\ 1 \le j \le n}} \left| a_{ij} \right|.$$



Figure 1.2: Comparison of different lattices in \mathbb{Z}^2 .

1.2 Lattices

Minkowski's Geometry of Numbers provides a mathematical toolbox for this thesis. One of the most important mathematical instruments that we use are the lattices. We refer the reader to [30, 13] for an extensive introduction to the geometry of numbers.

Definition 4. Let $\boldsymbol{b}_1, \ldots, \boldsymbol{b}_k \in \mathbb{R}^n$ be linearly independent vectors. The set

$$\Lambda = \left\{ \sum_{i=1}^k x_i \boldsymbol{b}_i : x_i \in \mathbb{Z} \right\}$$

is a *k*-dimensional lattice with basis b_1, \ldots, b_k . If n = k, then Λ is a fulldimensional lattice, as illustrated in Figure 1.2b.

Let $B = (\boldsymbol{b}_1, \dots, \boldsymbol{b}_k) \in \mathbb{R}^{n \times k}$ be a matrix with columns $\boldsymbol{b}_1, \dots, \boldsymbol{b}_k$, referred to as a *basis matrix*. We can define the lattice $\Lambda = \Lambda(B)$ as

$$\Lambda = \left\{ B\boldsymbol{x} : \boldsymbol{x} \in \mathbb{Z}^k \right\}.$$

Notice that, since the basis vectors $\mathbf{b}_1, \dots, \mathbf{b}_k$ are linearly independent, any point \mathbf{y} in the subspace spanned by the basis vectors can be written as a



Figure 1.3: Two different bases generate the same lattice.

unique linear combination $\mathbf{y} = x_1 \mathbf{b}_1 + x_2 \mathbf{b}_2 + \dots + x_k \mathbf{b}_k$, for some $x_1, x_2, \dots, x_k \in \mathbb{R}$. \mathbb{R} . Therefore, $\mathbf{y} \in \Lambda$ if and only if $x_j \in \mathbb{Z}$ for all $j \in \{1, 2, \dots, k\}$.

Note that the choice of a lattice basis is not unique, the same lattice can be generated by different bases, as we can see in the next example.

Example 1. Take

$$B = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, and C = \begin{pmatrix} 1 & 2 \\ 1 & 1 \end{pmatrix}.$$

Then as it is illustrated by Figure 1.3, B and C generate the same lattice \mathbb{Z}^2 , that is $\Lambda(B) = \Lambda(C) = \mathbb{Z}^2$.

In general, one can wonder when two different sets of linearly independent vectors generate the same lattice. The answer is given by the unimodular basis transformations.

Definition 5. Let $U \in \mathbb{Z}^{k \times k}$ be an integer square matrix. Then *U* is called a *unimodular matrix* if it has determinant det(*U*) = ±1.

Two basis matrices *B* and *C* generate the same lattice $\Lambda(B) = \Lambda(C)$ if and only if B = CU for a unimodular matrix *U*, see [46].

Next, we will introduce the *fundamental parallelepiped* of a lattice. Let $\Lambda \subset \mathbb{R}^n$ be a lattice with basis matrix $B = (\mathbf{b}_1, \dots, \mathbf{b}_k) \in \mathbb{R}^{n \times k}$. The fundamental parallelepiped of the lattice Λ associated with basis *B* is defined as

$$P(B) = \{B\mathbf{x} : \mathbf{x} \in \mathbb{R}^n, 0 \le \mathbf{x}_i < 1\}$$

It is worth pointing out that because the column vectors in *B* form a basis for Λ , it follows that the fundamental parallelepiped $\mathbb{P}(B)$ contains no other lattice points except the origin, i.e. $P(B) \cap \Lambda = \{\mathbf{0}\}$.

It is clear that the fundamental parallelepiped P(B) depends on the basis *B* of the lattice Λ . See Figure 1.4.

Definition 6. The *determinant* of the lattice Λ with a basis matrix *B* is defined as

$$\det(\Lambda) = \sqrt{\det(B^T B)},$$

where B^{\top} denotes the transpose of *B*.

In the special case when Λ is a full-dimensional lattice (that is k = n), the matrix *B* is a square matrix, and hence the determinant of the lattice is given by det(Λ) = |det(B)| = vol(P(B)).

We will now show that the determinant is well defined, that is the value of the determinant does not depend on the choice of a lattice basis. Let the two bases $B_1, B_2 \in \mathbb{R}^{n \times k}$ generate the same lattice Λ . We know that for some unimodular matrix $U \in \mathbb{Z}^{k \times k}$ the equality $B_2 = B_1 U$ holds. Calculating the determinant det(Λ) related to the bases B_1 and B_2 yields that

$$\det(\Lambda) = \sqrt{\det(B_2^T B_2)} = \sqrt{\det(U^T B_1^T B_1 U)} = \sqrt{\det(B_1^T B_1)}$$

holds. A *sublattice* of a full-dimensional lattice $\Lambda \subset \mathbb{R}^n$ is defined as a lattice



Figure 1.4: Different fundamental parallelepipeds dependent on different bases.

 $\widetilde{\Lambda}$ such that $\widetilde{\Lambda} \subset \Lambda$. If \boldsymbol{a} and $\overline{\boldsymbol{a}}$ are lattice points in Λ , and $\widetilde{\Lambda}$ is a sublattice of Λ , then we write

$$a \equiv \bar{a} \mod \tilde{\Lambda} \iff (a - \bar{a}) \in \tilde{\Lambda}.$$

This means that a and \bar{a} belong to the same residue class of Λ with respect to $\tilde{\Lambda}$. It is important to note that there are exactly det $\tilde{\Lambda}$ /det Λ different residue classes of Λ with respect to $\tilde{\Lambda}$.

1.3 Convex sets and convex bodies

Minkowski's Geometry of Numbers studies the interactions of lattices and convex sets. For a comprehensive introduction to convex sets, and to convex geometry in general, we refer the reader to [28].

Definition 7. A set $C \subset \mathbb{R}^n$ is called *convex* if for any $x_1, x_2 \in C$ and for any



Figure 1.5: Examples of convex and non-convex sets.

 $\theta \in [0, 1]$, we have

$$\theta \boldsymbol{x}_1 + (1-\theta)\boldsymbol{x}_2 \in C.$$

In other words, one can move from any point in a convex set *C* to any other point via a line segment within the set. See Figure 1.5.

In what follows, we collect some basic properties that preserve convexity (see e.g. [10])

- 1. The empty set \emptyset and the whole space \mathbb{R}^n are both convex.
- 2. The intersection of convex sets is convex.
- 3. Convexity is preserved by scaling. If $C \subset \mathbb{R}^n$ is a convex set and $\alpha \in \mathbb{R}$, then the set $\alpha C = \{\alpha x : x \in C\}$ is convex.
- 4. Convexity is preserved by summation. The *Minkowski's sum* of two sets C_1 and C_2 is defined as

$$C_1 + C_2 = \{x + y : x \in C_1, y \in C_2\}.$$

If C_1 and C_2 are convex, then $C_1 + C_2$ is convex.



Figure 1.6: Visualisation of the set of points S (left) and its corresponding convex hull (right).

5. The image of a convex set under a linear affine transformation is also a convex set.

Given points $x_1, x_2, ..., x_t$ in \mathbb{R}^n , a *convex combination* of these points is the point in the form

$$\theta_1 \mathbf{x_1} + \theta_2 \mathbf{x_2} + \cdots + \theta_t \mathbf{x_t}$$

where $\theta_1 + \theta_2 + \dots + \theta_t = 1$ and $\theta_i \ge 0$ for $i = 1, 2, \dots, t$.

Given a set *C* in \mathbb{R}^n , its *convex hull*, denoted *conv*(*C*), is the smallest with respect to inclusions convex set that contains *C*. Equivalently, *conv*(*C*) is the set of all convex combinations of points in *C*, See Figure 1.6.

Definition 8. A compact (that is closed and bounded) convex set $C \subset \mathbb{R}^n$ with nonempty interior is called a *convex body*.

We will denote the set of all convex bodies in \mathbb{R}^n by \mathscr{K}^n , and the set of all **0**-symmetric convex bodies by \mathscr{K}_0^n .



Figure 1.7: Example of Set of Points C (left) and its Difference Set D(C) (right).

Definition 9. Given a set $C \in \mathbb{R}^n$, the *difference set of* C is defined as

$$D(C) = C - C = \{x - y : x, y \in C\}.$$

Clearly, D(C) is **0**-symmetric. If $C \in \mathcal{K}^n$, then $D(C) \in \mathcal{K}^n_0$, see for example Figure 1.7.

1.4 Minkowski's first fundamental theorem

The first Minkowski's fundamental theorem considers the following question: when does an origin-symmetric convex set contain a non-zero lattice point?

Theorem 1 (Minkowski's 1st fundamental theorem). Let $S \subset \mathbb{R}^n$ be an **0**-symmetric convex set of volume vol(S), and let $L \subset \mathbb{R}^n$ be a lattice with determinant det(L). If vol(S) > 2^n det(L), then S contains a nonzero point $u \in L$.

This remarkable result paved the way for many discoveries in number theory, as well as implications in combinatorics, computational complexity,



Figure 1.8: A two-dimensional lattice *L* and an **0**-symmetric box *C* such that vol(C) > 4 det(L).

cryptography and other disciplines. Figure 1.8 provides a illustration for the Minkowski first fundamental theorem.

We will include in this thesis a proof of Theorem (1) based on the following beautiful result of Blichfeldt. We refer the reader to the classical book of Cassels [13] for a detailed introduction to the geometry of numbers.

Lemma 2 (Blichfeldt's Lemma). Let *L* be a lattice in \mathbb{R}^n and $S \subset \mathbb{R}^n$ a set of volume vol(*S*). Suppose that vol(*S*) > det(*L*), then there exist points $\mathbf{x}_1, \mathbf{x}_2 \in S$, $\mathbf{x}_1 \neq \mathbf{x}_2$, such that $\mathbf{x}_1 - \mathbf{x}_2 \in L$.

Proof. Let b_1, \ldots, b_n be any basis of the lattice *L*, and *F* be the corresponding fundamental parallelepiped of *L*. It follows that vol(F) = det(L). Furthermore, any point $x \in \mathbb{R}^n$ can be expressed as x = u + v, where $u \in L$ and $v \in F$. For each $u \in L$ let R(u) be the set of points v such that

$$v \in F, u + v \in S.$$

Then



Figure 1.9: The sum of the volumes of the sets R(u) for each u in L is equal to the volume of the set S.

$$\sum_{u\in L}\operatorname{vol}(R(u))=\operatorname{vol}(S).$$

Suppose that vol(S) > det(L) Then

$$\sum_{u\in L} \operatorname{vol}(R(u)) > \operatorname{vol}(F).$$

Since R(u) are all contained in F, see Figure 1.9, there must be at least one point $v_0 \in F$ which belongs to at least two of the R(u), say

$$\boldsymbol{v}_0 \in R(\boldsymbol{u}_1)$$
 and $\boldsymbol{v}_0 \in R(\boldsymbol{u}_2)$,

where $u_1 \neq u_2$. Then the points

$$x_1 = v_0 + u_1, \ x_2 = v_0 + u_2$$

are in *S* by definition of R(u) and

$$\boldsymbol{x}_1 - \boldsymbol{x}_2 = \boldsymbol{u}_1 - \boldsymbol{u}_2 \in L \setminus \{\boldsymbol{0}\}.$$

The lemma is proved.

We are now ready to prove the Minkowski's first fundamental theorem.

Proof. The set

$$\frac{1}{2}S = \left\{\frac{1}{2}\boldsymbol{x} : \boldsymbol{x} \in S\right\},\,$$

has volume

$$\operatorname{vol}\left(\frac{1}{2}S\right) = \frac{1}{2^n}\operatorname{vol}(S).$$

Therefore, by the statement of the theorem

$$\operatorname{vol}\left(\frac{1}{2}S\right) > \operatorname{det}(L),$$

This enables us to apply Lemma 2 to the set (1/2)S. In this case, we have

$$\mathbf{x}_1, \mathbf{x}_2 \in \frac{1}{2}S, \ \mathbf{x}_1 \neq \mathbf{x}_2,$$
 (1.1)

such that

$$\boldsymbol{x}_1 - \boldsymbol{x}_2 \in L.$$

Obviously,

$$x_1 = \frac{1}{2}u_1, \ x_2 = \frac{1}{2}u_2,$$

where $u_1, u_2 \in S, u_1 \neq u_2$. We have by (1.1)

$$\frac{1}{2}u_1 - \frac{1}{2}u_2 = \frac{1}{2}u_1 + \frac{1}{2}(-u_2) \in L.$$

Given that *S* is symmetric with respect to the origin, we have $-u_2 \in S$. Moreover, since *S* is convex, it follows that

$$\mathbf{0} \neq \frac{1}{2}\boldsymbol{u}_1 + \frac{1}{2}(-\boldsymbol{u}_2) \in S.$$

The geometry of this argument is illustrated by Figure 1.10. Therefore

$$\mathbf{0} \neq \frac{1}{2}\boldsymbol{u}_1 + \frac{1}{2}\left(-\boldsymbol{u}_2\right) \in S \cap L.$$

Theorem is proved.



Figure 1.10: Midpoint of the line segment between u_1 and u_2 .

1.5 Successive minima

Successive minima of convex bodies with respect to lattices were defined and investigated by Minkowski in the context of the geometry of numbers [35].

Definition 10. Let $C \in \mathscr{K}_0^n$ and $\Lambda \subset \mathbb{R}^n$ be a full-dimensional lattice. The *i*th *successive minimum* $\lambda_i(C, \Lambda)$, $1 \le i \le n$, of *C* with respect to the lattice Λ is defined by

$$\lambda_i(C, \Lambda) = \min\{\lambda > 0 : \dim(\lambda C \cap \Lambda) \ge i\}, \ 1 \le i \le n.$$

In other words, the *i*th successive minimum is the smallest positive dilation factor λ such that λC contains at least *i* linearly independent points of the lattice Λ .

Clearly we have,

$$0 < \lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_n.$$

Example 2. Consider the cube $C^{2}(1/2) = \{x \in \mathbb{R}^{2} : ||x||_{\infty} \le 1/2\}$, and the



Figure 1.11: The first and second successive minima of convex body *C* with respect to the lattices $\Lambda(B)$.

lattice $\Lambda(B)$ with basis matrix

$$B = \begin{pmatrix} 2 & 1 \\ 1 & -1 \end{pmatrix}.$$

The first and second successive minima of $C^2(1/2)$ with respect to $\Lambda(B)$ are $\lambda_1 = 2$ and $\lambda_2 = 4$. See Figure 1.11.

Any full-dimensional lattice Λ can be expressed as $\Lambda = A\mathbb{Z}^n$, where A is a non-singular $n \times n$ basis matrix. Consequently, we have $\lambda_i(C, \Lambda) = \lambda_i \left(A^{-1}C, \mathbb{Z}^n\right)$ and $|C \cap \Lambda| = |A^{-1}C \cap \mathbb{Z}^n|$.

Furthermore, there exists a basis b_1, \ldots, b_n for Λ , such that the linear hull of any *i* of *n* linearly independent lattice points a_1, \ldots, a_n is equal to the linear hull of b_1, \ldots, b_i for $1 \le i \le n$. This means that

$$lin\{\boldsymbol{a}_1,\ldots,\boldsymbol{a}_i\}=lin\{\boldsymbol{b}_1,\ldots,\boldsymbol{b}_i\}.$$

In particular, if $\mathbf{z}_i \in \mathbb{Z}^n$, $1 \le i \le n$, are *n* linearly independent lattice vectors that satisfy $\mathbf{z}_i \in \lambda_i(C, \mathbb{Z}^n)C$, then there exists a unimodular matrix $U \in \mathbb{Z}^{n \times n}$ such that

$$U\boldsymbol{z}_i \in (\lambda_i(UC, \mathbb{Z}^n)UC) \cap \ln\{\boldsymbol{e}_1, \dots, \boldsymbol{e}_i\}, \quad 1 \le i \le n,$$
(1.2)

where $\mathbf{e}_i \in \mathbb{R}^n$ denotes the *i*-th unit vector. Moreover, it can be observed that for *n* linearly independent lattice points $\mathbf{a}_1, \dots, \mathbf{a}_n$ of a lattice Λ satisfying $\mathbf{a}_i \in \lambda_i(C, \Lambda)C$, the definition of the successive minima implies

$$\operatorname{int}(\lambda_i(C,\Lambda)C) \cap \Lambda \subset \operatorname{lin}\{\mathbf{0}, \mathbf{a}_1, \dots, \mathbf{a}_{i-1}\} \cap \Lambda, \quad 1 \le i \le n,$$
(1.3)

where int denotes the interior.

1.6 Minkowski's theorem on successive minima

Minkowski's convex body theorem has many extensions and refinements. See for example Gruber and Lekkerkerker [29] and Lagarias [27]. An important refinement is the second fundamental theorem of Minkowski, also known as Minkowski's theorem on successive minima [28]. It gives optimal upper and lower bounds for the product of the successive minima of an origin-symmetric convex body with respect to a lattice.

Theorem 3 (Minkowski [35]). Let $C \in \mathscr{K}_0^n$ and $\Lambda \subset \mathbb{R}^n$ be a full-dimensional *lattice. Then we have,*

$$\frac{2^n}{n!} \det(\Lambda) \le \prod_{i=1}^n \lambda_i(C,\Lambda) \operatorname{vol}(C) \le 2^n \det(\Lambda)$$

The difficult part of the proof is to show the right-hand inequality. The original proof of Minkowski [35] is rather involved. Alternative proofs were

found by Bambah, Woods and Zassenhaus [8] and others. Below we include an elegant proof due to Henk [31].

Proof. Without loss of generality, we will assume that $\Lambda = \mathbb{Z}^n$. For convenience, we denote $\lambda_i = \lambda_i(C, \mathbb{Z}^n)$, and set $C_i = \frac{\lambda_i}{2}C$. Moreover, we assume that z_1, \ldots, z_n are *n* linearly independent lattice points such that $z_i \in \lambda_i C \cap \mathbb{Z}^n$, and the linear space of z_1, \ldots, z_i is equal to the linear space of e_1, \ldots, e_i , where $1 \le i \le n$. For shortness, we denote the linear space $\lim\{e_1, \ldots, e_i\}$ by L_i .

Let $M_q^n = \{ z \in \mathbb{Z}^n : |z_i| \le q, 1 \le i \le n \}$ be a set of lattice points in \mathbb{Z}^n where each coordinate is bounded by $q \in \mathbb{N}$. For $1 \le j \le n-1$, let $M_q^j = M_q^n \cap L_j$. Since *C* is a bounded set, there exists a constant γ , which depends only on *C*, such that

$$\operatorname{vol}\left(M_{q}^{n}+C_{n}\right)\leq(2q+\gamma)^{n}.$$
(1.4)

Using the definition of λ_1 , for any two distinct lattice points $z, \bar{z} \in \mathbb{Z}^n$, the intersection of $(z + int(C_1))$ and $(\bar{z} + int(C_1))$ is empty. This is because otherwise, we would obtain a contradiction $z - \bar{z} \in (int(C_1) - int(C_1)) \cap \mathbb{Z}^n = int(C_1 - C_1) \cap \mathbb{Z}^n = int(\lambda_1 C) \cap \mathbb{Z}^n = \{0\}$. Therefore, we have the following

$$\operatorname{vol}(M_q^n + C_1) = (2q+1)^n \operatorname{vol}(C_1) = (2q+1)^n \left(\frac{\lambda_1}{2}\right)^n \operatorname{vol}(C).$$
 (1.5)

We will now show that for $1 \le i \le n-1$, the following inequality holds

$$\operatorname{vol}\left(M_{q}^{n}+C_{i+1}\right) \geq \left(\frac{\lambda_{i+1}}{\lambda_{i}}\right)^{n-i} \operatorname{vol}\left(M_{q}^{n}+C_{i}\right).$$
(1.6)

Thus, we can assume that $\lambda_{i+1} > \lambda_i$ and consider two lattice points z and \bar{z} in \mathbb{Z}^n that differ in the last n-i coordinates, i.e., $(z_{i+1}, \ldots, z_n) \neq (\bar{z}_{i+1}, \ldots, \bar{z}_n)$. It follows that

$$[\boldsymbol{z} + \operatorname{int}(C_{i+1})] \cap [\bar{\boldsymbol{z}} + \operatorname{int}(C_{i+1})] = \emptyset.$$
(1.7)

If this is not the case, then the i + 1 linearly independent lattice points $z - \bar{z}, z_1, \ldots, z_i$ would belong to the interior of $\lambda_{i+1}C$, contradicting the minimality of λ_{i+1} . Hence, from equation (1.7), we obtain

$$\operatorname{vol}(M_{q}^{n} + C_{i+1}) = (2q+1)^{n-i} \operatorname{vol}(M_{q}^{i} + C_{i+1}),$$
$$\operatorname{vol}(M_{q}^{n} + C_{i}) = (2q+1)^{n-i} \operatorname{vol}(M_{q}^{i} + C_{i}).$$

To verify equation (1.6), it suffices to show

$$\operatorname{vol}\left(M_{q}^{i}+C_{i+1}\right) \geq \left(\frac{\lambda_{i+1}}{\lambda_{i}}\right)^{n-i} \operatorname{vol}\left(M_{q}^{i}+C_{i}\right).$$
(1.8)

Let f_1 and f_2 be linear maps from \mathbb{R}^n to \mathbb{R}^n defined as

$$f_1(\mathbf{x}) = \left(\frac{\lambda_{i+1}}{\lambda_i} x_1, \dots, \frac{\lambda_{i+1}}{\lambda_i} x_i, x_{i+1}, \dots, x_n\right)^{\top},$$

$$f_2(\mathbf{x}) = \left(x_1, \dots, x_i, \frac{\lambda_{i+1}}{\lambda_i} x_{i+1}, \dots, \frac{\lambda_{i+1}}{\lambda_i} x_n\right)^{\top}.$$

Using the fact that $M_q^i + C_{i+1} = f_2 \left(M_q^i + f_1 (C_i) \right)$, we have

$$\operatorname{vol}\left(M_{q}^{i}+C_{i+1}\right)=\left(\frac{\lambda_{i+1}}{\lambda_{i}}\right)^{n-i}\operatorname{vol}\left(M_{q}^{i}+f_{1}\left(C_{i}\right)\right).$$

To prove equation (1.8), we have to show

$$\operatorname{vol}\left(M_{q}^{i}+f_{1}\left(C_{i}\right)\right)\geq\operatorname{vol}\left(M_{q}^{i}+C_{i}\right).$$
(1.9)

Let L_i^{\perp} be the (n - i)-dimensional orthogonal complement of L_i , and it is obvious that for any $\mathbf{x} \in L_i^{\perp}$, there exists a $t(\mathbf{x}) \in L_i$ such that $C_i \cap (\mathbf{x} + L_i) \subset$ $(f_1(C_i) \cap (\mathbf{x} + L_i)) + t(\mathbf{x})$. Thus, we have

$$\left(M_q^i+C_i\right)\cap(\mathbf{x}+L_i)\subset\left[\left(M_q^i+f_1(C_i)\right)\cap(\mathbf{x}+L_i)\right]+t(\mathbf{x}).$$

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From this, we obtain

$$\operatorname{vol}(M_q^i + C_i) = \int_{\boldsymbol{x} \in L_i^{\perp}} \operatorname{vol}_i((M_q^i + C_i) \cap (\boldsymbol{x} + L_i)) d\boldsymbol{x}$$
$$\leq \int_{\boldsymbol{x} \in L_i^{\perp}} \operatorname{vol}_i((M_q^i + f_1(C_i)) \cap (\boldsymbol{x} + L_i)) d\boldsymbol{x}$$
$$= \operatorname{vol}(M_q^i + f_1(C_i)),$$

where $vol_i(\cdot)$ denotes the *i*-dimensional volume. This implies (1.9), so we have proven (1.6). Finally, from (1.4), (1.5) and (1.6) we have

$$(2q+\gamma)^{n} \ge \operatorname{vol}\left(M_{q}^{n}+C_{n}\right) \ge \left(\frac{\lambda_{n}}{\lambda_{n-1}}\right) \operatorname{vol}\left(M_{q}^{n}+C_{n-1}\right)$$
$$\ge \left(\frac{\lambda_{n}}{\lambda_{n-1}}\right) \left(\frac{\lambda_{n-1}}{\lambda_{n-2}}\right)^{2} \operatorname{vol}\left(M_{q}^{n}+C_{n-2}\right) \ge \cdots \cdots$$
$$\ge \left(\frac{\lambda_{n}}{\lambda_{n-1}}\right) \cdot \left(\frac{\lambda_{n-1}}{\lambda_{n-2}}\right)^{2} \cdot \cdots \cdot \left(\frac{\lambda_{2}}{\lambda_{1}}\right)^{n-1} \operatorname{vol}\left(M_{q}^{n}+C_{1}\right)$$
$$= \lambda_{n} \cdot \cdots \cdot \lambda_{1} \frac{\operatorname{vol}(C)}{2^{n}} (2q+1)^{n}$$

and so

$$\lambda_1 \cdot \ldots \cdot \lambda_n \operatorname{vol}(C) \leq 2^n \cdot \left(\frac{2q+\gamma}{2q+1}\right)^n.$$

Since this inequality holds for all $q \in \mathbb{N}$, the theorem is proven.

1.7 Number of lattice points in a convex body

and successive minima

Henk [31] used Minkowski's second theorem to give an upper bound on the number of lattice points in an origin symmetric convex body *C* in terms of the successive minima in the following way:

Theorem 4 (Henk [31]). Let $n \ge 2$, $C \in \mathscr{K}_0^n$ and $\Lambda \subset \mathbb{R}^n$ be a full-dimensional *lattice then*

$$|C \cap \Lambda| \le 2^{n-1} \prod_{i=1}^n \left(\frac{2}{\lambda_i(C,\Lambda)} + 1\right).$$

For completeness we include a proof of this result as given in [31].

Proof. Without loss of generality, let $\Lambda = \mathbb{Z}^n$, and we can assume that (cf. 1.2 and 1.3),

$$\operatorname{int}(\lambda_i(C,\mathbb{Z}^n)C) \cap \mathbb{Z}^n \subset \operatorname{lin}\{\mathbf{0}, \mathbf{e}_1, \dots, \mathbf{e}_{i-1}\} \cap \mathbb{Z}^n, \quad 1 \le i \le n.$$
(1.10)

Let's denote $q_i = \left\lfloor \frac{2}{\lambda_i(C,\Lambda)} + 1 \right\rfloor$, where $1 \le i \le n$, for abbreviation. We need to determine *n* numbers $u_i \in \mathbb{N}$ such that

$$u_n = q_n, \quad q_i \le u_i < 2q_i, \quad \text{and} \quad u_{i+1} \text{ divides } u_i, \quad 1 \le i \le n-1.$$
 (1.11)

To determine u_k , we consider two cases, assuming that we have already found u_n, \ldots, u_{k+1} with the properties mentioned earlier. If $u_{k+1} \ge q_k$, then we simply set $u_k = u_{k+1}$. This satisfies $q_k \le u_k = u_{k+1} < 2q_{k+1} \le 2q_k$, since $q_k \ge q_{k+1}$. Otherwise, if $u_{k+1} < q_k$. We can express q_k as $q_k = m \cdot u_{k+1} + r$, where $m \in \mathbb{N}$, $m \ge 1$, and $0 \le r < u_{k+1}$. In this case, we set $u_k = q_k + u_{k+1} - r$, and it can be easily verified that u_k satisfies the requirements of (1.11). Let $\widetilde{\Lambda} \subset \mathbb{Z}^n$ be the lattice generated by the vectors $u_1e_1, u_2e_2, \ldots, u_ne_n$. It follows that det $\widetilde{\Lambda}/\det \Lambda = u_1 \cdot u_2 \cdot \ldots \cdot u_n$, and combining this with the upper bounds on the values of u_i as given in Lemma 2.1, in [31] we obtain

$$|C \cap \Lambda| \le |2C \cap \widetilde{\Lambda}| \prod_{i=1}^{n} u_i < |2C \cap \widetilde{\Lambda}| 2^{n-1} \prod_{i=1}^{n} \left\lfloor \frac{2}{\lambda_i(C,\Lambda)} + 1 \right\rfloor.$$

To verify the theorem, it is sufficient to show that $2C \cap \widetilde{\Lambda} = \{0\}$. Let's assume that there exists a non-zero element $g \in 2C \cap \widetilde{\Lambda}$, and let *k* be the largest index

where **g** has a non-zero coordinate, i.e., $g_k \neq 0$, while $g_{k+1} = \cdots = g_n = 0$. Then we can express **g** as

$$g = z_1(u_1e_1) + z_2(u_2e_2) + \dots + z_k(u_ke_k) \in 2C,$$

where $z_i \in \mathbb{Z}$. Since u_k is a divisor of u_1, \ldots, u_{k-1} , and we have $2/u_k < \lambda_k(C, \mathbb{Z}^n)$ (1.11), we can deduce that

$$\frac{1}{u_k} \mathbf{g} \in \left(\frac{2}{u_k} C\right) \cap \mathbb{Z}^n \subset \operatorname{int}(\lambda_k(C, \mathbb{Z}^n) C) \cap \mathbb{Z}^n.$$

However, this violates condition (1.10) since $g_k \neq 0$. Therefore, we conclude that $2C \cap \widetilde{\Lambda} = \{0\}$, proving the theorem.

1.8 The generalised Hermite constant

The generalised Hermite constant $\gamma_{r,s}$, as defined by Rankin [39], is a positive real number defined as the infimum over all γ such that every lattice Λ of rank r in \mathbb{R}^r has a sublattice Γ of rank s and determinant

$$\det(\Gamma) \le \gamma^{1/2} (\det(\Lambda))^{s/r}. \tag{1.12}$$

In other words, $\gamma_{r,s}$ is the smallest constant for which this inequality holds for all lattices Λ of rank r and for all choices of sublattice Γ of rank s.

The generalised Hermite constant is an important constant in lattice theory and has applications in number theory, cryptography, and coding theory. The value of $\gamma_{r,s}$ is known exactly for some small values of r and s, but it is generally difficult to compute or estimate for larger values. We refer the reader to the papers [43, 50] for known results on the Rankin constant.

Note that in (1.12) $\gamma_{r,1} = \gamma_{r,r-1} = \gamma_r$ represents the ordinary Hermite constant [18].

Chapter 2

Compressed sensing for integer valued signals

2.1 Introduction

This chapter gives a brief introduction to the mathematical aspects of compressed sensing relevant to this thesis. It is based on the references [23, 24, 26, 33, 37, 42].

The main objective of compressed sensing is to design and efficiently solve underdetermined systems of linear equations

$$A\mathbf{x} = \mathbf{b},\tag{2.1}$$

with $A \in \mathbb{R}^{m \times n}$ and $b \in \mathbb{R}^{m}$, m < n, under an additional assumption that the solution $\mathbf{x} \in \mathbb{R}^{n}$ is "sparse". Here a solution vector \mathbf{x} is "sparse" or, more specifically, *s*-sparse if the number of its non-zero elements is less than or equal to *s*. In other words, $\|\mathbf{x}\|_{0} \leq s$. The matrix *A* has been traditionally referred to as a *measurement matrix*, \boldsymbol{b} as a vector of *m* measurements and a solution vector \boldsymbol{x} as a signal. In a basic scenario, there is an unknown *s*-sparse signal \boldsymbol{x} and we want to know what conditions for the measurement matrix *A* will allow us to uniquely and efficiently recover \boldsymbol{x} from the measurements \boldsymbol{b} . That is, in particular, \boldsymbol{x} must be the unique *s*-sparse solution to the system (2.1).

The feasibility of this problem in various practical settings is by now welljustified, and sparsity of data has been identified as a new paradigm in signal processing. Various sufficient conditions for precise recovery of the signal xhave been found in terms of the properties of the measurement matrix Aand the sparsity of x. We refer the reader to [19] for an extensive survey. In the following section, we briefly discuss probably the most studied recovery approach.

2.2 Unique recovery using basis pursuit

A natural way to recover sparse signals from an underdetermined linear system would be to solve an ℓ_0 -minimisation problem, that is the problem

$$\min \|\mathbf{x}\|_0 \quad \text{subject to} \quad A\mathbf{x} = \mathbf{b}. \tag{2.2}$$

This problem, however, is known to be NP-hard [36], which makes solving it unpractical. A popular and by now well-understood approach to avoid the NP-hardness barrier is to replace (2.2) with a convex optimisation problem

$$\min \|\boldsymbol{x}\|_1 \quad \text{subject to} \quad A\boldsymbol{x} = \boldsymbol{b}$$



Figure 2.1: The 1-sparse vector x_0 can be uniquely recovered via ℓ_1 -minimisation since the null space of *A* intersects the cone only at the origin.

which is known as *basis pursuit* [16]. As this problem is convex, it can be solved efficiently using convex optimisation methods. A necessary and sufficient condition under which a signal \mathbf{x}_0 is uniquely recovered by basis pursuit is given as follows: The set of all feasible solutions, $\mathbf{x}_0 + \ker(A)$, intersects with the set $\{\mathbf{x} : ||\mathbf{x}||_1 \le ||\mathbf{x}_0||_1\}$ exactly at \mathbf{x}_0 . This condition provides a useful geometric intuition about properties of measurement matrices to ensure uniqueness of the solution. One of those properties is the so-called Null Space Property (NSP). To describe it we need to introduce a useful notation. Let $M \in \mathbb{R}^{m \times n}$ and let $\tau = \{i_1, \ldots, i_k\} \subseteq \{1, \ldots, n\}$ with $i_1 < i_2 < \cdots < i_k$. We will denote by M_{τ} the $m \times k$ submatrix of M with columns indexed by τ . In the same manner, given $\mathbf{x} \in \mathbb{R}^n$, we will denote by \mathbf{x}_{τ} the vector $(x_{i_1}, \ldots, x_{i_k})^{\top}$. The complement of τ in $\{1, \ldots, n\}$ will be denoted as $\bar{\tau}$.



Figure 2.2: When the ℓ_1 ball intersects with the solution set of $A\mathbf{x} = \mathbf{b}$ in a single point (left), then ℓ_1 minimisation can provide the same solution as ℓ_0 minimisation.

Now the NSP is given by

$$\ker(A) \cap \left\{ w \in \mathbb{R}^n : \|w_{\tau}\|_1 \ge \|w_{\bar{\tau}}\|_1 \right\} = \{\mathbf{0}\},\$$

where $\tau = \operatorname{supp}(\boldsymbol{x}_0)$. See Figure 2.1.

It is well-known that, if *A* fulfills the NSP with respect to some subset $\tau \subset [n]$, every signal \mathbf{x}_0 supported on τ is the unique minimiser of (2.2) with $\mathbf{b} = A\mathbf{x}_0$ (cf. [17, 20, 21, 22, 24]). Figure 2.2 illustrates this concept.

By using random matrices *A* such as a matrix with independent and identically distributed Gaussian entries, it is possible to achieve a very high probability of *A* having the NSP and therefore of (2.2) to succeed, given that the number of measurements satisfies $m \ge Cs \log(n)$, where *s* is the sparsity of the signal x_0 and *C* some positive constant not depending on *s* and *n* [11].

2.3 Unique recovery of integer signals

In some applications, additional information on the signal *x* is known and one can say that the signal belongs to a certain *signal space S* that possesses certain additional structure. Hence, we face a natural question: how the structure of *S* can be exploited to improve recovery guarantees? In this thesis, we will focus on the integer-valued signals, which appear, for instance, in MIMO systems [41], wideband spectrum sensing [7] and error correcting codes [12]. We will work in the setting introduced by Fukshansky, Needell, and Sudakov [26].

Further, in this thesis we will consider the signals with bounded entries. We will focus on the conditions that guarantee the uniqueness of the solution.

Let us fix a set $S \subset \mathbb{Z}^n$ which will be our *signal space*. Fukshansky, Needell, and Sudakov [26] introduced and studied the problem of unique recovery of a signal $\mathbf{x}_0 \in S$ from a relatively small number of noisy linear integer measurements. Specifically, given a number of measurements m with m < n, we aim to construct an integer measurement matrix $A \in \mathbb{Z}^{m \times n}$ such that any signal $\mathbf{x}_0 \in S$ can be uniquely recovered from m measurements represented by the vector $\mathbf{b} \in \mathbb{R}^m$ of the form

$$\boldsymbol{b} = A\boldsymbol{x}_0 + \boldsymbol{e} \tag{2.3}$$

with an unknown *noise vector* $e \in \mathbb{R}^m$. To allow unique recovery we assume that

$$\|\boldsymbol{e}\|_2 < c\,,$$

where $\|\cdot\|_2$ denotes the ℓ_2 -norm and *c* is a suitably chosen constant. Note that, the number of measurements *m* is independent of the choice of the constant *c*.

Based on [26], we will use the following recovery approach. For a set $Q \subset \mathbb{Z}^n$ we denote by $\mathscr{R}(Q)$ the set of all integer matrices *A* such that

$$||Ay||_2 \ge 1 \text{ for any nonzero } y \in Q.$$
(2.4)

Recall that the *difference set* D(X) of a set $X \in \mathbb{R}^n$ consists of all points x - ywith $x, y \in X$. We set c = 1/2 and consider measurement matrices $A \in \mathscr{R}(Q)$ with Q = D(S). In this case, for any $e \in \mathbb{R}^m$ with $||e||_2 < c = 1/2$, the signal x_0 is the unique point of S satisfying the bound

$$||A\mathbf{x} - \mathbf{b}||_2 \le \frac{1}{2}.$$
 (2.5)

Indeed, for any $x \in S$, $x \neq x_0$, satisfying (2.5), we would have

$$||A(\mathbf{x} - \mathbf{x}_0)||_2 = ||A\mathbf{x} - \mathbf{b} - (A\mathbf{x}_0 - \mathbf{b})||_2$$
$$= ||A\mathbf{x} - \mathbf{b} + \mathbf{e}||_2$$
$$\leq ||A\mathbf{x} - \mathbf{b}||_2 + ||\mathbf{e}||_2 < 1,$$

contradicting (2.4). Therefore, x_0 can be recovered by any algorithm that, given input $b \in \mathbb{R}^m$ computes a vector $x \in S$ satisfying (2.5).

By \mathbb{Z}_l^n we will denote the set of *l*-sparse *n*-dimensional integer vectors:

$$\mathbb{Z}_l^n = \left\{ \boldsymbol{z} \in \mathbb{Z}^n : \|\boldsymbol{z}\|_0 \le l \right\}.$$

Given positive integers *n*, *r*, we denote by $C^n(r)$ the *n*-dimensional cube defined as $C^n(r) = \{ x \in \mathbb{R}^n : ||x||_{\infty} \le r \}.$

We will be interested in unique recovery of *l*-sparse signals with entries from a finite integer alphabet $[-r, r] \cap \mathbb{Z}$. Specifically, we will work with the signal space

$$S_l^n(r) = C^n(r) \cap \mathbb{Z}_l^n,$$
where 2l < n.

The space $S_l^n(r)$ is finite and hence allows using a single measurement for unique recovery of its signals. From the computational and error-correcting perspectives (see [26] for more details), the measurement should have as small as possible absolute integer entries. Hence, given $l, n \in \mathbb{Z}$ with $1 \le l <$ $n/2, r \in \mathbb{Z}_{>0}$ and letting $Q = D(S_l^n(r))$, we face the optimisation problem

$$\min\{\|H\|_{\infty} : H \in \mathbb{Z}^{1 \times n}, H \in \mathcal{R}(Q)\}.$$
(2.6)

In Chapter 4 we obtain general estimates for the minimum in (2.6).

2.4 Recent results on unique recovery over \mathbb{Z}_l^n

The papers [26, 25, 33, 32] consider the problem of unique recovery over the signal space \mathbb{Z}_l^n , where *l* is a positive integer with 2l < n. In this setting, the difference set $D(\mathbb{Z}_l^n)$ consists of 2*l*-sparse integer vectors, $D(\mathbb{Z}_l^n) = \mathbb{Z}_{2l}^n$. The unique recovery of signals from \mathbb{Z}_l^n involves constructing matrices $A \in \mathbb{Z}^{m \times n}$ with m = 2l and as large as possible *n*, that belong to $\mathscr{R}(\mathbb{Z}_m^n)$. From the computational and error-correcting perspectives (see [26] for more details), it is also desirable to fix or bound the maximum absolute entry $||A||_{\infty}$ of the matrix *A*.

Konyagin [32, Theorem 3] proved the following theorem.

Theorem 5. For integers $k, m \ge 2$, and integer n with

$$m < n \le \frac{c^{-m}k^{m/(m-1)}}{\log(k)}$$

there exists an integer $m \times n$ matrix $A \in \mathscr{R}(\mathbb{Z}_m^n)$ such that $||A||_{\infty} = k$, where c is an absolute constant.

The proof of Theorem 5 employs probabilistic arguments to show existence of the desired measurement matrices. Subsequently, Konyagin and Sudakov [33, Theorem 1.3] (see also Ryutin [42]) proved the following result using an explicit and easily computable construction.

Theorem 6. Let $k \in \mathbb{Z}_{>0}$, $m \in \mathbb{Z}_{>0}$, $m \ge 2$, and

$$m < n \le \max(k+1, k^{m/(m-1)}/2).$$
 (2.7)

Then there exists an $m \times n$ integer matrix $A \in \mathscr{R}(\mathbb{Z}_m^n)$ such that $||A||_{\infty} \leq k$.

Theorem 6 implies the following result.

Corollary 7. For any given $m, n \in \mathbb{Z}_{>0}$, $2 \le m < n$ there exists an $m \times n$ integer matrix $A \in \mathscr{R}(\mathbb{Z}_m^n)$ with

$$||A||_{\infty} \leq \Delta(m, n) = \min(n - 1, (2n)^{(m-1)/m}).$$

We will use Corollary 7 in the proofs of our results given in Chapter 4.

Chapter 3

Main tools: consolidation of linear Diophantine equations and Siegel's lemma

3.1 Consolidation/aggregation of linear Diophantine equations

The proof of one of the main results of this thesis (see Theorem 18) is based on consolidation/aggregation of linear Diophantine equations. This topic has been extensively studied in the literature. We refer the reader to the papers [38, 40].

Let $D \subset \mathbb{Z}^n$ be a set of integer points, $A \in \mathbb{Z}^{m \times n}$, $2 \le m < n$, be a matrix of rank rank(A) = m, and $b \in \mathbb{Z}^m$.

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Let $B \in \mathbb{Z}^{l \times m}$, l < m, be a matrix of rank l such that

$${x \in D : (BA)x - Bb = 0} = {x \in D : Ax - b = 0}$$

Following [40], we will call B an m-into-l consolidating matrix and (BA)x –

 $B\mathbf{b} = \mathbf{0}$ an *m*-into-*l* consolidation of $A\mathbf{x} - \mathbf{b} = \mathbf{0}$ with respect to the set *D*.

Let further $C \in \mathbb{Z}^{m \times (m-l)}$ be an integer matrix of rank(C) = m-l such that, for some consolidating matrix B, the columns of C span the *kernel* ker(B) = $\{x \in \mathbb{R}^m : Bx = 0\}$. That is, denoting by span_R(C) the subspace spanned by the columns of C, we have span_R(C) = ker(B). We will call C an *aggregating matrix* for Ax - b = 0 with respect to the set D.

We will write

$$F(\mathbf{x}) = A\mathbf{x} - \mathbf{b}$$
,

and denote by $F_i(x)$ the *i*th entry of the vector F(x), that is

$$F(\mathbf{x}) = (F_1(\mathbf{x}), \dots, F_m(\mathbf{x}))^\top.$$

Consider the set

$$F^{\circ} = \{F(x) : x \in D\} = \{Ax : x \in D\} - b$$

This is the image of *D* under the linear mapping determined by the matrix *A* translated by the vector $-\mathbf{b}$. The following well-known lemma describes very important properties of the consolidation/aggregation.

Lemma 8. Let $B \in \mathbb{Z}^{l \times m}$, l < m, be a matrix of rank l and $C \in \mathbb{Z}^{m \times (m-l)}$ be a matrix of rank m-l. Then

(i) *B* an *m*-into-*l* consolidating matrix for $A\mathbf{x} - \mathbf{b} = \mathbf{0}$ if and only if $F^{\circ} \cap \ker(B) \subset \{\mathbf{0}\}$.

(ii) *C* is an aggregating matrix of $A\mathbf{x} - \mathbf{b} = \mathbf{0}$ if and only if $F^{\circ} \cap \operatorname{span}_{\mathbb{R}}(C) \subset \{\mathbf{0}\}$.

We will need the following lemma, given in [6, Theorem 6]. For completeness, we include a proof of this result as given in [40, Example 4.1].

Lemma 9. Assume that $q_i \in \mathbb{Z}$ satisfy $|F_i(\mathbf{x})| < q_i$ for every $\mathbf{x} \in D$ such that $F_i(\mathbf{x}) = \cdots = F_{i-1}(\mathbf{x}) = 0$, $i = 1, \dots, m-1$. Then $F_1(\mathbf{x}) + q_1F_2(\mathbf{x}) + q_1q_2F_3(\mathbf{x}) + \cdots + q_1 \cdots q_{m-1}F_m(\mathbf{x}) = 0$ is an m-into-1 consolidation of $F(\mathbf{x}) = \mathbf{0}$.

Proof. We have to show that $B = (1, q_1, ..., q_1 \cdots q_{m-1})$ is an *m*-into-1 consolidating matrix for $F(\mathbf{x}) = \mathbf{0}$. Let $C = (c_{ij}) \in \mathbb{Z}^{m \times (m-1)}$ be defined by $c_{1,k} = q_1 \delta_{1,k}$ for k = 1, ..., m-1 and $c_{ij} = q_i \delta_{i,j} - \delta_{i-1,j}$ for i = 2, ..., m, j = 1, ..., m-1. Here $\delta_{i,j}$ stands for the Kronecker delta. That is

$$C = \begin{pmatrix} q_1 & 0 & \cdots & 0 & 0 \\ -1 & q_2 & \cdots & 0 & 0 \\ 0 & -1 & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & -1 & q_{m-1} \\ 0 & 0 & \cdots & 0 & -1 \end{pmatrix}.$$

We will show that *C* is an aggregating matrix for $F(\mathbf{x}) = \mathbf{0}$. It is sufficient to check that the inclusion $F^{\circ} \cap \operatorname{span}_{\mathbb{R}}(C) \subset \{\mathbf{0}\}$ in the part (ii) of Lemma 8 holds. Suppose

$$C\mathbf{v} = F(\mathbf{x}) \tag{3.1}$$

for some $v \in \mathbb{R}^{m-1}$ and $x \in D$.

Observe that the greatest common divisor of $m \times m$ subdeterminants of *C* is equal to one. It follows that the columns of *C* form a basis of the lattice

span_R(C) || \mathbb{Z}^{-1} . Next, by (3.1) we have $C \mathbf{v} \in \mathbb{Z}^{-1}$. Therefore, $\mathbf{v} \in \mathbb{Z}^{-1}$. The first coordinate of $C \mathbf{v}$ is $q_1 v_1$, hence $|q_1 v_1| = |F_1(\mathbf{x})| < |q_1|$ implies $v_1 = F_1(\mathbf{x}) = 0$. The second coordinate of $C \mathbf{v}$ is $q_2 v_2 - v_1 = q_2 v_2 = F_2(\mathbf{x})$. Now by definition of q_2 we have $|q_2 v_2| = |F_2(\mathbf{x})| < |q_2|$ and, consequently, $v_2 = 0$. Proceeding in this way we get $\mathbf{v} = \mathbf{0}$.

Finally, it is easy to see that the columns of *C* span ker(*B*). Hence *B* is an *m*-into-1 consolidating matrix for F(x) = 0.

3.2 Siegel's lemma

Consider a system of homogeneous linear equations

$$\begin{cases} a_{11}x_1 + \dots + a_{1n}x_n = 0, \\ \vdots \\ a_{m1}x_1 + \dots + a_{mn}x_n = 0. \end{cases}$$
(3.2)

Let $A = (a_{ij}) \in \mathbb{Z}^{m \times n}$, m < n. We can write the system (3.2) in the form:

$$A\mathbf{x} = \mathbf{0}.\tag{3.3}$$

Since m < n, the system (3.3) must have a non-trivial solution in integers. If the entries of *A* are small integers, then one can expect that there will be a solution in relatively small integers. This idea was applied by Thue in [49] to Diophantine approximations. Subsequently, Siegel [47, Bd. I, p. 213, Hilfssatz] was the first to state this principle formally.

Recall that we denote by $||A||_{\infty}$ the maximum absolute value of an entry of *A*, that is $||A||_{\infty} = \max_{ij} |a_{ij}|$. In what follows, without loss of generality we may assume that *A* has rank *m*. Following Siegel's work, Schmidt [45] states the following result.

Theorem 10 (Siegel's Lemma). *The system* (3.3) *has a solution* $x \in \mathbb{Z}^n$ *with*

$$0 < \|\mathbf{x}\|_{\infty} < 1 + (n\|A\|_{\infty})^{m/(n-m)}.$$
(3.4)

The exponent m/(n-m) on the right hand side of (3.4) is optimal.

In general, Siegel's lemma refers to a family of results that link the properties of lattice points in the kernel space of the matrix *A* and the characteristics of the matrix, such as its maximum absolute entry or certain functions of its subdeterminants.

Bombieri and Vaaler [9] proved, by using geometry of numbers, the following version of Siegel's lemma which involves n-m solutions to the system (3.3).

Theorem 11 (Bombieri and Vaaler [9]). The system (3.3) has n - m linearly independent integer solutions $\mathbf{x}_1, \ldots, \mathbf{x}_{n-m} \in \mathbb{Z}^n$, with

$$\prod_{i=1}^{n-m} \|\boldsymbol{x}_i\|_{\infty} \leq \frac{\sqrt{\det(AA^T)}}{\gcd(A)},$$

where gcd(A) is the greatest common divisor of all $m \times m$ subdeterminants of A.

In this thesis, Siegel's lemma in the form of Theorem 11 plays an important role in the proof of Theorem 19. In what follows we will give a short overview of selected results related to Siegel's lemma in line with a recent paper [4].

Recall that $C^n(1) = [-1, 1]^n$ and let $\ker(A) = \{x \in \mathbb{R}^n : Ax = 0\}$. Consider the section $S(A) = C^n(1) \cap \ker(A)$ of the cube $C^n(1)$ and the lattice $\Lambda(A) = \mathbb{Z}^n \cap \ker(A)$. The lattice $\Lambda(A)$ has determinant $\det(\Lambda(A)) = \sqrt{\det(AA^T)} / \gcd(A)$. The (n - m)-dimensional subspace $\ker(A)$ can be considered as a usual Euclidean (n-m)-dimensional space. This immediately extends the definition of successive minima to *o*-symmetric bounded convex sets with nonempty relative interior in ker(A) and (n - m)-dimensional lattices in ker(A). Theorem 11 is an immediate corollary of the following result.

Theorem 12. Let $A \in \mathbb{Z}^{m \times n}$, m < n, be an integral matrix of rank m. Then the inequality

$$\prod_{i=1}^{n-m} \lambda_i(S(A), \Lambda(A)) \le \det(\Lambda(A))$$
(3.5)

holds.

Proof of Theorem 12: Using a result of Vaaler [51], we have $vol_{n-m}(S(A)) \ge 2^{n-m}$. Hence, Minkowski's theorem on successive minima gives

$$\prod_{i=1}^{n-m} \lambda_i(S(A), \Lambda(A)) \le \frac{2^{n-m} \det(\Lambda(A))}{\operatorname{vol}_{n-m}(S(A))} \le \det(\Lambda(A)).$$

In what follows, we will discuss the special case m = 1, that is when *A* is just an *n*-dimensional nonzero row vector. Refinements for this case imply a slight improvement of the bound (4.13).

Theorem 11 implies that for every nonzero vector \mathbf{a} in \mathbb{Z}^n , $n \ge 2$, there exists a vector \mathbf{x} in \mathbb{Z}^n , such that

$$a \cdot x = 0, \quad 0 < \|x\|_{\infty}^{n-1} \le \sqrt{n} \|a\|_{\infty}.$$
 (3.6)

The exponent n-1 in the latter bound is optimal. Let us define

$$c(n) = \sup_{a \in \mathbb{Z}^n \setminus \{0\}} \inf_{\substack{x \in \mathbb{Z}^n \setminus \{0\} \\ a \cdot x = 0}} \frac{\|x\|_{\infty}^{n-1}}{\|a\|_{\infty}}.$$

That is c(n) is the optimal constant in the bound (3.6).

It is easy to see that c(2) = 1. Further, the equality c(3) = 4/3 is implicit in [15]. Namely, the inequality $c(3) \le 4/3$ is contained in [15, Lemma 4], while the inequality $c(3) \ge 4/3$ is a consequence of [15, Lemma 7]. We have also c(4) = 27/19. The inequality $c(4) \ge 27/19$ was proved by Chaladus in [14] and its counterpart $c(4) \le 27/19$ was obtained in [2] (see also [44]). For n > 4, the exact values of the constants c(n) remain unknown.

A general result that gives a geometric interpretation for the constant c(n) was given by Schinzel [44]. Given an origin-symmetric convex body K we denote by $\Delta(K)$ its critical determinant.

Theorem 13 (Schinzel [44]). For $n \ge 3$

$$c(n) = \sup \Delta(H_{\alpha_1,\ldots,\alpha_{n-3}})^{-1},$$

where $H_{\alpha_1,...,\alpha_{n-3}}$ is a generalised hexagon given by

$$H_{\alpha_1,...,\alpha_{n-3}} = \left\{ \boldsymbol{x} \in \mathbb{R}^{n-1} : \|\boldsymbol{x}\|_{\infty} \le 1, \left| \sum_{i=1}^{n-3} \alpha_i x_i + x_{n-2} + x_{n-1} \right| \le 1 \right\}$$

and the supremum is taken over all rational numbers $\alpha_1, \ldots, \alpha_{n-3}$ in the interval (0, 1].

Based on the values of c(n) for $n \le 4$, the following conjecture was proposed in [2].

Conjecture 14. The equality

$$c(n) = \Delta(H_{1,\dots,1})^{-1}$$

holds. Here $H_{1,...,1}$ is a generalised hexagon in \mathbb{R}^{n-1} .

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The bound in Theorem (11) immediately implies

$$c(n) \le \sqrt{n} \,. \tag{3.7}$$

In [2], the constant c(n) was estimated as

$$c(n) \le \sigma_n^{-1},\tag{3.8}$$

where σ_n is the sinc integral

$$\sigma_n = \frac{2}{\pi} \int_0^\infty \left(\frac{\sin t}{t}\right)^n dt \, .$$

The bound (3.8) asymptotically improves on (3.7) with factor $\sqrt{\pi/6}$. The numbers σ_n are rational, the sequences of numerators and denominators of $\sigma_n/2$ can be found in [48] (sequences A049330 and A049331).

Chapter 4

Unique recovery of finite-valued integer signals

This chapter presents the main results of the thesis, published in [1].

4.1 Unique recovery of sparse bounded integer signals

In this thesis, we will obtain general estimates for the minimum in (2.6). Using condition (2.4), to get an upper bound for (2.6), it is sufficient to find an $1 \times n$ measurement matrix H such that its kernel space does not share any nonzero integer points with the convex hull of $D(S_l^n(r))$. This straightforward approach, however, results in the estimate

$$||H||_{\infty} = O(r^{n-1}), \tag{4.1}$$

where the implicit constant depends on l and n.

The first result of this thesis shows that the exponent n - 1 in (4.1) can be replaced with 2l - 1. Let

$$p_{m,n}(r) = (m\Delta(m,n)r+1)^{m-1} + \Delta(m,n) \sum_{i=0}^{m-2} (m\Delta(m,n)r+1)^i,$$

where $\Delta(m, n) = \min(n - 1, (2n)^{(m-1)/m})$.

Theorem 15. For any $l, n \in \mathbb{Z}$ with $1 \le l < n/2$ and $r \in \mathbb{Z}_{>0}$ there exists an $1 \times n$ integer matrix H such that $H \in \mathscr{R}(Q)$ with $Q = D(S_l^n(r))$ and

$$\|H\|_{\infty} < p_{2l,n}(2r). \tag{4.2}$$

The proof of Theorem 15 is constructive. To obtain the bound (4.2) we combine known results on unique recovery over \mathbb{Z}_l^n , outlined in Section 2.4, with aggregation techniques, outlined in Section 3.1.

The second result gives a lower bound for the minimum in (2.6). Notably, it shows that the polynomial $p_{2l,n}(2r)$ in (4.2) cannot be replaced by a polynomial in r with degree smaller than 2l - 1.

Theorem 16. For any $l, n \in \mathbb{Z}$ with $1 \le l < n/2$, $r \in \mathbb{Z}_{>0}$ and $1 \times n$ integer matrix $H \in \mathscr{R}(Q)$ with $Q = D(S_l^n(r))$ the bound

$$\|H\|_{\infty} > \frac{r^{2l-1}}{\sqrt{2l}} \tag{4.3}$$

holds.

Based on Theorems 15 and 16 we pose the following question. Let us fix the sparsity level l and dimension n. In this setting, it would be interesting to find optimal upper bounds for minimal $||H||_{\infty}/r^{2l-1}$ when r tends to infinity. Specifically, given $l, n \in \mathbb{Z}_{>0}$ with $1 \le l < n/2$, to estimate

$$c_1(l,n) = \limsup \inf \frac{\|H\|_{\infty}}{r^{2l-1}},$$
 (4.4)

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where the supremum limit is taken over all positive integers r and the infimum is taken over all $1 \times n$ integer matrices $H \in \mathscr{R}(Q)$ with $Q = D(S_l^n(r))$. Theorems 15 and 16 give a large interval for values of this quantity

$$\frac{1}{\sqrt{2l}} \le c_1(l,n) \le (4l\Delta(2l,n))^{2l-1}$$

Although for finite signal spaces a single measurement is sufficient for unique recovery, one can ask whether allowing extra measurements would result in reducing measurements' entries. In this vein, we obtain the following general result. Let $Q \subset \mathbb{Z}^n$ be an arbitrary set. Suppose that we have an $1 \times n$ matrix $H \in \mathscr{R}(Q)$. We show that, for any integer m with 1 < m < n, there exists an $m \times n$ matrix $A = (a_{ij})$ such that $A \in \mathscr{R}(Q)$ and the maximum absolute entry $||A||_{\infty} = \max_{i,j} |a_{ij}|$ is sub-linear in $||H||_{\infty}$.

Recall that $\gamma_{r,s}$ denotes the generalised Hermite constant, introduced in Section 1.8.

Theorem 17. Let $Q \subset \mathbb{Z}^n$ and let H be an $1 \times n$ matrix such that $H \in \mathscr{R}(Q)$. For any integer m with 1 < m < n, there exists an $m \times n$ matrix A such that $A \in \mathscr{R}(Q)$ and

$$\|A\|_{\infty} \le c_2(m,n) \|H\|_{\infty}^{\frac{n-m}{n-1}},$$
(4.5)

where $c_2(m,n) = \gamma_{n-1,n-m}^{1/2} n^{(n-m)/(2(n-1))}$.

The proof of Theorem 17 makes use of results on rational subspaces obtained in [5]. Note that (4.5) improves the immediate bound $||A||_{\infty} \le ||H||_{\infty}$ when $||H||_{\infty} > c_2(m, n)^{(n-1)/(m-1)}$.



Figure 4.1: The lattice Λ does not intersect with the **0**-symmetric convex body *C* in any non-zero points. This means that the lattice Λ is admissible for *C*.

4.2 Admissible lattices of *m*-sparse *n*-cubes

By a *rational subspace* of \mathbb{R}^n we understand a subspace generated by integer vectors. A rational hyperplane can be written as $P = \{x \in \mathbb{R}^n : Hx = 0\}$, where $H = (H_{11}, \dots, H_{1n})$ is an $1 \times n$ integer matrix with $gcd(H) := gcd(H_{11}, \dots, H_{1n}) = 1$. We say that *P* has *height* $h(P) = ||H||_{\infty}$.

A lattice $\Lambda \subset \mathbb{R}^d$ is *(strictly) admissible* for a set $X \subset \mathbb{R}^d$ if $\Lambda \cap X = \{\mathbf{0}\}$. See, for example, Figure 4.1. For a comprehensive introduction to the theory of lattices we refer the reader to [13, 29].

Let *r* be a positive integer and *m* be a positive integer with 1 < m < n. We will consider an *m*-sparse *n*-dimensional cube

$$C_m^n(r) = \{ \mathbf{x} \in C^n(r) : \|\mathbf{x}\|_0 \le m \}.$$

An example of this *m*-sparse *n*-dimensional cube can be seen in Figure 4.2.

Constructing single measurements for unique recovery of sparse integer signals is closely linked to constructing admissible (n - 1)-dimensional lattices for $C_m^n(r)$. From the unique recovery perspective, it is desirable to find a

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Figure 4.2: The set of all 2-sparse signal S_2^3 inside a three dimensional cube $C^3(3) = \{x \in \mathbb{R}^3 : |x_i| \le 3 \text{ where } 1 \le i \le 3\}.$

rational hyperplane *P* of smallest possible height such that the lattice $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$. Similarly to (2.6), we consider the following optimisation problem. Given $m, n \in \mathbb{Z}$ with 1 < m < n and $r \in \mathbb{Z}_{>0}$, find

min{h(P) : P is a rational hyperplane in \mathbb{R}^n and the lattice $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$ }. (4.6)

The proofs of Theorems 15 and 16 will be based on the following estimates for the minimum in (4.6) that are of independent interest.

Theorem 18. For any $m, n \in \mathbb{Z}$ with 1 < m < n and $r \in \mathbb{Z}_{>0}$ there exists a rational hyperplane P in \mathbb{R}^n such that the lattice $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$ and

$$h(P) < p_{m,n}(r).$$
 (4.7)

To prove Theorem 18 we combine constructions from the proof of Theorem 6 with aggregation techniques outlined in Section 3.1. The next result shows that the polynomial $p_{m,n}(r)$ in (4.7) cannot be replaced by a polynomial in r of degree smaller than m-1.

Theorem 19. For any $m, n \in \mathbb{Z}$ with 1 < m < n, any $r \in \mathbb{Z}_{>0}$ and any rational hyperplane P in \mathbb{R}^n such that the lattice $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$ we have

$$h(P) > \frac{r^{m-1}}{\sqrt{m}} \tag{4.8}$$

holds.

Similarly to (4.4), for $m, n \in \mathbb{Z}$ with 1 < m < n, it would be interesting to estimate

$$c_3(m,n) = \limsup \inf \frac{h(P)}{r^{m-1}}, \qquad (4.9)$$

where the supremum limit is taken over all positive integers r and the infimum is taken over rational hyperplanes P in \mathbb{R}^n such that $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$. Theorems 18 and 19 imply the bounds

$$\frac{1}{\sqrt{m}} \leq c_3(m,n) \leq (m\Delta(m,n))^{m-1}.$$

4.3 Proofs of Theorem 18 and Theorem 15

We will begin with proving Theorem 18. The proof of Theorem 6 (Konyagin and Sudakov [33, Theorem 1.3]) gives two explicit constructions that can be used to obtain matrices $A \in \mathscr{R}(\mathbb{Z}_m^n)$ that satisfy conditions of Corollary 7. For completeness, we will outline these constructions here.

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Proof of Theorem 18

Observe first that $A \in \mathscr{R}(\mathbb{Z}_m^n)$ if and only if all $m \times m$ subdeterminants of A are nonzero. Therefore, if for some d satisfying m < n < d there exists an $m \times d$ matrix $A \in \mathscr{R}(\mathbb{Z}_m^d)$, then an $m \times n$ matrix in $\mathscr{R}(\mathbb{Z}_m^n)$ can be obtained by removing any d-n columns from A. Set first k = n-1. The first construction gives $A = (a_{ij}) \in \mathscr{R}(\mathbb{Z}_m^d)$ with $d \ge k+1$. The dimension d is chosen as an odd prime number satisfying $k+1 \le d \le 2k+1$. Subsequently, the entries of the matrix A are defined as $a_{ij} \equiv j^{i-1} \pmod{d}$ with $|a_{ij}| \le (d-1)/2 \le k$. In particular, for all j we have $a_{1j} = 1$. Next, set $k = (2n)^{(m-1)/m}$. The second construction gives $A = (a_{ij}) \in \mathscr{R}(\mathbb{Z}_m^d)$ with $d \ge k^{m/(m-1)}/2$. One can assume that $k^{m/(m-1)}/2 > k+1$ and, in particular, $k \ge 3$. The dimension d is now chosen as a prime number with $k^{m/(m-1)}/2 \le d \le k^{m/(m-1)}$. The entries of the matrix A satisfy $a_{ij} \equiv l_{ij}j^{i-1} \pmod{d}$, where l_{ij} are certain integers not divisible by d chosen in a such way that $|a_{ij}| \le k$. In particular, for all j one can take $l_{1j} = 1$, so that $a_{1j} = 1$.

In both constructions above, renumbering the rows of *A*, we may assume that $a_{mj} = 1$ for all *j*. Set $k = \Delta(m, n)$ and for s = m k r + 1 take

$$B=(1,s,\ldots,s^{m-1})$$

and

$$H = BA$$
.

We will show that the hyperplane P = ker(H) satisfies the conditions of Theorem 18. Let $F(\mathbf{x}) = A\mathbf{x}$ and let $F_i(\mathbf{x})$ denote the *i*th entry of $F(\mathbf{x})$, that is

$$F_1(\mathbf{x}) = a_{11}x_1 + \dots + a_{1n}x_n,$$

$$\vdots$$

$$F_m(\mathbf{x}) = a_{m1}x_1 + \dots + a_{mn}x_n.$$

For any $\mathbf{x} \in C_m^n(r)$ and any $i \in \{1, \ldots, m\}$, we have

$$F_i(\mathbf{x}) \le \|\mathbf{x}\|_0 \|A\|_{\infty} r \le m k r < s.$$

Lemma 9, applied with $D = C_m^n(r) \cap \mathbb{Z}^n$ and $q_i = s$ for i = 1, ..., m-1, implies that

$$\{\mathbf{x} \in C_m^n(r) \cap \mathbb{Z}^n : H\mathbf{x} = \mathbf{0}\} = \{\mathbf{x} \in C_m^n(r) \cap \mathbb{Z}^n : A\mathbf{x} = \mathbf{0}\}.$$
 (4.10)

Since $A \in \mathscr{R}(\mathbb{Z}_m^n)$ we have

$$\{\mathbf{x}\in C_m^n(r)\cap\mathbb{Z}^n: A\mathbf{x}=\mathbf{0}\}=\{\mathbf{0}\}.$$
(4.11)

Consequently, combining (4.10) and (4.11), the lattice $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$.

Finally, we obtain the bound

$$h(P) \le ||H||_{\infty} \le s^{m-1} + k \sum_{i=0}^{m-2} s^{i}$$

that implies (4.7).

Remark 1: For given sparsity level m, dimension n and cube size r the set $\{F(\mathbf{x}) : \mathbf{x} \in D\}$ constructed in the proof will likely allow a more accurate choice of parameters q_i in Lemma 9, resulting in an improvement on the bound (4.7). Further, aggregation techniques can be also applied to the matrices in $\Re(\mathbb{Z}_m^n)$ obtained using a probabilistic approach from the proof of Theorem 5 (Konyagin [32, Theorem 3]).

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Proof of Theorem 15

Let m = 2l and let P be a rational hyperplane in \mathbb{R}^n such that the lattice $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(2r) \cap \mathbb{Z}^n$. We can write $P = \ker(H)$ for an $1 \times n$ integer matrix H with $h(P) = ||H||_{\infty}$.

The inclusion

$$D(S_l^n(r)) \subset S_m^n(2r) = C_m^n(2r) \cap \mathbb{Z}^n$$

implies the condition (2.4) with $Q = D(S_l^n(r))$. Hence $H \in \mathcal{R}(Q)$. Finally, the bound (4.2) immediately follows from (4.7).

4.4 Examples of measurement matrices

In this section, we compute examples of $m \times n$ integer matrix A using the method outlined in the proof of Theorem 15. Specifically, we set the maximum absolute entry of A to k = n-1 and $k = \lceil (2n)^{(m-1)/m} \rceil$. Subsequently, we verify that all $m \times m$ square submatrices of A were nonsingular. Next, we fix a positive integer r and set s = m k r + 1. Then we construct a one-row matrix B which has powers of s as its entries, and multiply it by A to generate the measurement matrix H. Afterwards, we used a Python program to recover the signal \mathbf{x} by solving the equation $H\mathbf{x} = \mathbf{b}$, where \mathbf{b} is a measurement vector.

We tested two measurement matrices under different scenarios by modifying the sparsity level and the bound on entries of the signal x. In Examples 3 and 4, the Python program successfully recovered a unique signal x that satisfied Theorem 1. **Example 3.** Let n = 10 and m = 4. Using the method in the proof of Theorem 18, we construct a matrix A as follows:

$$A = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & -1 \\ 1 & 4 & 9 & 5 & 3 & 3 & 5 & 9 & 4 & 1 \\ 1 & 8 & 5 & 9 & 4 & 7 & 2 & 6 & 3 & -1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}.$$

The maximum absolute value is k = 9. Every $m \times m$ submatrix of A is nonsingular.

Next, we set r = 2, s = mkr + 1 = and compute 1×4 matrix B as :

$$B = \left(1 \quad s \quad s^2 \quad s^3\right) = \left(1 \quad 73 \quad 5329 \quad 389017\right).$$

We now create a sensing matrix H by multiplying A by B, which gives us:

$$H = \begin{pmatrix} 394420 & 431943 & 416322 & 437347 & 410557 & 426545 & 400047 \\ 421656 & 405305 & 383760 \end{pmatrix}.$$

Note that H satisfied (4.2).

Finally, we use the sensing matrix H to recover a signal

$$\boldsymbol{x}_0 = (-1, 0, 1, 0, 0, 0, 0, 0, 0, 0)^\top$$

by solving the system $H\mathbf{x} = \mathbf{b}$, where $\mathbf{b} = H\mathbf{x}_0 = 21902$. Notice that $||\mathbf{x}_0||_0 = 2$ and the absolute values of its entries are bounded by one. Using a Python program we checked that \mathbf{x}_0 is the unique solution over the signal space set $S_2^{10}(2)$. Thus, we have successfully recovered the sparse signal \mathbf{x} using the measurement matrix H and the set of linear equations $H\mathbf{x} = \mathbf{b}$. Interestingly, only replacing $S_2^{10}(2)$ with a substantially larger set $S_4^{10}(3)$ we obtain one extra solution to $H\mathbf{x} = 21902$:

$$\boldsymbol{x}_1 = (0, 0, 0, 3, -2, -3, 0, 0, 2, 0)^\top$$
.

In this situation $H \notin \mathscr{R}(D(S_4^{10}(3)))$.

Example 4. Next we construct an integer $m \times n$ matrix A with m = 4 and n = 12, with maximum absolute value $k = \lceil (2n)^{(m-1)/m} \rceil = 11$, such that all $m \times m$ submatrix are non-singular:

We set r = 2, s = mkr + 1 = and compute a 1×4 matrix $B = \begin{pmatrix} 1 & s & s^2 & s^3 \end{pmatrix}$, where s = 89. We then obtain the sensing matrix H by multiplying A by B:

$$H = \Big(712980 \ 768695 \ 713694 \ 697319 \ 768253 \ 769233 \ 745471$$

744493 713166 697859 744941 697136\Big).

One can verify that (4.2) is satisfied. We take

$$\boldsymbol{x}_0 = (0, 0, 0, -1, 0, 1, 0, 0, 0, 0, 0, 0)^\top$$

and solve the system $H\mathbf{x} = \mathbf{b}$, where $\mathbf{b} = H\mathbf{x}_0 = 71914$, to recover the sparse signal \mathbf{x} . A Python program was used to verify the uniqueness of the solution over $S_2^{12}(2)$.

Interestingly, replacing $S_2^{12}(2)$ by $S_4^{12}(4)$ gives extra solutions to $H\mathbf{x} = 71914$:

$$x_1 = (0, 3, 0, -2, -4, 0, 0, 3, 0, 0, 0, 0)^{\top},$$

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$$\mathbf{x}_2 = (0, -3, 0, 0, 4, 2, 0, -3, 0, 0, 0, 0)^\top$$

Clearly, by changing the signal space the unique recovery is no longer guaranteed. In this situation $H \notin \mathscr{R}(D(S_4^{12}(4)))$.

4.5 Proofs of Theorem 19 and Theorem 16

We will first prove Theorem 19.

Proof of Theorem 19

Let $A \in \mathbb{Z}^{m \times n}$, m < n, and let $\tau = \{i_1, \ldots, i_k\} \subseteq \{1, \ldots, n\}$ with $i_1 < i_2 < \cdots < i_k$. Recall that we denote by A_{τ} the $m \times k$ submatrix of A with columns indexed by τ . In the same manner, given $\mathbf{x} \in \mathbb{R}^n$, we denote by \mathbf{x}_{τ} the vector $(x_{i_1}, \ldots, x_{i_k})^{\top}$. The complement of τ in $\{1, \ldots, n\}$ is denoted as $\bar{\tau}$. For matrices A of rank m, the notation gcd(A) will be used for the greatest common divisor of all $m \times m$ subdeterminants of A.

The proof makes use of Theorem 11, which is a version of Siegel's Lemma obtained by Bombieri and Vaaler [9, Theorem 2].

Suppose, to derive a contradiction, that Theorem 19 does not hold. Then for some $m, n \in \mathbb{Z}_{>0}$ with 1 < m < n, and $r \in \mathbb{Z}_{>0}$ there exists a rational hyperplane P in \mathbb{R}^n such that $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$ and

$$h(P) \le \frac{r^{m-1}}{\sqrt{m}}$$
 (4.12)

There exists an $1 \times n$ integer matrix H such that $P = \ker(H)$ and $h(P) = ||H||_{\infty}$. Take $\tau = \{1, ..., m\}$. Observe that H cannot have zero entries, as otherwise its kernel *P* would contain the corresponding standard basis vectors. Hence, $H_{\tau} \neq \mathbf{0}$. By Theorem 11, applied with $M = H_{\tau}$, there exists an integer vector $\mathbf{x}_{\tau} \in \ker(H_{\tau})$ such that

$$0 < \|\boldsymbol{x}_{\tau}\|_{\infty}^{m-1} \le \frac{\|H_{\tau}\|_{2}}{\gcd(H_{\tau})} \le \sqrt{m} \|H_{\tau}\|_{\infty} \le \sqrt{m} h(P).$$
(4.13)

By the upper bound (4.12) we have

$$\|\boldsymbol{x}_{\tau}\|_{\infty} \leq r.$$

Consequently, the lifted vector

$$\begin{pmatrix} \mathbf{x}_{\tau} \\ \mathbf{0}_{\bar{\tau}} \end{pmatrix} \in C_m^n(r) \cap P \cap \mathbb{Z}^n,$$

contradicting the assumption that $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$. The obtained contradiction completes the proof of Theorem 19.

Remark 2: A minor improvement of (4.13) can be obtained using a refinement of Siegel's lemma proved in [3]. The latter refinement is also discussed for completeness in Section of this thesis; see the bound (3.8). Further, the last inequality in (4.13) can be slightly strengthened using the following observation. Since $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(1)$ and $m \ge 2$, we may assume that $H_{11} < H_{12} < \cdots < H_{1n}$. This allows choosing H_τ with $||H_\tau|| \le ||H||_{\infty} - n + m$.

Proof of Theorem 16

Take any $l, n \in \mathbb{Z}_{>0}$ with $1 \le l < n/2$, $r \in \mathbb{Z}_{>0}$ and any $1 \times n$ integer matrix $H \in \mathscr{R}(Q)$ with $Q = D(S_l^n(r))$. Consider the hyperplane $P = \ker(H)$. Set m = 2l and observe that

$$C_m^n(r)\cap\mathbb{Z}^n\subset Q$$

Therefore, the lattice $P \cap \mathbb{Z}^n$ is admissible for $C_m^n(r)$ and (4.8) implies (4.3).

4.6 Proof of Theorem 17

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The proof of Theorem 17 is based on the following result, which is a special case of Proposition 1 (ii) in [5].

Theorem 20. Let *S* be an one-dimensional rational subspace of \mathbb{R}^n . When 1 < m < n, there is a rational subspace $T \supset S$ of dimension *m* in \mathbb{R}^n with

$$\det(T \cap \mathbb{Z}^n) \le \gamma_{n-1,n-m}^{1/2} \det(S \cap \mathbb{Z}^n)^{(n-m)/(n-1)}.$$
(4.14)

The constant $\gamma_{n-1,n-m}^{1/2}$ here is best possible.

Take any $Q \subset \mathbb{Z}^n$ and suppose that we are given an integer $1 \times n$ matrix $H \in \mathscr{R}(Q)$. Let *m* be an integer with 1 < m < n and let *T* be the subspace from Theorem 20, applied to the rational subspace *S* of \mathbb{R}^n spanned by the row vector *H*.

By Theorem 11, applied with any $(n-m) \times n$ integer matrix M with T = ker(M), there exist m linearly independent integer vectors $g_1, \ldots, g_m \in T$ such that

$$\|\boldsymbol{g}_1\|_{\infty} \cdots \|\boldsymbol{g}_m\|_{\infty} \leq \frac{\sqrt{\det(MM^T)}}{\gcd(M)} = \det(T \cap \mathbb{Z}^n).$$
(4.15)

For a proof of the last equality in (4.15) we refer the reader to [45, Corollaries 5I-J]. Now we can form a matrix *A* with rows $\boldsymbol{g}_1^{\top}, \ldots, \boldsymbol{g}_m^{\top}$, so that $S \subset$ span_{\mathbb{R}}(A^{\top}). Observe that ker(A) \subset ker(H) and hence, $A \in \mathcal{R}(E)$.

Finally, combining (4.15), (4.14) and the bound $det(S \cap \mathbb{Z}^n) \leq \sqrt{n} ||H||_{\infty}$, we get the estimate (4.5):

$$\|A\|_{\infty} \leq \gamma_{n-1,n-m}^{1/2} n^{(n-m)/(2(n-1))} \|H\|_{\infty}^{(n-m)/(n-1)}.$$

Chapter 5

Probability of uniqueness of sparse solutions with integer bounded entries

In this chapter, we consider the uniqueness of sparse integer solutions of underdetermined linear systems from a probabilistic point of view. In the case of unbounded solutions, Konyagin [32] proved the following two theorems.

Theorem 21 (Konyagin [32]). Consider an $m \times m$ random matrix M with independent elements, each of which takes any integer value from -k to k with probability 1/(2k+1). Then the probability of the determinant of M being zero is at most $C^{m^2}k^{-m}(\log k + 1)^{m-1}$, where C > 0 is an absolute constant.

This result is linked to sparse solutions of a system $A\mathbf{x} = \mathbf{b}$ with $A \in \mathbb{Z}^{m \times n}$, m < n, since any *l*-sparse solution with $l \le m$ of the latter system corresponds to a singular $m \times m$ submatrix M of A. Theorem 21 was subsequently used to prove Theorem 5, which we recall here for convenience of the reader.

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Theorem 5 [Konyagin [32]]

For integers $k, m \ge 2$, and integer n with

$$m < n \le \frac{c^{-m}k^{m/(m-1)}}{\log(k)}$$

there exists an integer $m \times n$ matrix $A \in \mathscr{R}(\mathbb{Z}_m^n)$ such that $||A||_{\infty} = k$, where c is an absolute constant.

The assertions of Theorem 21 and 5 are nontrivial if *k* is sufficiently large and $m \le c \log k$, c > 0.

Our goal is to obtain a result similar to Theorem 21 which utilises the boundedness of the integer signals and apply it to estimate the probability of uniqueness of bounded integer sparse signals (see Theorem 25). A future direction of work is to use Theorem 25 to obtain an analog of Theorem 5 in the bounded setting.

The proofs of Theorems 24 and 25 below are based on the proof of Theorem 21. We will start with developing auxiliary tools required for dealing with bounded integer signals. Given $A \in \mathbb{Z}^{m \times n}$ and $b \in \mathbb{Z}^m$, we will denote by $\Gamma(A, b)$ the set of integer points in the affine subspace

$$H(A, \mathbf{b}) = \{\mathbf{x} \in \mathbb{R}^n : A\mathbf{x} = \mathbf{b}\},\$$

that is

$$\Gamma(A, \mathbf{b}) = H(A, \mathbf{b}) \cap \mathbb{Z}^n$$
.

The set $\Gamma(A, b)$ is either empty or an affine lattice of the form $\Gamma(A, b) = r + \Gamma(A)$, where r is any integer vector with Ar = b and $\Gamma(A) = \Gamma(A, 0)$ is the lattice formed by all integer points in the kernel ker(A) = { $x \in \mathbb{R}^n : Ax = 0$ } of the matrix A.

Lemma 22. Let $n > m \ge 2l$, and $A \in \mathbb{Z}^{m \times n}$. If for any $m \times 2l$ submatrix M of A

$$\Gamma(M) \cap D(S_1^n(r)) = \{\mathbf{0}\}$$
(5.1)

then any $\mathbf{x}_0 \in S_l^n(r)$ is the unique solution to the system $A\mathbf{x} = A\mathbf{x}_0$ over $S_l^n(r)$.

Proof. Take any $\mathbf{x}_0 \in S_l^n(r)$. Suppose, to derive a contradiction, that there exists a vector $\mathbf{x}_1 \in S_l^n(r)$ such that $\mathbf{x}_0 \neq \mathbf{x}_1$ and $A\mathbf{x}_0 = A\mathbf{x}_1$. Set $\mathbf{y} = \mathbf{x}_0 - \mathbf{x}_1 \neq \mathbf{0}$ and $\tau = \text{supp}(\mathbf{y})$. Take any set of indices $\sigma \subset [n]$ such that $|\sigma| = 2l$ and $\tau \subset \sigma$. Then

$$\mathbf{y}_{\sigma} \in \Gamma(A_{\sigma}) \cap D(S_{l}^{n}(r)),$$

contradicting (5.1).

In what follows, we will consider the case m = 2l < n. Given $U \in \mathbb{Z}^{m \times m}$ with rows $\boldsymbol{u}_1^T, \ldots, \boldsymbol{u}_m^T$ and rank $d \leq m-1$, we will denote by $\Sigma(U)$ the lattice formed by all integer points in the subspace $\operatorname{span}_{\mathbb{R}}(\boldsymbol{u}_1, \ldots, \boldsymbol{u}_m)$ spanned by $\boldsymbol{u}_1, \ldots, \boldsymbol{u}_m$. That is

$$\Sigma(U) = \operatorname{span}_{\mathbb{R}}(\boldsymbol{u}_1, \ldots, \boldsymbol{u}_m) \cap \mathbb{Z}^m.$$

Let $\lambda_i(\Sigma(U)) = \lambda_i(C^m(1), \Sigma(U))$ be the *i*th successive minimum of the cube $C^m(1)$ with respect to $\Sigma(U)$. Note that $1 \le \lambda_1(\Sigma(U)) \le \cdots \le \lambda_d(\Sigma(U))$.

Lemma 23. Suppose that rank(U) = m - 1 and

$$\lambda_1(\Sigma(U)) > R = (2r\sqrt{m})^{1/(m-1)}.$$

Then for $K = S_m^n(r)$ we have

$$\Gamma(U) \cap D(K) = \{\mathbf{0}\}.$$

Proof. Since $\lambda_1(\Sigma(U)) > R$, the lattice $\Sigma(U)$ is admissible for the cube section $S = C^m(R) \cap \operatorname{span}_{\mathbb{R}}(u_1, \dots, u_m)$. Therefore,

$$\det(\Sigma(U)) > \Delta(S), \tag{5.2}$$

where $\Delta(S)$ is the critical determinant of the set *S*. Further, by Theorem 5 in Section 17.3 of [29],

$$\Delta(S) \ge \frac{\operatorname{vol}(S)}{2^{m-1}}.$$
(5.3)

The cube-slicing inequality of Vaaler [51] implies that $vol(S) \ge (2R)^{m-1}$. Therefore, by (5.2) and (5.3),

$$\det(\Sigma(U)) > R^{m-1}.$$
(5.4)

Note that $\Sigma(U) = \Gamma(U)^{\perp} := \{ \mathbf{x} \in \mathbb{Z}^m : \mathbf{x}^\top \mathbf{y} = 0 \text{ for all } \mathbf{y} \in \Gamma(U) \}$. Hence, by [34], det($\Sigma(U)$) = det($\Gamma(U)$) and, using (5.4), we have

$$\det(\Gamma(U)) > R^{m-1}.$$
(5.5)

Therefore, the one-dimensional lattice $\Gamma(U)$ is generated by a basis vector \mathbf{v} with $\|\mathbf{v}\|_2 > R^{m-1} = 2r\sqrt{m}$. Since $D(K) \subset C^m(2r)$, we have $\Gamma(U) \cap D(K) = \{\mathbf{0}\}$.

Theorem 24. Fix $R \in \mathbb{R}_{>1}$, $k \in \mathbb{Z}_{>0}$ and consider a random matrix $M = (a_{ij}) \in \mathbb{Z}^{m \times m}$ with independent entries that take integer values from -k to k with probability $(2k + 1)^{-1}$. Then the probability that M is a singular matrix with $\lambda_1(\Sigma(M)) \leq R$ is at most

$$\frac{15^{m^2-m}R(\log(k)+1)^{m-2}}{k^m}.$$
(5.6)

Proof. Consider the set

$$\mathcal{M} = \{ M \in \mathbb{Z}^{m \times m} : \operatorname{rank}(M) \le m - 1, \|M\|_{\infty} \le k, \lambda_1(\Sigma(M)) \le R \}.$$

We will obtain an upper bound for $|\mathcal{M}|$ using the following strategy. Let \mathscr{P} be the set of primitive m-1 dimensional lattices $\Omega \subset \mathbb{Z}^m$ with $\lambda_1(C^m(1), \Omega) \leq R$ and $\lambda_{m-1}(C^m(1), \Omega) \leq k$. First, we will construct a function $\phi : \mathcal{M} \to \mathscr{P}$. Second, we split \mathscr{P} into subsets of lattices $\mathscr{P}_{\lambda_1,\dots,\lambda_{m-1}}$ that correspond to the fixed values of successive minima $\lambda_1, \dots, \lambda_{m-1}$ and write

$$\mathscr{P} = igcup_{1 \leq \lambda_1 \leq R} \mathscr{P}_{\lambda_1,...,\lambda_{m-1}} \, .$$

 $1 \leq \lambda_2, \dots, \lambda_{m-1} \leq k$

Then we use the estimate

$$|\mathcal{M}| \leq \sum_{\substack{1 \leq \lambda_1 \leq R \\ 1 \leq \lambda_2, \dots, \lambda_{m-1} \leq k}} |\mathcal{P}_{\lambda_1, \dots, \lambda_{m-1}}| \cdot L_{\lambda_1, \dots, \lambda_{m-1}}, \qquad (5.7)$$

where $L_{\lambda_1,\dots,\lambda_{m-1}} = \max_{\Omega \in \mathscr{P}_{\lambda_1,\dots,\lambda_{m-1}}} |\phi^{-1}(\Omega)|$. The desired bound is then obtained by estimating $|\mathscr{P}_{\lambda_1,\dots,\lambda_{m-1}}|$ and $L_{\lambda_1,\dots,\lambda_{m-1}}$.

Given $M \in \mathcal{M}$, we construct the lattice $\phi(M)$ as follows. If rank(M) = m-1, we set $\phi(M) = \Sigma(M)$. Suppose now that rank(M) < m-1 and take any basis $\mathbf{b}_1, \ldots, \mathbf{b}_l$ of the lattice $\Sigma(M)$. There exist the standard basis vectors $\mathbf{e}_{i_1}, \ldots, \mathbf{e}_{i_{m-1-l}}$ such that the linear subspace *S* spanned by $\mathbf{b}_1, \ldots, \mathbf{b}_l, \mathbf{e}_{i_1}, \ldots, \mathbf{e}_{i_{m-1-l}}$ has dimension m-1. We set $\phi(M) = S \cap \mathbb{Z}^m$. Note that $\phi(M) \in \mathcal{P}$ as in this case $\lambda_1(C^m(1), \phi(M)) = 1$ and $\lambda_{m-1}(C^m(1), \phi(M)) \leq k$.

Let us now fix the values $\lambda_1, \ldots, \lambda_{m-1}$ and estimate the size of the set $\mathscr{P}_{\lambda_1, \ldots, \lambda_{m-1}}$. Given a lattice $\Omega \in \mathscr{P}_{\lambda_1, \ldots, \lambda_{m-1}}$, there exist linearly independent vec-

tors $\boldsymbol{x}_1, \ldots, \boldsymbol{x}_{m-1} \in \Omega$ such that

$$\|\boldsymbol{x}_i\|_{\infty} = \lambda_i(C^m(1), \Omega).$$
(5.8)

The lattice Ω is the intersection of the m-1 dimensional subspace S of \mathbb{R}^m spanned by $\mathbf{x}_1, \ldots, \mathbf{x}_{m-1}$ and \mathbb{Z}^m . Consequently, Ω is uniquely determined by $\mathbf{x}_1, \ldots, \mathbf{x}_{m-1}$.

The number of vectors $\boldsymbol{x} \in \mathbb{Z}^m$ with $\|\boldsymbol{x}\|_{\infty} = \lambda$ is

$$(2\lambda+1)^m - (2\lambda-1)^m < 2m(2\lambda+1)^{m-1} < 3^m m \lambda^{m-1} < 5^m \lambda^{m-1}$$
.

Therefore,

$$|\mathscr{P}_{\lambda_1,\dots,\lambda_{m-1}}| \leq 5^{m^2-m} \prod_{i=1}^{m-1} \lambda_i^{m-1}.$$

Further,

$$\begin{aligned} |\mathscr{P}| &\leq 5^{m^2 - m} \sum_{\lambda_1 = 1}^R \sum_{\lambda_2 = 1}^k \cdots \sum_{\lambda_{m-1} = 1}^k \prod_{i=1}^{m-1} \lambda_i^{m-1} \\ &\leq 5^{m^2 - m} \sum_{\lambda_2 = 1}^k \cdots \sum_{\lambda_{m-1} = 1}^k \left(\sum_{\lambda_1 = 1}^R \prod_{i=1}^{m-1} \lambda_i^{m-1} \right) \\ &\leq 5^{m^2 - m} R^m \sum_{\lambda_2 = 1}^k \cdots \sum_{\lambda_{m-1} = 1}^k \prod_{i=2}^{m-1} \lambda_i^{m-1}. \end{aligned}$$

Next, for fixed values $\lambda_1, \ldots, \lambda_{m-1}$, we will estimate $L_{\lambda_1, \ldots, \lambda_{m-1}}$. For a lattice $\Omega \in \mathscr{P}_{\lambda_1, \ldots, \lambda_{m-1}}$ and, consequently, for vectors $\mathbf{x}_1, \ldots, \mathbf{x}_{m-1}$ satisfying (5.8), $|\phi^{-1}(\Omega)|$ is the number of matrices $M \in \mathscr{M}$ with $\Omega = \phi(M)$. Observe that the number of possible *j*th rows of such a matrix *M* is bounded from above by $|\Omega \cap C^m(k)|$.

Let $\mu_i = \lambda_i(C^m(k), \Omega)$, so that

$$\mu_i = \frac{\lambda_i}{k} \, .$$

$$|\Omega \cap C^m(k)| \le 2^{m-2} \prod_{i=1}^{m-1} \left(\frac{2}{\mu_i} + 1\right).$$

Hence

$$|\Omega \cap C^{m}(k)| \le 6^{m-1} \prod_{i=1}^{m-1} \frac{1}{\mu_{i}} = 6^{m-1} k^{m-1} \prod_{i=1}^{m-1} \lambda_{i}^{-1}.$$

Therefore, the number of matrices $M \in \mathcal{M}$ with $\Omega = \phi(M)$ is at most

$$|\Omega \cap C^m(k)|^m \le 6^{m^2 - m} k^{m^2 - m} \prod_{i=1}^{m-1} \lambda_i^{-m}.$$

Using (5.7), we get the bound

$$\begin{aligned} |\mathcal{M}| &\leq 30^{m^2 - m} k^{m^2 - m} \sum_{\lambda_1 = 1}^R \sum_{\lambda_2 = 1}^k \cdots \sum_{\lambda_{m-1} = 1}^k \prod_{i=1}^{m-1} \lambda_i^{-1} \\ &\leq 30^{m^2 - m} k^{m^2 - m} R \sum_{\lambda_2 = 1}^k \cdots \sum_{\lambda_{m-1} = 1}^k \prod_{i=2}^{m-1} \lambda_i^{-1} \,. \end{aligned}$$

Next,

$$\sum_{\lambda_2=1}^k \cdots \sum_{\lambda_{m-1}=1}^k \prod_{i=2}^{m-1} \lambda_i^{-1} = \left(\sum_{\lambda=1}^k \lambda^{-1}\right)^{m-2} \le (\log(k) + 1)^{m-2}$$

It is now sufficient to observe that the number of matrices $M \in \mathbb{Z}^{m \times m}$ with $|M| \le k$ is $(2k+1)^{m^2}$. Therefore, the required probability is

$$\leq 30^{m^2 - m} k^{m^2 - m} R(\log(k) + 1)^{m-2} (2k+1)^{-m^2}$$

$$\leq 15^{m^2 - m} k^{-m} R(\log(k) + 1)^{m-2}.$$

The bound (5.6) is smaller than the bound for the unbounded case in Theorem 21 provided that $R < \log(k) + 1$.

Theorem 25. Consider a random matrix $A = (a_{ij}) \in \mathbb{Z}^{m \times n}$, m = 2s < n with independent entries that take integer values from -k to k with probability $(2k + 1)^{m \times n}$

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1)⁻¹. Then the probability that any $\mathbf{x}_0 \in S_m^n(r)$ is the unique solution to the system $A\mathbf{x} = A\mathbf{x}_0$ over $S_m^n(r)$ is at least

$$1 - \frac{\binom{n}{m} \cdot 15^{m^2 - m} R(\log(k) + 1)^{m-2}}{k^m},$$
(5.9)

where $R = (2r\sqrt{m})^{1/(m-1)}$.

Proof. By Theorem 24 and the union bound, the probability that there exists a singular submatrix M of A with $\lambda_1(\Omega(M)) \le R = (2r\sqrt{m})^{1/(m-1)}$ is at most

$$\binom{n}{m} \cdot 15^{m^2-m} k^{-m} R(\log(k)+1)^{m-2},$$

Therefore, by Lemma 23, with probability at least (6.1) the matrix A satisfies condition (5.1). The result now follows by Lemma 22.

Chapter 6

Conclusions

In this final chapter, we give a brief overview of the contributions obtained in this thesis and, additionally, outline possible directions for future research.

6.1 Conclusions

The main results of the thesis are included in Chapter 4. Firstly, for integer bounded *l*-sparse *n*-dimensional signals, we obtained general estimates for the minimum size of a measurement matrix that corresponds to the minimum in (2.6). A straightforward approach, based on constructing admissible lattices for convex hulls of the sparse sets, would only show existence of a measurement $1 \times n$ matrix *H* with

$$\|H\|_{\infty}=O(r^{n-1})$$

where *r* is the entry size and the implicit constant depends on the sparsity level *l* and dimension *n*. Here 2l < n. We showed that the exponent n - 1 in (4.1) can be replaced with 2l - 1. Specifically, for a polynomial $p_{2l,n}(r)$ of

degree 2l - 1, Theorem 15 gives the bound

$$||H||_{\infty} < p_{2l,n}(2r).$$

Secondly, we give a lower bound for the minimum in (2.6). Roughly speaking, for any measurement matrix *H* in our setting, Theorem 16 shows that

$$\|H\|_{\infty} > \frac{r^{2l-1}}{\sqrt{2l}}$$

This implies that the polynomial $p_{2l,n}(2r)$ above cannot be replaced by a polynomial in r with degree smaller than 2l - 1.

Although for finite signal spaces a single measurement is sufficient for unique recovery, it is interesting to know whether allowing extra measurements would result in reducing measurements' entries. We indeed confirmed that such a reduction possible.

Avoiding technical details, given any $1 \times n$ measurement matrix H, Theorem 17 states existence of $m \times n$ measurement matrices A with entries sublinear in $||H||_{\infty}$:

$$||A||_{\infty} \leq c_2(m,n) ||H||_{\infty}^{\frac{n-m}{n-1}},$$

where $c_2(m, n)$ is a constant.

Another group of our results considers admissible lattices of *m*-sparse *n*cubes $C_m^n(r)$. In fact, constructing single measurements for unique recovery of sparse integer signals is directly linked to constructing admissible (n -1)-dimensional lattices for $C_m^n(r)$. Such lattices can be described in terms of rational hyperplanes. Then the height of a hyperplane *P*, denoted in the thesis by h(P) corresponds to the maximum absolute entry of a measurement matrix associated with *P*.

6.1. CONCLUSIONS

Theorem 18 shows existence of an admissible lattice determined by a rational subspace *P* of a relatively small height. Specifically,

$$h(P) < p_{m,n}(r).$$

As its counterpart, Theorem 19 shows that the polynomial $p_{m,n}(r)$ in the latter bound cannot be replaced by a polynomial in r of degree smaller than m-1:

$$h(P) > \frac{r^{m-1}}{\sqrt{m}}$$

holds.

Additionally, in Chapter 5, we consider the uniqueness of sparse integer solutions of underdetermined linear systems from a probabilistic perspective. Our main contribution here gives a result, similar to Theorem 21, which utilises the boundedness of the integer signals. Avoiding technical details, Theorem 25 gives a lower bound for the probability of obtaining a measurement matrix for integer signals from *m*-sparse hypercubes. This probability is at least

$$1 - \frac{\binom{m}{n} \cdot 15^{m^2 - m} R(\log(k) + 1)^{m-2}}{k^m},$$

where *r* is the cube size, *k* bounds measurement matrix absolute entries and $R = (2r\sqrt{m})^{1/(m-1)}$. Hence, as *k* grows, the probability of unique recovery rapidly tends to one.
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