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A SHARP UPPER BOUND FOR THE LATTICE PROGRAMMING GAP

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ABSTRACT. Given a full-dimensional lattice $\Lambda \subset \mathbb{Z}^d$ and a vector $\mathbf{l} \in \mathbb{Q}_{>0}^d$, we consider the family of the lattice problems

$$(0.1) \quad \text{Minimize } \{\mathbf{l} \cdot \mathbf{x} : \mathbf{x} \equiv \mathbf{r} \pmod{\Lambda}, \mathbf{x} \in \mathbb{Z}_{\geq 0}^d\}, \quad \mathbf{r} \in \mathbb{Z}^d.$$

The *lattice programming gap* $\text{gap}(\Lambda, \mathbf{l})$ is the largest value of the minima in (0.1) as \mathbf{r} varies over \mathbb{Z}^d . We obtain a sharp upper bound for $\text{gap}(\Lambda, \mathbf{l})$.

1. INTRODUCTION AND STATEMENT OF RESULTS

For linearly independent $\mathbf{b}_1, \dots, \mathbf{b}_k$ in \mathbb{R}^d , the set $\Lambda = \{\sum_{i=1}^k x_i \mathbf{b}_i, x_i \in \mathbb{Z}\}$ is a k -dimensional *lattice* with *basis* $\mathbf{b}_1, \dots, \mathbf{b}_k$ and *determinant* $\det(\Lambda) = (\det[\mathbf{b}_i \cdot \mathbf{b}_j]_{1 \leq i, j \leq k})^{1/2}$, where $\mathbf{b}_i \cdot \mathbf{b}_j$ is the standard inner product of the basis vectors \mathbf{b}_i and \mathbf{b}_j . The points $\mathbf{x}, \mathbf{y} \in \mathbb{R}^d$ are *equivalent modulo* Λ , denoted as $\mathbf{x} \equiv \mathbf{y} \pmod{\Lambda}$, if the difference $\mathbf{x} - \mathbf{y}$ is a point of Λ .

For a positive rational vector $\mathbf{l} \in \mathbb{Q}_{>0}^d$, a d -dimensional integer lattice $\Lambda \subset \mathbb{Z}^d$ and an integer vector $\mathbf{r} \in \mathbb{Z}^d$ we consider the lattice problem

$$(1.1) \quad \text{Minimize } \{\mathbf{l} \cdot \mathbf{x} : \mathbf{x} \equiv \mathbf{r} \pmod{\Lambda}, \mathbf{x} \in \mathbb{Z}_{\geq 0}^d\}.$$

Let $m(\Lambda, \mathbf{l}, \mathbf{r})$ denote the value of the minimum in (1.1). We are interested in the *lattice programming gap* $\text{gap}(\Lambda, \mathbf{l})$ of (1.1) defined as

$$(1.2) \quad \text{gap}(\Lambda, \mathbf{l}) = \max_{\mathbf{r} \in \mathbb{Z}^d} m(\Lambda, \mathbf{l}, \mathbf{r}).$$

The lattice programming gaps were introduced and studied for sublattices of all dimensions in \mathbb{Z}^d by Hosten and Sturmfels [14]. Computing $\text{gap}(\Lambda, \mathbf{l})$ is known to be NP-hard when d is a part of input (see [1]). For fixed d the value of $\text{gap}(\Lambda, \mathbf{l})$ can be computed in polynomial time (see Section 3 in [14], [10] and [9]).

The lower and upper bounds for $\text{gap}(\Lambda, \mathbf{l})$ in terms of the parameters Λ, \mathbf{l} were given in [1]. The lower bound is known to be sharp. In this paper we improve on the upper bound and show that the obtained bound is attained for parameters Λ, \mathbf{l} that satisfy certain arithmetic properties.

Let $|\cdot|$ denote the Euclidean norm and let γ_d be the d -dimensional Hermite constant (see e.g. Section IX.7 in [7]). In [1] it was shown that for any $\mathbf{l} \in \mathbb{Q}_{>0}^d$, $d \geq 2$, and any d -dimensional lattice $\Lambda \subset \mathbb{Z}^d$

$$(1.3) \quad \text{gap}(\Lambda, \mathbf{l}) \leq \frac{d\gamma_d^{d/2} \det(\Lambda)(\sum_{i=1}^d l_i + |\mathbf{l}|)}{2} - \sum_{i=1}^d l_i.$$

The bound (1.3) was obtained using a geometric argument based on estimating the covering radius of a simplex, associated with the vector \mathbf{l} , via the covering radius of the unit d -dimensional ball. Note that by a result of Blichfeldt (see e.g. §38 in Chapter 6 of [13]) $\gamma_d \leq 2 \left(\frac{d+2}{\sigma_d} \right)^{2/d}$, where σ_d is the volume of the unit d -ball; thus $\gamma_d = O(d)$. It follows from results in [2, Section 6] that the order $\text{gap}(\Lambda, \mathbf{l}) = O_{d,\mathbf{l}}(\det(\Lambda))$, where the constant depends on d and \mathbf{l} , cannot be improved.

Let $\|\cdot\|_\infty$ denote the maximum norm. In this paper we use coverings that are based on the arithmetic properties of the integer lattices and improve the bound (1.3) as follows.

Theorem 1.1. *For any $\mathbf{l} \in \mathbb{Q}_{>0}^d$, $d \geq 2$, and any d -dimensional lattice $\Lambda \subset \mathbb{Z}^d$*

$$(1.4) \quad \text{gap}(\Lambda, \mathbf{l}) \leq (\det(\Lambda) - 1) \|\mathbf{l}\|_\infty.$$

Using a link between the lattice programming gaps and the Frobenius numbers we also show that the bound (1.4) is sharp.

Theorem 1.2. *For $d \geq 2$ and any positive integer D there exist $\mathbf{l} \in \mathbb{Z}_{>0}^d$ and a lattice $\Lambda \subset \mathbb{Z}^d$ of determinant $\det(\Lambda) = D$ such that*

$$(1.5) \quad \text{gap}(\Lambda, \mathbf{l}) = (D - 1) \|\mathbf{l}\|_\infty.$$

2. COVERINGS OF \mathbb{R}^d AND LATTICE PROGRAMMING GAPS

Recall that the *Minkowski sum* $X + Y$ of the sets $X, Y \subset \mathbb{R}^d$ consists of all points $\mathbf{x} + \mathbf{y}$ with $\mathbf{x} \in X$ and $\mathbf{y} \in Y$. For a set $K \subset \mathbb{R}^d$ and a lattice $\Lambda \subset \mathbb{R}^d$, the Minkowski sum $K + \Lambda$ is a *packing* if the translates of K are mutually disjoint, a *covering* if $\mathbb{R}^d = K + \Lambda$ and a *tiling* if it is both packing and covering, simultaneously.

Let Λ be a lattice in \mathbb{R}^d with basis $\mathbf{b}_1, \dots, \mathbf{b}_d$. Let Λ_i denote the lattice generated by the first i basis vectors $\mathbf{b}_1, \dots, \mathbf{b}_i$ and let $\pi_i : \mathbb{R}^d \rightarrow \text{span}_{\mathbb{R}}(\Lambda_{i-1})^\perp$ be the orthogonal projection onto the subspace $\text{span}_{\mathbb{R}}(\Lambda_{i-1})^\perp$ orthogonal to $\mathbf{b}_1, \dots, \mathbf{b}_{i-1}$.

The vectors $\hat{\mathbf{b}}_i = \pi_i(\mathbf{b}_i)$ can be obtained using the Gram-Schmidt orthogonalisation of $\mathbf{b}_1, \dots, \mathbf{b}_d$:

$$\begin{aligned} \hat{\mathbf{b}}_1 &= \mathbf{b}_1, \\ \hat{\mathbf{b}}_i &= \mathbf{b}_i - \sum_{j=1}^{i-1} \mu_{i,j} \hat{\mathbf{b}}_j, \quad j = 2, \dots, d, \end{aligned}$$

where $\mu_{i,j} = (\mathbf{b}_i \cdot \hat{\mathbf{b}}_j) / |\hat{\mathbf{b}}_j|^2$.

Define the box $B = B(\mathbf{b}_1, \dots, \mathbf{b}_d)$ as

$$B = [0, \hat{\mathbf{b}}_1) \times \cdots \times [0, \hat{\mathbf{b}}_d).$$

We will need the following well-known and useful observation.

Lemma 2.1. *$B + \Lambda$ is a tiling of \mathbb{R}^d .*

Tilings of \mathbb{R}^d with lattice translates of B were implicitly used already in the classical Babai's nearest lattice point algorithm (see [3] and Theorem 5.3.26 in [11]) and in the work of Lagarias, Lenstra and Schorr on Korkin-Zolotarev bases (see the proof of Theorem 2.6 in [16]). Lemma 2.1 was also explicitly stated (with translated B) by Cai and Nerurkar (see [6], Lemma 2). A proof of this result can be obtained by modifying the proof of Theorem 5.3.26 in [11]. We also remark that for the purposes of this paper we only need the coverings of \mathbb{R}^d by the lattice translates of the closure of B .

In what follows, \mathcal{K}^d will denote the space of all d -dimensional *convex bodies*, i.e., closed bounded convex sets with non-empty interior in the d -dimensional Euclidean space \mathbb{R}^d . Let also \mathcal{L}^d denote the set of all d -dimensional lattices in \mathbb{R}^d . For $K \in \mathcal{K}^d$ and $\Lambda \in \mathcal{L}^d$ the *covering radius* of K with respect to Λ is the smallest positive number ρ such that any point $\mathbf{x} \in \mathbb{R}^d$ is covered by $\rho K + \Lambda$, that is

$$\rho(K, \Lambda) = \min\{\rho > 0 : \mathbb{R}^d = \rho K + \Lambda\}.$$

For further information on covering radii in the context of the geometry of numbers see e.g. Gruber [12] and Gruber and Lekkerkerker [13].

Given $\mathbf{l} \in \mathbb{Q}_{>0}^d$, consider the simplex $\Delta_{\mathbf{l}} = \{\mathbf{x} \in \mathbb{R}_{\geq 0}^d : \mathbf{l} \cdot \mathbf{x} \leq 1\}$. As it was shown in [1], the lattice programming gap can be expressed via the covering radius of $\Delta_{\mathbf{l}}$ with respect to Λ :

$$(2.1) \quad \text{gap}(\Lambda, \mathbf{l}) = \rho(\Delta_{\mathbf{l}}, \Lambda) - \sum_{i=1}^d l_i.$$

3. PROOF OF THEOREM 1.1

We will obtain an upper bound for $\text{gap}(\Lambda, \mathbf{l})$ in terms of \mathbf{l} and certain parameters of the lattice Λ that will imply (1.4).

By Theorem I (A) and Corollary 1 in Chapter I of Cassels [7], there exists a basis $\mathbf{b}_1, \dots, \mathbf{b}_d$ of the lattice Λ of the form

$$(3.1) \quad \begin{aligned} \mathbf{b}_1 &= v_{11} \mathbf{e}_1, \\ \mathbf{b}_2 &= v_{21} \mathbf{e}_1 + v_{22} \mathbf{e}_2, \\ &\vdots \\ \mathbf{b}_d &= v_{d1} \mathbf{e}_1 + \cdots + v_{dd} \mathbf{e}_d, \end{aligned}$$

where \mathbf{e}_i are the standard basis vectors of \mathbb{Z}^d , the coefficients v_{ij} are integers, $v_{ii} > 0$ and $0 \leq v_{ij} < v_{jj}$.

Lemma 3.1. *We have*

$$(3.2) \quad \text{gap}(\Lambda, \mathbf{l}) \leq l_1 v_{11} + \cdots + l_d v_{dd} - \sum_{i=1}^d l_i.$$

Proof. Note that the Gram-Schmidt orthogonalisation of $\mathbf{b}_1, \dots, \mathbf{b}_d$ has the form

$$(3.3) \quad \hat{\mathbf{b}}_1 = v_{11} \mathbf{e}_1, \hat{\mathbf{b}}_2 = v_{22} \mathbf{e}_2, \dots, \hat{\mathbf{b}}_d = v_{dd} \mathbf{e}_d.$$

Hence, the box $B = B(\mathbf{b}_1, \dots, \mathbf{b}_d)$ can be written as

$$B = [0, v_{11}) \times \cdots \times [0, v_{dd}).$$

By Lemma 2.1, $B + \Lambda$ is a tiling of \mathbb{R}^d . In particular, $B + \Lambda$ covers \mathbb{R}^d .

Since $B \subset (l_1 v_{11} + \cdots + l_d v_{dd}) \Delta_{\mathbf{l}}$, we have

$$\rho(\Delta_{\mathbf{l}}, \Lambda) \leq l_1 v_{11} + \cdots + l_d v_{dd}.$$

By (2.1), the bound (3.2) holds. □

Consider the simplex $\Delta = \text{conv} \{\mathbf{1}, \mathbf{p}_1, \dots, \mathbf{p}_d\}$, where $\text{conv} \{\cdot\}$ denotes the convex hull, $\mathbf{1}$ is the all-one vector and

$$\begin{aligned} \mathbf{p}_1 &= (\det(\Lambda), 1, \dots, 1)^t, \\ \mathbf{p}_2 &= (1, \det(\Lambda), \dots, 1)^t, \\ &\vdots \\ \mathbf{p}_d &= (1, 1, \dots, \det(\Lambda))^t. \end{aligned}$$

It is easy to see that

$$(3.4) \quad \{\mathbf{x} \in \mathbb{R}_{\geq 1}^d : x_1 \cdots x_d = \det(\Lambda)\} \subset \Delta.$$

Since Δ is a convex bounded polyhedron, the maximum of the linear function $\mathbf{l} \cdot \mathbf{x}$ over Δ is attained at one of its vertices $\mathbf{1}, \mathbf{p}_1, \dots, \mathbf{p}_d$. Therefore

$$(3.5) \quad \max\{\mathbf{l} \cdot \mathbf{x} : \mathbf{x} \in \Delta\} = (\det(\Lambda) - 1) \|\mathbf{l}\|_{\infty} + \sum_{i=1}^d l_i.$$

Since $v_{11} \cdots v_{dd} = \det(\Lambda)$, we obtain by (3.4) and (3.5)

$$(3.6) \quad l_1 v_{11} + \cdots + l_d v_{dd} \leq (\det(\Lambda) - 1) \|\mathbf{l}\|_{\infty} + \sum_{i=1}^d l_i.$$

By (3.2) and (3.6) we obtain (1.4).

4. PROOF OF THEOREM 1.2

In this section we will use classical results of Brauer [4] and Brauer and Seelbinder [5] to prove Theorem 1.2. In the course of the proof we also show that the bound (3.2) in Lemma 3.1 is sharp.

Let $\mathbf{a} = (a_1, \dots, a_{d+1})^t \in \mathbb{Z}_{>0}^{d+1}$ be a positive integer vector with coprime entries, that is $\gcd(a_1, \dots, a_{d+1}) = 1$. Consider the lattice $\Lambda = \Lambda(\mathbf{a})$ defined as

$$\Lambda = \{\mathbf{x} \in \mathbb{Z}^d : a_2x_1 + \dots + a_{d+1}x_d \equiv 0 \pmod{a_1}\}.$$

Note that $\det(\Lambda) = a_1$ (see e.g. Corollary 3.2.20 in [8]).

Let

$$f_1 = a_1, f_2 = \gcd(a_1, a_2), \dots, f_{d+1} = \gcd(a_1, a_2, \dots, a_{d+1}) = 1.$$

Consider the basis $\mathbf{b}_1, \dots, \mathbf{b}_d$ of the lattice Λ given by (3.1). The next lemma shows that the Gram-Schmidt box $B(\mathbf{b}_1, \dots, \mathbf{b}_d)$ is entirely determined by the parameters f_i .

Lemma 4.1. *The box $B = B(\mathbf{b}_1, \dots, \mathbf{b}_d)$ has the form*

$$B = \left[0, \frac{f_1}{f_2}\right) \times \left[0, \frac{f_2}{f_3}\right) \times \dots \times \left[0, \frac{f_d}{f_{d+1}}\right).$$

Proof. By the definition of the box B and (3.3), it is enough to show that

$$(4.1) \quad v_{11} = \frac{f_1}{f_2}, v_{22} = \frac{f_2}{f_3}, \dots, v_{dd} = \frac{f_d}{f_{d+1}}.$$

Recall that Λ_i denotes the sublattice of Λ generated by the first i basis vectors $\mathbf{b}_1, \dots, \mathbf{b}_i$. We can write Λ_i in the form

$$\Lambda_i = \left\{ (x_1, \dots, x_i, 0, \dots, 0) \in \mathbb{Z}^d : \frac{a_2}{f_{i+1}}x_1 + \dots + \frac{a_{i+1}}{f_{i+1}}x_i \equiv 0 \pmod{\frac{a_1}{f_{i+1}}} \right\}.$$

Hence, $\det(\Lambda_i) = a_1/f_{i+1}$. On the other hand, (3.1) implies that $\det(\Lambda_i) = v_{11}v_{22} \dots v_{ii}$. Since $\det(\Lambda) = v_{11}v_{22} \dots v_{dd} = a_1$, we have $f_{i+1} = v_{i+1}v_{i+1} \dots v_{dd}$ for $i \leq d-1$, which immediately implies (4.1). □

The *Frobenius number* $F(\mathbf{a})$ associated with the integer vector \mathbf{a} is the largest integer number which cannot be represented as a nonnegative integer combination of the a_i 's. The problem of finding $F(\mathbf{a})$ has a long history and is traditionally referred to as the *Frobenius problem*, see e. g. [18].

Set $\mathbf{l}(\mathbf{a}) = (a_2, \dots, a_{d+1})^t$. It is known (see e.g. proof of Theorem 1.1 in [1] and Section 5.1 in [17]) that

$$(4.2) \quad \text{gap}(\Lambda(\mathbf{a}), \mathbf{l}(\mathbf{a})) = F(\mathbf{a}) + a_1.$$

Note also that, in this special case, (2.1) follows from Theorem 2.5 of Kannan [15].

By Lemma 4.1, the bound (3.2) for $\text{gap}(\Lambda(\mathbf{a}), \mathbf{l}(\mathbf{a}))$ given by Lemma 3.1 can be obtained by replacing $F(\mathbf{a})$ on the right hand side of (4.2) by the estimate

$$(4.3) \quad F(\mathbf{a}) \leq C(\mathbf{a}) := a_2 \frac{f_1}{f_2} + \cdots + a_{d+1} \frac{f_d}{f_{d+1}} - \sum_{i=1}^{d+1} a_i$$

given in Brauer [4]. It should be remarked here that Brauer [4] rather worked with the quantity $F^+(\mathbf{a}) = F(\mathbf{a}) + \sum_{i=1}^{d+1} a_i$, the largest number which cannot be represented as a *positive* integer combination of the a_i 's. Brauer [4] and, subsequently, Brauer and Seelbinder [5] proved that the bound (4.3) is sharp and obtained the following necessary and sufficient condition for the equality $F(\mathbf{a}) = C(\mathbf{a})$.

Lemma 4.2 (see Theorem 5 in [4] and Theorem 1 in [5]). *Let $\mathbf{a} = (a_1, \dots, a_{d+1})^t \in \mathbb{Z}_{>0}^{d+1}$, $d \geq 2$, with $\gcd(a_1, \dots, a_{d+1}) = 1$. Then $F(\mathbf{a}) = C(\mathbf{a})$ if and only if for $m = 3, 4, \dots, d+1$ the integer a_m/f_m is representable in the form*

$$(4.4) \quad \frac{a_m}{f_m} = \sum_{i=1}^{m-1} \frac{a_i}{f_{m-1}} y_{mi}$$

with integers $y_{mi} \geq 0$.

For $s = 2, 3, \dots, d+1$, let

$$\mathbf{a}^{(s)} = \left(\frac{a_1}{f_s}, \dots, \frac{a_s}{f_s} \right)^t.$$

The condition (4.4) is satisfied, in particular, if

$$\frac{a_m}{f_m} > F(\mathbf{a}^{(m-1)}).$$

Hence the bound (3.2) in Lemma 3.1 is sharp and the vectors \mathbf{a} satisfying (4.4) can be easily constructed. To show that (1.4) is sharp, we will use a special case of Lemma 4.2, that regards the optimality of the Schur's upper bound for the Frobenius number (see [4]). Suppose that a vector $\mathbf{a} \in \mathbb{Z}_{>0}^{d+1}$ with coprime entries satisfies the following conditions:

$$(4.5) \quad \begin{aligned} (i) & \quad D = a_1 \leq a_2 \leq \cdots \leq a_{d+1}, \\ (ii) & \quad a_2 \equiv a_3 \equiv \cdots \equiv a_r \pmod{a_1} \text{ for some index } r \geq 3, \\ (iii) & \quad a_{r+1} = a_{r+2} = \cdots = a_{d+1}. \end{aligned}$$

By Theorem 3 in [4] (cf. Theorem 4 *ibid.*) conditions (4.5) imply that $F(\mathbf{a}) = a_1 a_{d+1} - a_1 - a_{d+1}$. Hence $\text{gap}(\Lambda(\mathbf{a}), \mathbf{l}(\mathbf{a})) = (a_1 - 1)a_{d+1} = (D - 1)\|\mathbf{l}\|_\infty$. The theorem is proved.

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