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Explaining Arguments’ Strength: Unveiling the Role of Attacks and Supports

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Abstract

Quantitatively explaining the strength of arguments under gradual semantics has recently received increasing attention. Specifically, several works in the literature provide quantitative explanations by computing the attribution scores of arguments. These works disregard the importance of attacks and supports, even though they play an essential role when explaining arguments’ strength. In this paper, we propose a novel theory of *Relation Attribution Explanations (RAEs)*, adapting Shapley values from game theory to offer fine-grained insights into the role of attacks and supports in quantitative bipolar argumentation towards obtaining the arguments’ strength. We show that RAEs satisfy several desirable properties. We also propose a probabilistic algorithm to approximate RAEs efficiently. Finally, we show the application value of RAEs in fraud detection and large language models case studies.

1 Introduction

Explainable Artificial Intelligence (XAI) has received increasing attention in fields such as finance and healthcare, which demand a reliable and legitimate reasoning process. Argumentation Frameworks (AFs), e.g. as first studied in [Dung, 1995], are promising tools in the XAI field [Mittelstadt *et al.*, 2019] due to their transparency and interpretability, as well as their ability to support reasoning about conflicting information [Čyras *et al.*, 2021; Albini *et al.*, 2020; Potyka, 2021; Potyka *et al.*, 2023; Ayoobi *et al.*, 2023]. In Quantitative Bipolar AFs (QBAFs) [Baroni *et al.*, 2015], each argument has a *base score*, and its final *strength* is computed by *gradual semantics* based on the strength of its attackers and supporters [Baroni *et al.*, 2019]. QBAFs can be deployed to support several applications. For example, [Cocarascu *et al.*, 2019] build QBAFs to rate movies by aggregating movie reviews. The QBAFs have a hierarchical structure, where the goodness of movies is at the top and influenced by arguments about criteria like the quality of acting and directing. These criteria/arguments, in turn, can be affected by subcriteria/subarguments like the performance of particular actors. In this application, the base scores of arguments are obtained

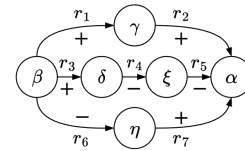


Figure 1: Graphical view of (elements of) a QBAF resulting from aggregating movie reviews (here, nodes are arguments, edges labelled + are supports, edges labelled - are attacks, and the r_i are identifiers for the edges (for ease of reference)).

from reviews via a natural language processing pipeline; finally, a gradual semantics is applied to determine the final strength of movies as their rating scores.

While the gradual semantics of a QBAF provides an assessment of arguments (e.g., when using QBAFs for aggregating movie reviews, the rating scores of movies), we may also be interested in an intuitive understanding of the underlying reasoning process. This leads to an interesting research question initially raised by [Delobelle and Villata, 2019]: **given an argument of interest (*topic argument*) in a QBAF, how to explain the reasoning outcome (i.e., the strength) of this topic argument?**

Most current approaches in the literature address this question by defining *argument-based attribution explanations* [Delobelle and Villata, 2019; Čyras *et al.*, 2022; Yin *et al.*, 2023], which explain the strength of the topic argument by assigning *attribution scores* to arguments: the greater the attribution score, the greater the argument’s contribution to the topic argument. However, in many cases, more fine-grained *relation-based attribution explanations* (RAEs) may be beneficial, or even necessary. For illustration, consider Figure 1, and assume that the QBAF (partially) depicted therein results from aggregating movie reviews as in [Cocarascu *et al.*, 2019], where α is a movie to be rated (topic argument)¹. Here, the *review* β has a positive argument attribution score by supporting the *famous actor* γ and the *influential director* δ , which attacks *bad directing* ξ , but this argument view conceals the fact that β also weakens α by attacking its *genre* η , which supports the topic argument. In contrast, (our) RAEs give more fine-grained insights: although β has a positive contribution via r_1

¹We give concrete values for the RAEs in Figure 1 in arxiv.org/abs/2404.14304.

and r_3 to α , it also has a negative contribution via r_6 .

Motivated by the aforementioned considerations, we make the following contributions:

- We propose a novel theory of RAEs (Section 4).
- We study desirable properties of RAEs under several gradual semantics (Section 5).
- We propose a probabilistic algorithm to efficiently approximate RAEs (Section 6).
- We carry out two case studies to demonstrate the practical usefulness of RAEs (Section 7).

The proofs of all results are in arxiv.org/abs/2404.14304.

2 Related Work

[Čyras *et al.*, 2022] propose the general idea of *contribution functions* that compute quantitative *contributions* from one argument to another under a given gradual semantics for QBAFs and study three such functions, described below.

The **removal-based** contribution function proposed by [Delobelle and Villata, 2019] measures how the strength of the topic argument changes if an argument is removed. In general, removal-based explanations are simple and intuitive for users to understand without a high cognitive burden. However, a problem with them is that removing an argument will also remove paths from its predecessor to the topic argument. The measure can therefore overestimate the contribution of an argument. To solve this problem, [Delobelle and Villata, 2019] propose to cut off the direct relations to an argument before removing it, to obtain the mere contribution of this argument.

The **gradient-based** contribution function captures the *sensitivity* of the topic argument w.r.t. another argument. It is based on the partial derivative of the topic argument’s strength w.r.t. the base score of another argument. Arguments with high sensitivity are seen as important. Following this idea, [Yin *et al.*, 2023] further explored the gradient-based contribution function under the *DF-QuAD* gradual semantics [Rago *et al.*, 2016] and studied its properties in this setting.

The **Shapley-based** contribution function uses the Shapley value from coalitional game theory [Shapley, 1951] to assign contribution scores. Each argument in a QBAF is seen as a *player* that can contribute to the strength of the topic argument. Although the Shapley-based contribution function is theoretically well-founded, it is significantly harder to compute than removal and gradient-based methods. Our RAEs are based on Shapley values as well, and we work around the complexity problem by proposing an approximation algorithm.

Other work focuses on restricted types of QBAFs. In particular, [Amgoud *et al.*, 2017] propose a contribution function for attack-only QBAFs (where arguments can only decrease the strength of the topic argument) and explain an argument’s strength by assigning attribution scores to its *direct* attacks. For instance, in Figure 1, assume again that α is the topic argument that needs to be explained. The spirit of the method is to attribute the strength of α to r_2 , r_5 and r_7 , which are directly connected to α .² Instead, in our RAEs, we take all edges into

²Here, although Figure 1 is not an attack-only QBAF, we focus on the spirit of the contribution function of [Amgoud *et al.*, 2017].

account because they all contribute to α . Let us take r_1 as an example, representing that β contributes to the strength of α by strengthening the supporter γ of α : the contribution of r_1 may even be greater than that of r_5 and r_7 , which means that *indirect* edges may play an important role.

3 Preliminaries

To begin with, we recall the definition of QBAFs. We focus on QBAFs with strength values in the domain $\mathbb{I} = [0, 1]$

Definition 1 (QBAF). A Quantitative Bipolar AF (QBAF) is a quadruple $\mathcal{Q} = \langle \mathcal{A}, \mathcal{R}^-, \mathcal{R}^+, \tau \rangle$ where:

- \mathcal{A} is a set of arguments;
- $\mathcal{R}^- \subseteq \mathcal{A} \times \mathcal{A}$ is a binary attack relation;
- $\mathcal{R}^+ \subseteq \mathcal{A} \times \mathcal{A}$ is a binary support relation;
- \mathcal{R}^- and \mathcal{R}^+ are disjoint;
- $\tau : \mathcal{A} \rightarrow \mathbb{I}$ is a function assigning base scores to arguments.

QBAFs are often denoted graphically (see Figure 1 as an example), where arguments are nodes and edges show the attack or support relations, labelled by $-$ and $+$, respectively. The base scores can be seen as apriori strengths of arguments when ignoring all other arguments (and is omitted from graphical representations, as in Figure 1). Seeing QBAFs as graphs allows us to use standard notions such as that of *path*.

In the remainder, unless specified otherwise, we assume as given a generic QBAF $\mathcal{Q} = \langle \mathcal{A}, \mathcal{R}^-, \mathcal{R}^+, \tau \rangle$ for $\mathbb{I} = [0, 1]$. Also, we let $\mathcal{R} = \mathcal{R}^- \cup \mathcal{R}^+$ and, for any $\alpha \in \mathcal{A}$, $\mathcal{R}(\alpha) = \{\beta \in \mathcal{A} \mid (\alpha, \beta) \in \mathcal{R}\}$ denotes the set of all outgoing edges from α .

Gradual semantics evaluate QBAFs by a function $\sigma : \mathcal{A} \rightarrow \mathbb{I}$ that assigns a (final) *strength* to every argument (e.g. see [Leite and Martins, 2011; Baroni *et al.*, 2015; Amgoud and Ben-Naim, 2018]). In most approaches, σ is defined via an iterative process that initializes all strength values with the base scores and then updates the strength values based on the strength of attackers and supporters. The final strength is the limit of this process. The process is guaranteed to converge for acyclic graphs after at most $n = |\mathcal{A}|$ iterations³ [Potyka, 2019] and, in practice, also quickly converges for cyclic graphs [Potyka, 2018]. Since we aim to explain the strength, which is only possible when it is defined, we will assume convergence for all arguments in the remainder, amounting to the following.

Definition 2 (Well-definedness). A gradual semantics σ is well-defined for \mathcal{Q} iff $\sigma(\alpha)$ exists for every $\alpha \in \mathcal{A}$.

Example 1. Consider the QBAF in Figure 2 where the base scores of all arguments are set to 0.5. Then, the *DF-QuAD* gradual semantics [Rago *et al.*, 2016], denoted by σ^{DF} , determines the following strengths:⁴ $\sigma^{DF}(\alpha) = 0.8046875$, $\sigma^{DF}(\beta) = \sigma^{DF}(\gamma) = 0.375$, $\sigma^{DF}(\delta) = 0.25$, $\sigma^{DF}(\zeta) = 0.5$.

We will often need to restrict QBAFs to a subset of the edges or change the base score function, as follows.

³Since the strength of an argument can only be affected by its parents, it actually suffices to update each argument only once by following a topological ordering of the arguments [Potyka, 2019].

⁴We omit details on how gradual semantics determine strengths, as we focus on explaining these strengths. For details see arxiv.org/abs/2404.14304.

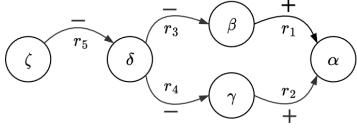


Figure 2: An example of QBAF (where all base scores are set to 0.5).

Definition 3. For $\mathcal{S} \subseteq \mathcal{R}$, let $\mathcal{Q}^{|\mathcal{S}|} = \langle \mathcal{A}, \mathcal{R}^- \cap \mathcal{S}, \mathcal{R}^+ \cap \mathcal{S}, \tau \rangle$. For $\tau' : \mathcal{A} \rightarrow \mathbb{I}$ a base score function, let $\mathcal{Q}^{|\tau'}| = \langle \mathcal{A}, \mathcal{R}^-, \mathcal{R}^+, \tau' \rangle$. Then, for any $\alpha \in \mathcal{A}$, we let $\sigma_{\mathcal{S}}(\alpha)$ denote the strength of α in $\mathcal{Q}^{|\mathcal{S}|}$ and $\sigma_{\tau'}(\alpha)$ denote the strength of α in $\mathcal{Q}^{|\tau'}|$.

For illustration, in Figure 2, suppose $\mathcal{S} = \{r_1, r_3, r_5\}$. Then r_2 and r_4 are not considered when computing $\sigma_{\mathcal{S}}(\alpha)$.

We will consider the following *monotonicity* property of gradual semantics, which is a variant of various notions proposed in the literature (see [Baroni *et al.*, 2019]) suitable for our setting. Roughly speaking, it states that base scores and relations monotonically influence arguments as one would intuitively expect.

Definition 4 (Monotonicity). A gradual semantics σ is monotonic iff for any $\alpha, \beta \in \mathcal{A}$ such that $\alpha \neq \beta$ and $\mathcal{R}(\beta) = \{(\beta, \alpha)\}$, for any $\tau' : \mathcal{A} \rightarrow \mathbb{I}$:

1. If $(\beta, \alpha) \in \mathcal{R}^-$, then $\sigma(\alpha) \leq \sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}(\alpha)$;
2. If $(\beta, \alpha) \in \mathcal{R}^+$, then $\sigma(\alpha) \geq \sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}(\alpha)$;
3. If $(\beta, \alpha) \in \mathcal{R}^-$, $\tau(\beta) \leq \tau'(\beta)$ and $\tau(\gamma) = \tau'(\gamma)$ for all $\gamma \in \mathcal{A} \setminus \{\beta\}$, then $\sigma(\alpha) \geq \sigma_{\tau'}(\alpha)$;
4. If $(\beta, \alpha) \in \mathcal{R}^+$, $\tau(\beta) \leq \tau'(\beta)$ and $\tau(\gamma) = \tau'(\gamma)$ for all $\gamma \in \mathcal{A} \setminus \{\beta\}$, then $\sigma(\alpha) \leq \sigma_{\tau'}(\alpha)$.

DF-QuAD, Quadratic Energy (QE) [Potyka, 2018], Restricted Euler-based (REB) [Amgoud and Ben-Naim, 2018] and all commonly considered gradual semantics in the literature satisfy monotonicity in acyclic QBAFs. We conjecture that these semantics are also monotonic for cyclic QBAFs.

Proposition 1. *DF-QuAD, QE, REB satisfy monotonicity in acyclic QBAFs.*

Conjecture 1. *DF-QuAD, QE, REB satisfy monotonicity in cyclic QBAFs.*

4 Relation Attribution Explanations

In order to explain the strength of a topic argument in a QBAF, we quantify the contribution of all edges to the topic argument. In order to find a fair and reasonable attribution method for quantifying these contributions, we build up on the Shapley-value as in [Amgoud *et al.*, 2017; Čyras *et al.*, 2022]. We define our *relation attribution explanations* as follows.

Definition 5 (RAEs). Let $\alpha \in \mathcal{A}$ be a topic argument and $r \in \mathcal{R}$. We define the Relation Attribution Explanation (RAE) from r to α under σ as:

$$\phi_{\sigma}^{\alpha}(r) = \sum_{\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}} \frac{(|\mathcal{R}| - |\mathcal{S}| - 1)! |\mathcal{S}|!}{|\mathcal{R}|!} [\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)].$$

Intuitively, $\phi_{\sigma}^{\alpha}(r)$ looks at every subset of edges (\mathcal{S}) and computes the marginal contribution of r

($[\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)]$). This marginal contribution is weighted by the probability that a random permutation of the edges starts with the subset (\mathcal{S}) and is followed by r . The main difference between our definition and that in [Amgoud *et al.*, 2017] lies in the potential “causes” of topic arguments. We attribute the strength of a topic argument to all edges (direct and indirect causes) in the QBAF while [Amgoud *et al.*, 2017] attribute it only to the directly incoming edges (direct causes). Furthermore, our definition is suitable not only for attacks but also for supports [Cayrol and Lagasque-Schiech, 2013], which are important in applications [Delobelle and Villata, 2019].

Qualitatively, we distinguish three different (*relation*) contributions based on the sign of $\phi_{\sigma}^{\alpha}(r)$.

Definition 6 ((Relation) Contribution). Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$.

1. If $\phi_{\sigma}^{\alpha}(r) > 0$, we say r has a positive contribution to α ;
2. If $\phi_{\sigma}^{\alpha}(r) < 0$, we say r has a negative contribution to α ;
3. If $\phi_{\sigma}^{\alpha}(r) = 0$, we say r has a neutral contribution to α .

Example 2 (Cont). Consider again the QBAF in Figure 2 under the DF-QuAD gradual semantics as in Example 1. Let α be the topic argument. We compute RAEs by Definition 5: $\phi_{\sigma}^{\alpha}(r_1) = 0.16875 > 0$, $\phi_{\sigma}^{\alpha}(r_2) = 0.16875 > 0$, $\phi_{\sigma}^{\alpha}(r_3) \approx -0.0318 < 0$, $\phi_{\sigma}^{\alpha}(r_4) \approx -0.0318 < 0$, $\phi_{\sigma}^{\alpha}(r_5) \approx 0.0307 > 0$. Hence, r_1 , r_2 and r_5 have a positive contribution to α while r_3 and r_4 have a negative one. We can also see that r_1 and r_2 have a more positive contribution to α than r_5 . We visualize the RAEs in Figure 3.

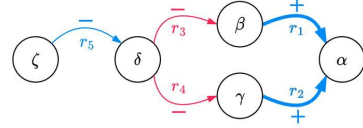


Figure 3: Contributions, drawn from RAEs, for topic argument α for the QBAF in Figure 2. (Blue/red edges denote *positive/negative* contributions, respectively. The thickness of edges represents the magnitude of their contributions, i.e. their RAE value.)

5 Properties

We now study some properties of RAEs. We start with *Shapley-based properties* that basically adapt to our setting properties of Shapley values and then move to *argumentative properties* that we deem interesting in our setting.

5.1 Shapley-based Properties

Similar to [Amgoud *et al.*, 2017], we first transfer the four basic properties of Shapley-values [Shapley, 1951] to our setting. *Efficiency* is recognized as a desirable property for attribution methods [Ancona *et al.*, 2017]. In our context, it states that the sum of all RAEs corresponds to the deviation of the topic argument’s strength $\sigma(\alpha)$ from its base score $\tau(\alpha)$ (namely the explanation distributes the responsibility for the difference among the edges). To prove this property, we assume that the semantics satisfies the *Stability* property [Amgoud and Ben-Naim, 2018], which states that the final strength of an argument is its base score whenever it has no incoming edges.

Proposition 2 (Efficiency). *If σ satisfies Stability, then for all $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$: $\sigma(\alpha) = \tau(\alpha) + \sum_{r \in \mathcal{R}} \phi_\sigma^\alpha(r)$.*

As an illustration, in Example 2, $\sigma^{DF}(\alpha) = 0.8046875$ equals to the the sum of $\tau(\alpha) = 0.5$ and the RAEs for all five edges in \mathcal{Q} : $\sum_{r \in \mathcal{R}} \phi_{\sigma^{DF}}^\alpha(r) = 0.3046875$.

Efficiency has an interesting implication that is called *Justification* in [Cyras et al., 2022]. This demands that whenever the strength of the topic argument differs from its base score, then there is a non-zero RAE explaining the difference.

Corollary 1 (Justification). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$.*

1. *If $\sigma(\alpha) > \tau(\alpha)$, then $\exists r \in \mathcal{R}$ such that $\phi_\sigma^\alpha(r) > 0$;*
2. *If $\sigma(\alpha) < \tau(\alpha)$, then $\exists r \in \mathcal{R}$ such that $\phi_\sigma^\alpha(r) < 0$.*

As an illustration, in Example 2, we have $\sigma^{DF}(\alpha) > \tau(\alpha)$ and r_1 with $\phi_{\sigma^{DF}}^\alpha(r_1) = 0.16875 > 0$ justifies the difference.

Dummy, also known as *Missingness* in [Lundberg and Lee, 2017], guarantees that if an edge does not make any contribution to the topic argument, then its RAE is 0.

Proposition 3 (Dummy). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$. If $\sigma_{S \cup \{r\}}(\alpha) = \sigma_S(\alpha)$ holds for all $S \subseteq \mathcal{R}$, then $\phi_\sigma^\alpha(r) = 0$.*

As an illustration, in Example 2, if we explain $\sigma^{DF}(\beta)$ (i.e. the topic argument is β), then $\phi_{\sigma^{DF}}^\beta(r_1) = 0$.

Symmetry states that if two edges share the same contribution to the topic argument, then their RAEs are equal.

Proposition 4 (Symmetry). *Let $\alpha \in \mathcal{A}$ and $r_i, r_j \in \mathcal{R}$ with $r_i \neq r_j$. If $\sigma_{S \cup \{r_i\}}(\alpha) = \sigma_{S \cup \{r_j\}}(\alpha)$ holds for any $S \subseteq \mathcal{R} \setminus \{r_i, r_j\}$, then $\phi_\sigma^\alpha(r_i) = \phi_\sigma^\alpha(r_j)$.*

As an illustration, in Example 2, r_1 and r_2 have symmetrical effects, thus $\phi_{\sigma^{DF}}^\alpha(r_1) = \phi_{\sigma^{DF}}^\alpha(r_2) = 0.1608$.

Dominance states that if one edge always makes a larger contribution than another, then this should be reflected in the magnitude of the RAE.

Proposition 5 (Dominance). *Let $\alpha \in \mathcal{A}$ and $r_i, r_j \in \mathcal{R}$ with $r_i \neq r_j$. If $\exists S' \subseteq \mathcal{R} \setminus \{r_i, r_j\}$ such that $\sigma_{S' \cup \{r_i\}}(\alpha) > \sigma_{S' \cup \{r_j\}}(\alpha)$ and $\forall S'' \subseteq \mathcal{R} \setminus \{r_i, r_j\}$ ($S' \neq S''$) such that $\sigma_{S'' \cup \{r_i\}}(\alpha) \geq \sigma_{S'' \cup \{r_j\}}(\alpha)$, then $\phi_\sigma^\alpha(r_i) > \phi_\sigma^\alpha(r_j)$.*

As an illustration, in Example 2, let $\tau(\beta) = 1.0$ while all other base scores remain 0.5. For $S' = \{r_3, r_4, r_5\}$ and $\forall S'' \subseteq \{r_3, r_4, r_5\}$ ($S' \neq S''$), we have $\sigma_{S' \cup \{r_1\}}(\alpha) > \sigma_{S' \cup \{r_2\}}(\alpha)$ and $\sigma_{S'' \cup \{r_1\}}(\alpha) \geq \sigma_{S'' \cup \{r_2\}}(\alpha)$, thus $\phi_{\sigma^{DF}}^\alpha(r_1) = 0.3375 > \phi_{\sigma^{DF}}^\alpha(r_2) = 0.1292$.

5.2 Argumentative Properties

We now study some argumentative properties that we deem interesting in our setting. When assessing these properties, we distinguish three edge types in QBAFs based on the form and number of paths to the topic argument.

Definition 7 (Edge Types). *Let $\alpha, \beta, \gamma \in \mathcal{A}$, $\alpha \neq \beta$. Then*

1. *(β, γ) is a direct edge w.r.t. α if $(\beta, \gamma) \in \mathcal{R}$ and there is only one path from γ to α in \mathcal{Q} (and $\gamma = \alpha$);*
2. *(β, γ) is an indirect edge w.r.t. α if there is only one path from γ to α in \mathcal{Q} (and $\gamma \neq \alpha$);*
3. *(β, γ) is a multifold edge w.r.t. α if there is more than one path from γ to α in \mathcal{Q} (and $\gamma \neq \alpha$).*

Example 3 (Cont). *Given the QBAF in Figure 2, r_1 and r_2 are direct edges w.r.t. α as they bring direct support to α ; r_3 and r_4 are indirect edges w.r.t. α as they are on single paths to α (while not bringing support or attack to it); r_5 is a multifold edge w.r.t. α because it starts two different paths (r_5, r_3, r_1 and r_5, r_4, r_2) to α .*

The first argumentative property is *Sign Correctness*, demanding that the sign of an edge reflects its polarity.

Property 1 (Sign Correctness). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$.*

1. *If $r \in \mathcal{R}^-$, then $\phi_\sigma^\alpha(r) \leq 0$;*
2. *If $r \in \mathcal{R}^+$, then $\phi_\sigma^\alpha(r) \geq 0$.*

Naturally, sign correctness cannot be satisfied if the gradual semantics does not behave in the intended way, but it does so if it satisfies monotonicity, for direct edges. Note that we are defining this and later properties wrt. specific arguments and edges, considering then their satisfaction wrt. classes of edges, e.g. direct edges as in the next result.

Proposition 6. *Let r be a direct edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies sign correctness if σ satisfies monotonicity.*

For indirect edges, we need to make several case differentiations, as the meaning of edges can be inverted along paths (e.g. an attacker of an attacker, actually serves as a supporter).

Proposition 7. *Let r be an indirect edge w.r.t. α . Suppose the path sequence from r to α is r, r_1, \dots, r_n ($n \geq 1$). Let $\lambda = |\{r_1, \dots, r_n\} \cap \mathcal{R}^-|$. Then the following statements hold if σ satisfies monotonicity.*

1. *If $r \in \mathcal{R}^-$ and λ is odd, then $\phi_\sigma^\alpha(r) \geq 0$;*
2. *If $r \in \mathcal{R}^-$ and λ is even, then $\phi_\sigma^\alpha(r) \leq 0$;*
3. *If $r \in \mathcal{R}^+$ and λ is odd, then $\phi_\sigma^\alpha(r) \leq 0$;*
4. *If $r \in \mathcal{R}^+$ and λ is even, then $\phi_\sigma^\alpha(r) \geq 0$.*

Example 4 (Cont). *Consider the QBAF in Figure 2 and the RAEs in Example 2. $r_1 \in \mathcal{R}^+$ is a direct edge w.r.t. α , hence $\phi_{\sigma^{DF}}^\alpha(r_1) \geq 0$; while $r_3 \in \mathcal{R}^-$ is an indirect edge w.r.t. α , and λ is 0 (even), hence $\phi_{\sigma^{DF}}^\alpha(r_3) \leq 0$.*

Essentially, these results show that RAEs correctly explain the behavior of direct and indirect edges under monotonicity. For multifold edges, however, monotonicity may not help.

Proposition 8. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate sign correctness even if σ satisfies monotonicity.*

Counterfactuality is a natural property which states that the strength of a topic argument will not be increased (decreased) if an edge with positive (negative) contribution is removed.

Property 2 (Counterfactuality). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$.*

1. *If $\phi_\sigma^\alpha(r) < 0$, then $\sigma(\alpha) \leq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$;*
2. *If $\phi_\sigma^\alpha(r) > 0$, then $\sigma(\alpha) \geq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$.*

Proposition 9. *Let r be a direct or indirect edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies counterfactuality if σ satisfies monotonicity.*

Example 5 (Cont). *Consider the QBAF in Figure 2 and the RAEs in Example 2. r_1 is a direct edge w.r.t. α and $\phi_{\sigma^{DF}}^\alpha(r_1) > 0$. If r_1 is removed, then $\sigma^{DF}(\alpha) = 0.8046875 > \sigma_{\mathcal{R} \setminus \{r_1\}}^{DF}(\alpha) = 0.6875$. r_3 is an indirect edge w.r.t. α and $\phi_{\sigma^{DF}}^\alpha(r_3) < 0$. If r_3 is removed, then $\sigma^{DF}(\alpha) = 0.8046875 < \sigma_{\mathcal{R} \setminus \{r_3\}}^{DF}(\alpha) = 0.84375$.*

Proposition 10. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate counterfactuality even if σ satisfies monotonicity.*

From a debugging angle, it is worth exploring how the RAE can be adjusted by a user. We find that the RAE of an edge (β, γ) is closely related to the base score of its *source argument* β , in the sense of the properties of *Qualitative Invariability* and *Quantitative Variability* defined below.

Qualitative Invariability states that an edge with positive RAE will never make a negative contribution to the topic argument even if the base score of its source argument changes.

Property 3 (Qualitative Invariability). *Let $\alpha, \beta \in \mathcal{A}$ and $r \in \mathcal{R}(\beta)$. Let ϕ_δ denote $\phi_\sigma^\alpha(r)$ when setting $\tau(\beta)$ to some $\delta \in \mathbb{I}$.*

1. *If $\phi_\sigma^\alpha(r) < 0$, then $\forall \delta \in \mathbb{I}, \phi_\delta \leq 0$;*
2. *If $\phi_\sigma^\alpha(r) > 0$, then $\forall \delta \in \mathbb{I}, \phi_\delta \geq 0$.*

Proposition 11. *Let r be a direct or indirect edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies qualitative invariability if σ satisfies monotonicity.*

Example 6 (Cont). *Consider the QBAF in Figure 2 and the RAEs in Example 2. Since r_1 is a direct edge w.r.t. α and $\phi_{\sigma_{DF}}^\alpha(r_1) > 0$, then even if $\tau(\beta)$ is changed to some other value δ , the new RAE $\phi_\delta \geq 0$ still holds by Proposition 11.*

Proposition 12. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate qualitative invariability even if σ satisfies monotonicity.*

Quantitative variability states that the RAE of an edge will not be increased (decreased) if the base score of its source argument is decreased (increased).

Property 4 (Quantitative Variability). *Let $\alpha, \beta \in \mathcal{A}$, $\alpha \neq \beta$, and $\mathcal{R}(\beta) = \{r\}$. Let ϕ_δ denote $\phi_\sigma^\alpha(r)$ when setting $\tau(\beta)$ to some $\delta \in \mathbb{I}$.*

1. *If $\delta < \tau(\beta)$, then $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$;*
2. *If $\delta > \tau(\beta)$, then $|\phi_\delta| \geq |\phi_\sigma^\alpha(r)|$.*

We assess the satisfaction of this property when r is a direct, indirect or manifold edge w.r.t. the topic argument.

Proposition 13. *Let r be a direct or indirect edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies quantitative variability if σ satisfies monotonicity.*

Example 7 (Cont). *Consider the QBAF in Figure 2 and the RAEs in Example 2. Since r_1 is a direct edge w.r.t. α and $\phi_{\sigma_{DF}}^\alpha(r_1) > 0$, then if $\tau(\beta)$ is increased to some $\delta > \tau(\beta)$, the new RAE ϕ_δ will not decrease by Proposition 13.*

Proposition 14. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate quantitative variability even if σ satisfies monotonicity.*

Let us note that many properties may be violated for the multifold case even if the underlying gradual semantics is monotonic. This is because monotonicity only guarantees the *direct* effects of attacks and supports on an argument. Since a single edge can be involved in a large number of paths to the topic argument, one cannot make a reasonable demand about its effect without a long list of case differentiations. We therefore focus on the special cases where there is only a single path from the argument under investigation to the topic argument (i.e. the cases with direct and indirect edges).

Algorithm 1 An Approximation Algorithm for RAEs

Input: A QBAF $\mathcal{Q} = \langle \mathcal{A}, \mathcal{R}^-, \mathcal{R}^+, \tau \rangle$; a gradual semantics σ ; a topic argument α ; the sample size N .

Output: Approximate RAEs *attribution_dict*.

```

1: attribution_dict  $\leftarrow \{\}$            % empty dictionary
2: for  $r$  in  $\mathcal{R}$  do
3:   sum  $\leftarrow 0$ 
4:   while  $N > 0$  do
5:      $\mathcal{S} \leftarrow \text{random\_sample}(\mathcal{R} \setminus \{r\})$ 
6:     sum  $\leftarrow \text{sum} + [\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)]$ 
7:      $N \leftarrow N - 1$ 
8:   attribution_dict[ $r$ ]  $\leftarrow \text{sum}/N$ 
9: return attribution_dict

```

6 Approximating RAEs Probabilistically

Here, we look at how we can compute RAEs efficiently.

Computing RAEs involves computing the final strength values of arguments. The runtime for computing these values depends on the graph structure of the QBAF and on the gradual semantics. Let $n = |\mathcal{R}|$ and $m = |\mathcal{A}|$. The strength values can be computed in linear time $\mathcal{O}(m + n)$ for acyclic QBAFs [Potyka, 2019, Proposition 3.1]. For some pathological examples of cyclic QBAFs, the strength computation may not converge resulting in infinite runtime [Mossakowski and Neuhaus, 2018; Potyka, 2019]. However, for randomly generated cyclic QBAFs, the strength values typically converge in subquadratic time [Potyka, 2018, Figure 7]. In fact, if the outdegree of arguments in the QBAF is not too large, the strength values are guaranteed to converge in linear time [Potyka, 2019, Proposition 3.3]. To avoid a large number of case differentiations, we just denote the runtime for computing strength values by $T(m, n)$ in the following.

Concerning RAEs, let us first note that we can compute them exactly in exponential time by inspecting all subsets of edges (excluding cases where determining strength values fails to converge).

Proposition 15 (Computing RAEs exactly). *RAEs can be computed in time $\mathcal{O}(n \cdot 2^n \cdot T(m, n))$.*

For larger QBAFs, we can apply approximation methods. It is folklore in algorithmic game theory that Shapley values can be seen as expected values. Under this view, the probability of a subset $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$ is defined by $P(\mathcal{S}) = \frac{(|\mathcal{R}| - |\mathcal{S}| - 1)! |\mathcal{S}|!}{|\mathcal{R}|!}$ and the *marginal contribution* of \mathcal{S} is defined by the function $m(\mathcal{S}) = \sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)$. Our RAEs then correspond to the expected value $E_P[m]$ of the marginal contribution function under P . This interpretation allows approximating the Shapley values by the approximation Algorithm 1, returning a dictionary *attribution_dict* used iteratively to accumulate pairs assigning estimate values to edges. By the *Law of Large Numbers* (e.g., Theorem 5.1 in [Gut and Gut, 2005]), the estimates converge in probability to the true Shapley values. That is, for every $\epsilon > 0$, the probability that the estimates deviate by more than ϵ from the true value approaches 0 as the number of samples approaches infinity. By the *Central Limit Theorem* (e.g., Theorem 5.2 in [Gut and Gut, 2005]), the distribution of the samples approaches a normal distribution with mean

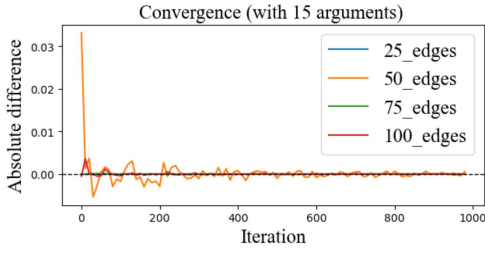


Figure 4: Convergence of Algorithm 1 for random cyclic QBAFs with 15 arguments and various numbers of edges.

equal to the true Shapley values. This means, in particular, that the estimator is unbiased. However, the variance can be quite large and is better evaluated empirically. We thus do not give a precise formula and simply state the following guarantees.

Proposition 16. *The estimates generated by Algorithm 1 converge in probability to the true Shapley values.*

Proposition 17 (Approximating RAEs). *If the number of samples for each edge is N , then approximate RAEs can be generated in time $\mathcal{O}(n \cdot N \cdot T(m, n))$.*

In particular, when the QBAF is acyclic or meets the conditions on the outdegree of arguments in [Potyka, 2019, Proposition 3.3], we can compute RAEs in time $\mathcal{O}(n \cdot N \cdot (n+m))$.

Proposition 16 guarantees that Algorithm 1 converges to the true RAEs but does not tell us how many iterations we require to reach a good approximation. In order to evaluate the convergence speed empirically, we conducted experiments with randomly generated QBAFs of increasing size. Figure 4 shows, for cyclic QBAFs (see arxiv.org/abs/2404.14304 for acyclic QBAFs), how the absolute difference (y-axis) between estimates at every 10-th iteration evolves with an increasing number of samples (x-axis), pointing to convergence within a few hundreds iterations. For each iteration, it approximately took 14ms and 0.9ms for cyclic and acyclic QBAFs, respectively, with 15 arguments and 25 edges. We give hardware specifications and additional experiments for runtime, acyclic and differently-sized QBAFs in arxiv.org/abs/2404.14304.

7 Case studies

Finally, we carry out two case studies, including a large QBAF and a non-tree QBAF, to show some practical use of our RAEs.

7.1 Case Study 1: Fraud Detection

Background Automatic fraud detection plays an important role in e-commerce. [Chi *et al.*, 2021] propose to use QBAFs for fraud detection because of their intrinsic interpretability. We take the QBAF from [Chi *et al.*, 2021], shown in Figure 5, where argument 1 (‘It is a fraud case’) is the topic argument, and arguments 2 – 48 represent evidence for or against this case. The specific content of these arguments can be found in arxiv.org/abs/2404.14304 or in the original [Chi *et al.*, 2021].

Settings We set the base score for each argument to 0.5 in line with [Chi *et al.*, 2021]. Since we do not consider in this paper edge-weighted argumentation [Dunne *et al.*, 2011], we apply DF-QuAD semantics here instead of O-QuAD [Chi *et al.*, 2021] which is a variant of DF-QuAD

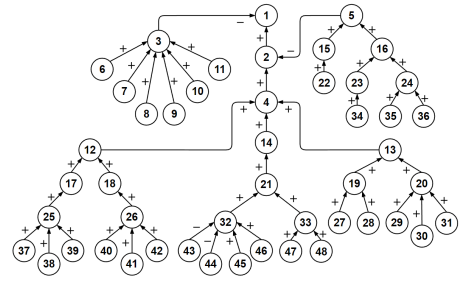


Figure 5: Fraud Detection example from [Chi *et al.*, 2021].

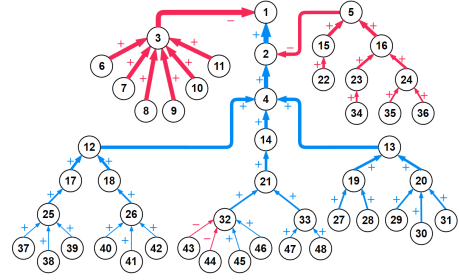


Figure 6: Contributions and RAEs for the Fraud Detection example. (We use conventions as described in the caption of Figure 3).

with weights on edges. The given case is considered fraud if and only if $\sigma^{DF}(1) > \tau(1) = 0.5$. Under DF-QuAD, we have $\sigma^{DF}(1) \approx 0.2544 < \tau(1)$, which means the case is not considered fraud. Since there are 47 edges in Figure 5, computing RAEs exactly is prohibitively expensive. Thus, we apply the approximate Algorithm 1, setting the sample size N to 1000. We chose N experimentally to be large enough to guarantee that the estimates converge.⁵

Explanations We apply our RAEs and the contributions derived from them to give quantitative explanations for $\sigma^{DF}(1)$ (see Figure 6; for more details see arxiv.org/abs/2404.14304).

Figure 6 shows that 18 red edges make a negative contribution while 29 blue edges make a positive contribution to argument 1, and negative contributions overwhelm the positive ones. Among the positive contributions, (2, 1) makes the largest, with $\text{RAE} = 2.55 \cdot 10^{-1}$, because it directly supports argument 1. (40, 26) makes the smallest contribution with $\text{RAE} = 2.83 \cdot 10^{-5}$ because it is indirect and far away from argument 1. Among the negative contributions, (3, 1) makes the largest ($\text{RAE} = -4.56 \cdot 10^{-1}$) since argument 3 directly attacks argument 1, whereas (43, 32) makes the smallest ($\text{RAE} = -5.84 \cdot 10^{-5}$) as there is only one (odd number) attack from argument 43 to argument 1.

In Figure 6, edges close to argument 1 make a greater contribution than those further away. This is because removing close edges will also remove their predecessors on the path to the topic argument. The RAEs of edges such as (4, 2) and (12, 4) show they also play a role, which is different from the

⁵As the additional experiments for large QBAFs in arxiv.org/abs/2404.14304 show, the estimates converge typically after a few hundreds iterations even with more than a thousand edges, so $N = 1000$ here is appropriate.

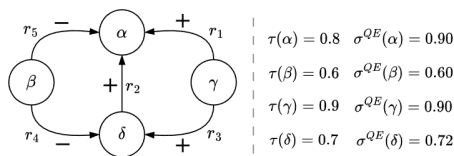


Figure 7: LLMs example.

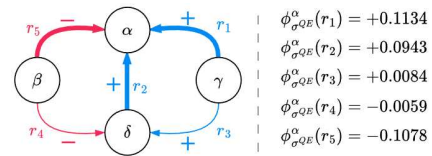


Figure 8: Contributions and RAEs for the LLMs example .

contribution function in [Amgoud *et al.*, 2017].

Note that direct edges do not always make greater contribution than those further away, especially in multifold scenarios. In such cases, we believe explaining the strength by considering all edges rather than only direct edges is a better choice.

7.2 Case Study 2: Large Language Models (LLMs)

Background and Settings LLMs’ ability to process and generate text can contribute to the development of various AI models [Naveed *et al.*, 2023] and help address the knowledge acquisition bottleneck. The idea is that we can query an LLM with a particular claim and use the answers to build up a QBAF. The QBAF can then be used to visualize (potentially contradictory) arguments that the LLM generated and compute final strength of these arguments, and RAEs can be used to explain the relevance of particular edges for the strength of the claim (seen as the topic argument). To present our RAE approach with QBAFs containing more intricate relationships among arguments than the simple tree-like structure resulting in case study 1, we force a certain structure and generate a non-tree QBAF by ChatGPT(GPT-3.5) [OpenAI, 2022], for the claim ‘*It is easy for children to learn a foreign language well*’ (topic argument α), prompted to create arguments satisfying the following requirements:⁶

1. Provide one argument β attacking α and two arguments γ and δ supporting α .
2. Let β and γ attack and support δ , respectively.
3. Give confidence for all arguments, ranging from 0 to 1.

We obtained the following arguments and confidence values (which we use as base scores).

β (0.6): *Learning a foreign language requires cognitive maturity, which children lack. Hence, it’s difficult for them to excel.*

γ (0.9): *Studies show that young children possess higher neuroplasticity, making language learning more effective.*

δ (0.7): *Children immersed in a foreign language environment from an early age have better language acquisition.*

We used the QE semantics (σ^{QE}) to compute the strength of arguments and visualize the QBAF and strengths in Figure 7.

Explanations Figure 8 visualises RAEs and contributions and gives a ranking of the edges based on their contribution. There are two paths from β to α : $p_1 = r_5$ and $p_2 = r_4, r_2$. The cumulative contributions of p_1 and p_2 are -0.1078 and

0.0884 , respectively, obtained by adding up the RAEs on the path. Thus, p_1 and p_2 make, respectively, a negative and positive contribution to α . Also, p_1 makes a greater contribution when considering absolute values. Although p_2 positively contributes to α , we find r_4 makes a negative contribution on this path, which is not obvious if we only compute argument-based attributions. Indeed, we believe that RAEs are better suited than argument-based attribution explanations in this scenario because they provide a deeper insight into the way arguments affect one another along different reasoning paths.

Property Verification Let us check the satisfaction of some properties introduced previously, under σ^{QE} used in this case study. First, the sum of all RAEs (0.10) corresponds to the deviation from $\tau(\alpha) = 0.80$ to $\sigma^{QE}(\alpha) = 0.90$, which satisfies *efficiency*. β directly attacks α so r_5 has a negative RAE while γ directly supports α thus r_1 has a positive RAE by *sign correctness*. According to *counterfactuality*, if r_1 is removed, then $\sigma^{QE}(\alpha)$ will decrease (to 0.80). If the $\tau(\gamma)$ is increased from 0.9 to 0.95, then $\phi_{\sigma^{QE}}^{\alpha}(r_1)$ and $\phi_{\sigma^{QE}}^{\alpha}(r_3)$ are still positive by *qualitative invariability*, and $\phi_{\sigma^{QE}}^{\alpha}(r_1)$ increases from 0.1134 to 0.1182 and $\phi_{\sigma^{QE}}^{\alpha}(r_3)$ increases from 0.008389 to 0.008438 by *quantitatively variability*.

8 Conclusion

We introduced RAEs to quantitatively explain the role of attack and support relations under gradual semantics for QBAFs, resulting in more fine-grained insights into the contribution of arguments, along different reasoning paths, than argument-based attribution explanations. We proposed several properties for RAEs, including some adapted from properties of Shapley values and some defined ex-novo. The satisfaction and violation of these properties theoretically shows that our RAEs are reasonable and faithful explanations, which is crucial to explanation methods. We also proposed an efficient probabilistic algorithm to approximate RAEs, proved theoretical convergence guarantees and demonstrated experimentally that it converges quickly. Finally, we carried out two case studies to evaluate and show the practical use of our RAEs.

Our work paves the way to many future directions. First, it would be interesting to explore joint Shapley values [Zhang *et al.*, 2021] for sets of attacks and supports and to investigate interactions among edges. Second, it would be worth exploring formal relationships between RAEs and argument-based attribution explanations. Third, it would be interesting to generalize our RAEs to edge-weighted QBAFs [Amgoud *et al.*, 2017]. Lastly, it would be important to carry out user studies as explanations should be easily understood and accepted by human users [Chen *et al.*, 2022].

⁶The prompt and response are given in arxiv.org/abs/2404.14304.

Acknowledgments

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Supplementary Material for

“Explaining Arguments’ Strength: Unveiling the Role of Attacks and Supports”

A Computational Details for Gradual Semantics

A.1 DF-QuAD Semantics

For any argument $\alpha \in \mathcal{A}$, $\sigma^{DF}(\alpha)$ is defined as follows:

$$\sigma^{DF}(\alpha) = \begin{cases} \tau(\alpha) - \tau(\alpha) \cdot (v_{\alpha-att} - v_{\alpha-sup}) & \text{if } v_{\alpha-att} \geq v_{\alpha-sup} \\ \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v_{\alpha-sup} - v_{\alpha-att}) & \text{if } v_{\alpha-att} < v_{\alpha-sup} \end{cases}$$

where $v_{\alpha-att}$ is the aggregation strength of all attackers against α , while $v_{\alpha-sup}$ is the aggregation strength of supporters for α . $v_{\alpha-att}$ and $v_{\alpha-sup}$ are defined as follows:

$$v_{\alpha-att} = 1 - \prod_{\{\beta \in \mathcal{A} | (\beta, \alpha) \in \mathcal{R}^-\}} (1 - \sigma^{DF}(\beta));$$

$$v_{\alpha-sup} = 1 - \prod_{\{\beta \in \mathcal{A} | (\beta, \alpha) \in \mathcal{R}^+\}} (1 - \sigma^{DF}(\beta)).$$

A.2 QE Semantics

For any argument $\alpha \in \mathcal{A}$, $\sigma^{QE}(\alpha)$ is defined as follows:

$$\sigma^{QE}(\alpha) = \begin{cases} \tau(\alpha) - \tau(\alpha) \cdot \frac{E_\alpha^2}{1+E_\alpha^2} & \text{if } E_\alpha \leq 0 \\ \tau(\alpha) + (1 - \tau(\alpha)) \cdot \frac{E_\alpha^2}{1+E_\alpha^2} & \text{if } E_\alpha > 0 \end{cases}$$

where E_α is the aggregation strength of all attackers and supporters of α . E_α is defined as follows:

$$E_\alpha = \sum_{\{\beta \in \mathcal{A} | (\beta, \alpha) \in \mathcal{R}^+\}} \sigma^{QE}(\beta) - \sum_{\{\beta \in \mathcal{A} | (\beta, \alpha) \in \mathcal{R}^-\}} \sigma^{QE}(\beta).$$

A.3 REB Semantics

For any argument $\alpha \in \mathcal{A}$, $\sigma^{REB}(\alpha)$ is defined as follows:

$$\sigma^{REB}(\alpha) = 1 - \frac{1 - \tau^2(\alpha)}{1 + \tau(\alpha) \cdot e^{E_\alpha}}$$

where E_α is the aggregation strength of all attackers and supporters of α . E_α is defined as follows:

$$E_\alpha = \sum_{\{\beta \in \mathcal{A} | (\beta, \alpha) \in \mathcal{R}^+\}} \sigma^{REB}(\beta) - \sum_{\{\beta \in \mathcal{A} | (\beta, \alpha) \in \mathcal{R}^-\}} \sigma^{REB}(\beta).$$

B Proofs for Section 3

Proposition 1. *DF-QuAD, QE, REB satisfy monotonicity in acyclic QBAFs.*

Proof. Proof for **DF-QuAD** semantics:

Let $v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-att}$ and $v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-sup}$ denote the aggregation value of α 's attackers and supporters after removing (β, α) , respectively. Let $v'_{\alpha-sup}$ and $v'_{\alpha-att}$ denote the aggregation value of α 's attackers and supporters based on the base score function τ' , respectively.

1. Since $\mathcal{R}(\beta) = \{(\beta, \alpha)\}$, if edge $(\beta, \alpha) \in \mathcal{R}^-$ is removed, then α loses an attacker β and the strength of all other attackers and supporters of α remain. Hence, $v_{\alpha-sup} = v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-sup}$ and $v_{\alpha-att} \geq v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-att}$ according to the DF-QuAD in Section A.1.

In order to see how $\sigma^{DF}(\alpha)$ changes, we need to discuss the relative magnitude of $v_{\alpha-att}$ and $v_{\alpha-sup}$. There are three cases:

(a) $v_{\alpha-att} \geq v_{\alpha-sup}$ and $v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-att} \geq v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-sup}$.

In this case, $\sigma^{DF}(\alpha) = \tau(\alpha) - \tau(\alpha) \cdot (v_{\alpha-att} - v_{\alpha-sup})$ and $\sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha) = \tau(\alpha) - \tau(\alpha) \cdot (v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-att} - v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-sup})$, therefore, we have $\sigma^{DF}(\alpha) \geq \sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha)$.

(b) $v_{\alpha-att} \geq v_{\alpha-sup}$ and $v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-att} \leq v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-sup}$.

In this case, $\sigma^{DF}(\alpha) = \tau(\alpha) - \tau(\alpha) \cdot (v_{\alpha-att} - v_{\alpha-sup})$ and $\sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha-sup} -$

$v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha - att}$, therefore, we have $\sigma^{DF}(\alpha) \geq \sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha)$.

(c) $v_{\alpha - att} \leq v_{\alpha - sup}$ and $v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha - att} \leq v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha - sup}$.

In this case, $\sigma^{DF}(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v_{\alpha - sup} - v_{\alpha - att})$ and $\sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha - sup} - v_{\mathcal{R} \setminus \{(\beta, \alpha)\} \alpha - att})$, therefore, we have $\sigma^{DF}(\alpha) \geq \sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha)$.

In summary, in any of these three cases, we know that if $(\beta, \alpha) \in \mathcal{R}^-$, then $\sigma^{DF}(\alpha) \leq \sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha)$.

2. Analogously to the proof in 1.

3. Since $\mathcal{R}(\beta) = \{(\beta, \alpha)\}$ and $(\beta, \alpha) \in \mathcal{R}^-$, if $\tau(\beta) \leq \tau'(\beta)$ and $\tau(\gamma) = \tau'(\gamma)$ for all $\gamma \in \mathcal{A} \setminus \{\beta\}$, then $\sigma^{DF}(\beta) \leq \sigma_{\tau'}^{DF}(\beta)$ and $\sigma^{DF}(\gamma) = \sigma_{\tau'}^{DF}(\gamma)$ for all $\gamma \in \mathcal{A} \setminus \{\alpha, \beta\}$. Thus, $v'_{\alpha - att} \geq v_{\alpha - att}$ and $v'_{\alpha - sup} = v_{\alpha - sup}$.

Similarly, in order to see how $\sigma^{DF}(\alpha)$ changes, we need to discuss the relative magnitude of $v_{\alpha - att}$ and $v_{\alpha - sup}$. There are three cases:

(a) $v_{\alpha - att} \geq v_{\alpha - sup}$ and $v'_{\alpha - att} \geq v'_{\alpha - sup}$.

In this case, $\sigma^{DF}(\alpha) = \tau(\alpha) - \tau(\alpha) \cdot (v_{\alpha - att} - v_{\alpha - sup})$ and $\sigma_{\tau'}^{DF}(\alpha) = \tau(\alpha) - \tau(\alpha) \cdot (v'_{\alpha - att} - v'_{\alpha - sup})$, therefore, we have $\sigma^{DF}(\alpha) \geq \sigma_{\tau'}^{DF}(\alpha)$.

(b) $v_{\alpha - att} \geq v_{\alpha - sup}$ and $v'_{\alpha - att} \leq v'_{\alpha - sup}$.

In this case, $\sigma^{DF}(\alpha) = \tau(\alpha) - \tau(\alpha) \cdot (v_{\alpha - att} - v_{\alpha - sup})$ and $\sigma_{\tau'}^{DF}(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v'_{\alpha - sup} - v'_{\alpha - att})$, therefore, we have $\sigma^{DF}(\alpha) \geq \sigma_{\tau'}^{DF}(\alpha)$.

(c) $v_{\alpha - att} \leq v_{\alpha - sup}$ and $v'_{\alpha - att} \leq v'_{\alpha - sup}$.

In this case, $\sigma^{DF}(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v_{\alpha - sup} - v_{\alpha - att})$ and $\sigma_{\tau'}^{DF}(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v'_{\alpha - sup} - v'_{\alpha - att})$, therefore, we have $\sigma^{DF}(\alpha) \geq \sigma_{\tau'}^{DF}(\alpha)$.

In summary, in any of these three cases, we know that if $(\beta, \alpha) \in \mathcal{R}^-$, $\tau(\beta) \leq \tau'(\beta)$ and $\tau(\gamma) = \tau'(\gamma)$ for all $\gamma \in \mathcal{A} \setminus \{\beta\}$, then $\sigma^{DF}(\alpha) \leq \sigma_{\mathcal{R} \setminus \{(\beta, \alpha)\}}^{DF}(\alpha)$.

4. Analogously to the proof in 3.

The proof for **QE** and **REB** semantics are analogous to the proof for DF-QuAD semantics. \square

C Proofs for Section 5

C.1 Shapley-based Properties

Proposition 2 (Efficiency). *If σ satisfies Stability, then for all $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$:*

$$\sigma(\alpha) = \tau(\alpha) + \sum_{r \in \mathcal{R}} \phi_{\sigma}^{\alpha}(r).$$

Proof. According to Definition 5, for any $r \in \mathcal{R}$ we have $\phi_{\sigma}^{\alpha}(r) = \sum_{S \subseteq \mathcal{R} \setminus \{r\}} \frac{(|\mathcal{R}| - |S| - 1)! |S|!}{|\mathcal{R}|!} [\sigma_{S \cup \{r\}}(\alpha) - \sigma_S(\alpha)]$. This definition is the compact form of the Shapley value. In order to prove the claim, we use the equivalent permutation representation [Maschler *et al.*, 2013].

$$\phi_{\sigma}^{\alpha}(r) = \frac{1}{|\mathcal{R}|!} \sum_{\pi \in \Pi(\mathcal{R})} [\sigma_{\pi^{-1}(r) \cup \{r\}}(\alpha) - \sigma_{\pi^{-1}(r)}(\alpha)], \quad (1)$$

where $\Pi(\mathcal{R})$ denotes the set of all permutations of \mathcal{R} and, for all $r \in \mathcal{R}$, $\pi^{-1}(r)$ denotes the set of edges in \mathcal{R} which precede r in the order π . Then, we have

$$\begin{aligned} \sum_{r \in \mathcal{R}} \phi_{\sigma}^{\alpha}(r) &= \sum_{r \in \mathcal{R}} \frac{1}{|\mathcal{R}|!} \sum_{\pi \in \Pi(\mathcal{R})} [\sigma_{\pi^{-1}(r) \cup \{r\}}(\alpha) - \sigma_{\pi^{-1}(r)}(\alpha)] \\ &= \frac{1}{|\mathcal{R}|!} \sum_{r \in \mathcal{R}} \sum_{\pi \in \Pi(\mathcal{R})} [\sigma_{\pi^{-1}(r) \cup \{r\}}(\alpha) - \sigma_{\pi^{-1}(r)}(\alpha)] \\ &= \frac{1}{|\mathcal{R}|!} \sum_{\pi \in \Pi(\mathcal{R})} \sum_{r \in \mathcal{R}} [\sigma_{\pi^{-1}(r) \cup \{r\}}(\alpha) - \sigma_{\pi^{-1}(r)}(\alpha)] \end{aligned} \quad (2)$$

Most terms in the inner sum cancel. To see this, consider an arbitrary permutation π and, for ease of notation, assume that it

induces the ordering (r_1, r_2, \dots, r_n) . Then, we have

$$\begin{aligned}
\sum_{r \in \mathcal{R}} [\sigma_{\pi^{-1}(r) \cup \{r\}}(\alpha) - \sigma_{\pi^{-1}(r)}(\alpha)] &= \sigma_{\{r_1\}}(\alpha) - \sigma_{\emptyset}(\alpha) \\
&+ \sigma_{\{r_1, r_2\}}(\alpha) - \sigma_{\{r_1\}}(\alpha) \\
&+ \sigma_{\{r_1, r_2, r_3\}}(\alpha) - \sigma_{\{r_1, r_2\}}(\alpha) \\
&+ \dots \\
&+ \sigma_{\{r_1, r_2, r_3, \dots, r_{n-1}\}}(\alpha) - \sigma_{\{r_1, r_2, r_3, \dots, r_{n-2}\}}(\alpha) \\
&+ \sigma_{\{r_1, r_2, r_3, \dots, r_n\}}(\alpha) - \sigma_{\{r_1, r_2, r_3, \dots, r_{n-1}\}}(\alpha) \\
&= \sigma_{\{r_1, r_2, r_3, \dots, r_n\}}(\alpha) - \sigma_{\emptyset}(\alpha) \\
&= \sigma(\alpha) - \tau(\alpha),
\end{aligned} \tag{3}$$

where the last equality follows from *Stability*. Since there are $|\mathcal{R}|!$ different orderings and π is taken arbitrarily, we have

$$\begin{aligned}
\sum_{r \in \mathcal{R}} \phi_{\sigma}^{\alpha}(r) &= \frac{1}{|\mathcal{R}|!} \sum_{\pi \in \Pi(\mathcal{R})} [\sigma(\alpha) - \tau(\alpha)] \\
&= \frac{1}{|\mathcal{R}|!} |\mathcal{R}|! [\sigma(\alpha) - \tau(\alpha)] \\
&= \sigma(\alpha) - \tau(\alpha)
\end{aligned} \tag{4}$$

Thus, $\sigma(\alpha) = \tau(\alpha) + \sum_{r \in \mathcal{R}} \phi_{\sigma}^{\alpha}(r)$, which completes the proof. \square

Corollary 1 (Justification). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$.*

1. *If $\sigma(\alpha) > \tau(\alpha)$, then $\exists r \in \mathcal{R}$ such that $\phi_{\sigma}^{\alpha}(r) > 0$;*
2. *If $\sigma(\alpha) < \tau(\alpha)$, then $\exists r \in \mathcal{R}$ such that $\phi_{\sigma}^{\alpha}(r) < 0$.*

Proof. 1. According to Proposition 2, $\sigma(\alpha) > \tau(\alpha)$ means $\sum_{r \in \mathcal{R}} \phi_{\sigma}^{\alpha}(r) > 0$, hence there must be at least one $r \in \mathcal{R}$ such that $\phi_{\sigma}^{\alpha}(r) > 0$.

2. According to Proposition 2, $\sigma(\alpha) < \tau(\alpha)$ means $\sum_{r \in \mathcal{R}} \phi_{\sigma}^{\alpha}(r) < 0$, hence there must be at least one $r \in \mathcal{R}$ such that $\phi_{\sigma}^{\alpha}(r) < 0$. \square

Proposition 3 (Dummy). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$. If $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) = \sigma_{\mathcal{S}}(\alpha)$ holds for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$, then*

$$\phi_{\sigma}^{\alpha}(r) = 0.$$

Proof. According to Definition 5, if $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) = \sigma_{\mathcal{S}}(\alpha)$ holds for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$, then $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha) = 0$. Thus, $\sum_{\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}} \frac{(|\mathcal{R}| - |\mathcal{S}| - 1)! |\mathcal{S}|!}{|\mathcal{R}|!} [\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)] = 0$, hence $\phi_{\sigma}^{\alpha}(r) = 0$, which completes the proof. \square

Proposition 4 (Symmetry). *Let $\alpha \in \mathcal{A}$ and $r_i, r_j \in \mathcal{R}$ with $r_i \neq r_j$. If $\sigma_{\mathcal{S} \cup \{r_i\}}(\alpha) = \sigma_{\mathcal{S} \cup \{r_j\}}(\alpha)$ holds for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r_i, r_j\}$, then*

$$\phi_{\sigma}^{\alpha}(r_i) = \phi_{\sigma}^{\alpha}(r_j).$$

Proof. According to Definition 5, if for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r_i, r_j\}$, we have $\sigma_{\mathcal{S} \cup \{r_i\}}(\alpha) = \sigma_{\mathcal{S} \cup \{r_j\}}(\alpha)$, then we have $\sigma_{\mathcal{S} \cup \{r_i, r_j\}}(\alpha) - \sigma_{\mathcal{S} \cup \{r_i\}}(\alpha) = \sigma_{\mathcal{S} \cup \{r_j, r_i\}}(\alpha) - \sigma_{\mathcal{S} \cup \{r_j\}}(\alpha)$. We let $\mathcal{S}_i = \mathcal{S} \cup \{r_i\}$ and $\mathcal{S}_j = \mathcal{S} \cup \{r_j\}$. Thus, $\sigma_{\mathcal{S}_i \cup \{r_j\}}(\alpha) - \sigma_{\mathcal{S}_i}(\alpha) = \sigma_{\mathcal{S}_j \cup \{r_i\}}(\alpha) - \sigma_{\mathcal{S}_j}(\alpha)$. Therefore, $\sum_{\mathcal{S}_i \subseteq \mathcal{R} \setminus \{r_j\}} \frac{(|\mathcal{R}| - |\mathcal{S}_i| - 1)! |\mathcal{S}_i|!}{|\mathcal{R}|!} [\sigma_{\mathcal{S}_i \cup \{r_j\}}(\alpha) - \sigma_{\mathcal{S}_i}(\alpha)] = \sum_{\mathcal{S}_j \subseteq \mathcal{R} \setminus \{r_i\}} \frac{(|\mathcal{R}| - |\mathcal{S}_j| - 1)! |\mathcal{S}_j|!}{|\mathcal{R}|!} [\sigma_{\mathcal{S}_j \cup \{r_i\}}(\alpha) - \sigma_{\mathcal{S}_j}(\alpha)]$, hence $\phi_{\sigma}^{\alpha}(r_i) = \phi_{\sigma}^{\alpha}(r_j)$, which completes the proof. \square

Proposition 5 (Dominance). *Let $\alpha \in \mathcal{A}$ and $r_i, r_j \in \mathcal{R}$ with $r_i \neq r_j$. If $\exists \mathcal{S}' \subseteq \mathcal{R} \setminus \{r_i, r_j\}$ such that $\sigma_{\mathcal{S}' \cup \{r_i\}}(\alpha) > \sigma_{\mathcal{S}' \cup \{r_j\}}(\alpha)$ and $\forall \mathcal{S}'' \subseteq \mathcal{R} \setminus \{r_i, r_j\}$ ($\mathcal{S}' \neq \mathcal{S}''$) such that $\sigma_{\mathcal{S}'' \cup \{r_i\}}(\alpha) \geq \sigma_{\mathcal{S}'' \cup \{r_j\}}(\alpha)$, then*

$$\phi_{\sigma}^{\alpha}(r_i) > \phi_{\sigma}^{\alpha}(r_j).$$

Proof. For \mathcal{S}' we have $\sigma_{\mathcal{S}' \cup \{r_i, r_j\}}(\alpha) - \sigma_{\mathcal{S}' \cup \{r_j\}}(\alpha) > \sigma_{\mathcal{S}' \cup \{r_i, r_j\}}(\alpha) - \sigma_{\mathcal{S}' \cup \{r_i\}}(\alpha)$; for any \mathcal{S}'' we have $\sigma_{\mathcal{S}'' \cup \{r_i, r_j\}}(\alpha) - \sigma_{\mathcal{S}'' \cup \{r_j\}}(\alpha) \geq \sigma_{\mathcal{S}'' \cup \{r_i, r_j\}}(\alpha) - \sigma_{\mathcal{S}'' \cup \{r_i\}}(\alpha)$. According to Definition 5, we sum up \mathcal{S}' and \mathcal{S}'' , then we have $\sum_{\mathcal{S} \subseteq \mathcal{R} \setminus \{r_i\}} \frac{(|\mathcal{R}| - |\mathcal{S}| - 1)! |\mathcal{S}|!}{|\mathcal{R}|!} [\sigma_{\mathcal{S} \cup \{r_i\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)] > \sum_{\mathcal{S} \subseteq \mathcal{R} \setminus \{r_j\}} \frac{(|\mathcal{R}| - |\mathcal{S}| - 1)! |\mathcal{S}|!}{|\mathcal{R}|!} [\sigma_{\mathcal{S} \cup \{r_j\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)]$, hence $\phi_{\sigma}^{\alpha}(r_i) > \phi_{\sigma}^{\alpha}(r_j)$, which completes the proof. \square

C.2 Argumentative Properties

Property 1 (Sign Correctness). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$.*

1. *If $r \in \mathcal{R}^-$, then $\phi_\sigma^\alpha(r) \leq 0$;*
2. *If $r \in \mathcal{R}^+$, then $\phi_\sigma^\alpha(r) \geq 0$.*

Proposition 6. *Let r be a direct edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies sign correctness if σ satisfies monotonicity.*

Proof. For any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$, we have $\frac{(|\mathcal{R}|-|\mathcal{S}|-1)!!|\mathcal{S}!}{|\mathcal{R}|!} > 0$. Since σ satisfies monotonicity, $r \in \mathcal{R}^-$ implies $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) \leq \sigma_{\mathcal{S}}(\alpha)$. According to Definition 5, $\phi_\sigma^\alpha(r) \leq 0$. Analogously, if $r \in \mathcal{R}^+$, then $\phi_\sigma^\alpha(r) \geq 0$. \square

Proposition 7. *Let r be an indirect edge w.r.t. α . Suppose the path sequence from r to α is $r, r_1, \dots, r_n (n \geq 1)$. Let $\lambda = |\{r_1, \dots, r_n\} \cap \mathcal{R}^-|$. Then the following statements hold if σ satisfies monotonicity.*

1. *If $r \in \mathcal{R}^-$ and λ is odd, then $\phi_\sigma^\alpha(r) \geq 0$;*
2. *If $r \in \mathcal{R}^-$ and λ is even, then $\phi_\sigma^\alpha(r) \leq 0$;*
3. *If $r \in \mathcal{R}^+$ and λ is odd, then $\phi_\sigma^\alpha(r) \leq 0$;*
4. *If $r \in \mathcal{R}^+$ and λ is even, then $\phi_\sigma^\alpha(r) \geq 0$.*

Proof. Case 1 (If $r \in \mathcal{R}^-$ and λ is odd): for any $\mathcal{S} \subseteq \mathcal{R}$, if $\{r_1, \dots, r_n\} \not\subseteq \mathcal{S}$, then $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) = \sigma_{\mathcal{S}}(\alpha)$ as the path from r to α remains broken, that is $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha) = 0$; if $\{r_1, \dots, r_n\} \subseteq \mathcal{S}$, then r will not decrease $\sigma(\alpha)$ if σ satisfies monotonicity, because r passes an odd number of attacks on the path to α , hence $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha) \geq 0$. Therefore, $\phi_\sigma^\alpha(r) = \sum_{\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}} \frac{(|\mathcal{R}|-|\mathcal{S}|-1)!!|\mathcal{S}!}{|\mathcal{R}|!} [\sigma_{\mathcal{S} \cup \{r\}}(\alpha) - \sigma_{\mathcal{S}}(\alpha)] \geq 0$. Analogously, if $r \in \mathcal{R}^-$ and λ is even, then $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) \leq \sigma_{\mathcal{S}}(\alpha)$ thus $\phi_\sigma^\alpha(r) \leq 0$; if $r \in \mathcal{R}^+$ and λ is odd, then $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) \leq \sigma_{\mathcal{S}}(\alpha)$ thus $\phi_\sigma^\alpha(r) \leq 0$; if $r \in \mathcal{R}^+$ and λ is even, then $\sigma_{\mathcal{S} \cup \{r\}}(\alpha) \geq \sigma_{\mathcal{S}}(\alpha)$ thus $\phi_\sigma^\alpha(r) \geq 0$. \square

Proposition 8. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate sign correctness even if σ satisfies monotonicity.*

Proof. A counterexample is shown in Figure 3, where $(\zeta, \delta) \in \mathcal{R}^-$ is a multifold edge w.r.t. α , but $\phi_{\sigma_{DF}}^\alpha((\zeta, \delta)) \approx 0.0307 > 0$. \square

Property 2 (Counterfactuality). *Let $\alpha \in \mathcal{A}$ and $r \in \mathcal{R}$.*

1. *If $\phi_\sigma^\alpha(r) < 0$, then $\sigma(\alpha) \leq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$;*
2. *If $\phi_\sigma^\alpha(r) > 0$, then $\sigma(\alpha) \geq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$.*

Proposition 9. *Let r be a direct or indirect edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies counterfactuality if σ satisfies monotonicity.*

Proof. Direct: According to Proposition 6, we have $r \in \mathcal{R}^- \Rightarrow \phi_\sigma^\alpha(r) \leq 0$, hence $\phi_\sigma^\alpha(r) > 0 \Rightarrow r \notin \mathcal{R}^-$, that is, $\phi_\sigma^\alpha(r) > 0 \Rightarrow r \in \mathcal{R}^+$. Since we assume monotonicity of σ , if $r \in \mathcal{R}^+$, then $\sigma(\alpha) \geq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$. Analogously, if $\phi_\sigma^\alpha(r) < 0$, then $\sigma(\alpha) \leq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$.

Indirect: According to Proposition 7, When λ is even, $r \in \mathcal{R}^- \Rightarrow \phi_\sigma^\alpha(r) \leq 0$; $r \in \mathcal{R}^+ \Rightarrow \phi_\sigma^\alpha(r) \geq 0$. Therefore, $\phi_\sigma^\alpha(r) > 0 \Rightarrow r \in \mathcal{R}^+$; $\phi_\sigma^\alpha(r) < 0 \Rightarrow r \in \mathcal{R}^-$. Analogously, when λ is odd, we have the conclusion $\phi_\sigma^\alpha(r) > 0 \Rightarrow r \in \mathcal{R}^-$; $\phi_\sigma^\alpha(r) < 0 \Rightarrow r \in \mathcal{R}^+$. Therefore, if $\phi_\sigma^\alpha(r) > 0$, then it means either $r \in \mathcal{R}^+$ and λ is even, or $r \in \mathcal{R}^-$ and λ is odd. In both cases, according to the proof in Proposition 7, we know $\sigma(\alpha) \geq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$. Similarly, if $\phi_\sigma^\alpha(r) < 0$, then it means either $r \in \mathcal{R}^-$ and λ is even, or $r \in \mathcal{R}^+$ and λ is odd. In both cases, according to the proof in Proposition 7, we know $\sigma(\alpha) \leq \sigma_{\mathcal{R} \setminus \{r\}}(\alpha)$. \square

Proposition 10. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate counterfactuality even if σ satisfies monotonicity.*

Proof. A counterexample is shown in Figure 9. σ is given by DF-QuAD since it satisfies monotonicity. Let us set $\tau(\mu) = 0.4$ and $\tau(\alpha) = \tau(\beta) = \tau(\gamma) = \tau(\zeta) = \tau(\rho) = 0$. (μ, ρ) is a multifold edge w.r.t. α because it passes 3 (more than 1) paths to reach α . By DF-QuAD semantics, we have $\sigma_{DF}^{DF}(\alpha) = 0.384$. By Definition 5, we have $\phi_{\sigma_{DF}}^\alpha((\mu, \rho)) \approx 0.0458 > 0$. If (μ, ρ) is removed, then we apply DF-QuAD again and we have $\sigma_{\mathcal{R} \setminus \{(\mu, \rho)\}}^{DF}(\alpha) = 0.4$. Therefore, when $\phi_{\sigma_{DF}}^\alpha((\mu, \rho)) > 0$, $\sigma_{DF}^{DF}(\alpha) = 0.384 < \sigma_{\mathcal{R} \setminus \{(\mu, \rho)\}}^{DF}(\alpha) = 0.4$, which violates counterfactuality. \square

Property 3 (Qualitative Invariability). *Let $\alpha, \beta \in \mathcal{A}$ and $r \in \mathcal{R}(\beta)$. Let ϕ_δ denote $\phi_\sigma^\alpha(r)$ when setting $\tau(\beta)$ to some $\delta \in \mathbb{I}$.*

1. *If $\phi_\sigma^\alpha(r) < 0$, then $\forall \delta \in \mathbb{I}, \phi_\delta \leq 0$;*
2. *If $\phi_\sigma^\alpha(r) > 0$, then $\forall \delta \in \mathbb{I}, \phi_\delta \geq 0$.*

Proposition 11. *Let r be a direct or indirect edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies qualitative invariability if σ satisfies monotonicity.*

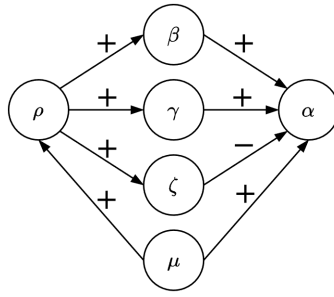


Figure 9: A counterexample.

Proof. Direct: According to the proof in Proposition 9, we have $\phi_\sigma^\alpha(r) > 0 \Rightarrow r \in \mathcal{R}^+$. According to Proposition 6, we have $r \in \mathcal{R}^+ \Rightarrow \phi_\sigma^\alpha(r) \geq 0$. Therefore, if $\phi_\sigma^\alpha(r) > 0$, then $\forall \delta \in \mathbb{I}, \phi_\delta \geq 0$. Analogously, if $\phi_\sigma^\alpha(r) < 0$, then $\forall \delta \in \mathbb{I}, \phi_\delta \leq 0$.

Indirect: According to the proof in Proposition 9, we have:

1. $\phi_\sigma^\alpha(r) < 0 \Rightarrow (r \in \mathcal{R}^- \wedge \lambda \text{ is even}) \vee (r \in \mathcal{R}^+ \wedge \lambda \text{ is odd})$;
2. $\phi_\sigma^\alpha(r) > 0 \Rightarrow (r \in \mathcal{R}^- \wedge \lambda \text{ is odd}) \vee (r \in \mathcal{R}^+ \wedge \lambda \text{ is even})$.

In case 1, according to Proposition 7, we have $(r \in \mathcal{R}^- \wedge \lambda \text{ is even}) \Rightarrow \phi_\sigma^\alpha(r) \leq 0$ and $(r \in \mathcal{R}^+ \wedge \lambda \text{ is odd}) \Rightarrow \phi_\sigma^\alpha(r) \leq 0$. Therefore, in both case, $\forall \delta \in \mathbb{I}, \phi_\delta \leq 0$. Analogously to the proof in case 2. \square

Proposition 12. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate qualitative invariability even if σ satisfies monotonicity.*

Proof. A counterexample is shown in Figure 9. σ is given by DF-QuAD since it satisfies monotonicity. Let us set $\tau(\mu) = 0.4$ and $\tau(\alpha) = \tau(\beta) = \tau(\gamma) = \tau(\zeta) = \tau(\rho) = 0$. (μ, ρ) is a multifold edge w.r.t. α because it passes 3 (more than 1) paths to reach α . By Definition 5, we have $\phi_{\sigma_{DF}}^\alpha((\mu, \rho)) \approx 0.0458 > 0$. Let $\delta = 0.7$, then we have $\phi_\delta \approx -0.0283 < 0$. Since $\phi_{\sigma_{DF}}^\alpha((\mu, \rho))$ does not keep positive for any δ , thus qualitative invariability is violated. \square

Property 4 (Quantitative Variability). *Let $\alpha, \beta \in \mathcal{A}$, $\alpha \neq \beta$, and $\mathcal{R}(\beta) = \{r\}$. Let ϕ_δ denote $\phi_\sigma^\alpha(r)$ when setting $\tau(\beta)$ to some $\delta \in \mathbb{I}$.*

1. *If $\delta < \tau(\beta)$, then $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$;*
2. *If $\delta > \tau(\beta)$, then $|\phi_\delta| \geq |\phi_\sigma^\alpha(r)|$.*

Proposition 13. *Let r be a direct or indirect edge w.r.t. α . $\phi_\sigma^\alpha(r)$ satisfies quantitative variability if σ satisfies monotonicity.*

Proof. Direct: We let $\sigma_S^\delta(\alpha)$ denote the strength of α in \mathcal{Q}^{ls} when setting $\tau(\beta)$ to some $\delta \in \mathbb{I}$.

Suppose $r \in \mathcal{R}^+$. As we assume monotonicity of σ , for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$, if $\delta < \tau(\beta)$, then $\sigma_{\mathcal{S} \cup \{r\}}^\delta(\alpha) \leq \sigma_{\mathcal{S} \cup \{r\}}(\alpha)$. Since $\sigma_S^\delta(\alpha) = \sigma_S(\alpha)$, according to Definition 5, we have $0 < \phi_\delta \leq \phi_\sigma^\alpha(r)$, i.e., $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$.

Suppose $r \in \mathcal{R}^-$. By monotonicity of σ , for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$ if $\delta < \tau(\beta)$, then $\sigma_{\mathcal{S} \cup \{r\}}^\delta(\alpha) \geq \sigma_{\mathcal{S} \cup \{r\}}(\alpha)$. According to Definition 5, since $\sigma_S^\delta(\alpha) = \sigma_S(\alpha)$, we have $0 > \phi_\delta \geq \phi_\sigma^\alpha(r)$, i.e., $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$. Therefore, if $\delta < \tau(\beta)$, then $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$. Analogously, if $\delta > \tau(\beta)$, then $|\phi_\delta| \geq |\phi_\sigma^\alpha(r)|$.

Indirect: We let $\sigma_S^\delta(\alpha)$ denote the strength of α in \mathcal{Q}^{ls} when setting $\tau(\beta)$ to some $\delta \in \mathbb{I}$.

1. Suppose $r \in \mathcal{R}^+$ and λ is even. Under the guarantee of monotonicity, for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$, if $\delta < \tau(\beta)$, then $\sigma_{\mathcal{S} \cup \{r\}}^\delta(\alpha) \leq \sigma_{\mathcal{S} \cup \{r\}}(\alpha)$. Since $\sigma_S^\delta(\alpha) = \sigma_S(\alpha)$, according to Definition 5, we have $0 < \phi_\delta \leq \phi_\sigma^\alpha(r)$, i.e., $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$.
2. Suppose $r \in \mathcal{R}^-$ and λ is odd. Under the guarantee of monotonicity, for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$, if $\delta < \tau(\beta)$, then $\sigma_{\mathcal{S} \cup \{r\}}^\delta(\alpha) \leq \sigma_{\mathcal{S} \cup \{r\}}(\alpha)$. Since $\sigma_S^\delta(\alpha) = \sigma_S(\alpha)$, according to Definition 5, we have $0 < \phi_\delta \leq \phi_\sigma^\alpha(r)$, i.e., $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$.
3. Suppose $r \in \mathcal{R}^+$ and λ is odd. Under the guarantee of monotonicity, for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$ if $\delta < \tau(\beta)$, then $\sigma_{\mathcal{S} \cup \{r\}}^\delta(\alpha) \geq \sigma_{\mathcal{S} \cup \{r\}}(\alpha)$. Since $\sigma_S^\delta(\alpha) = \sigma_S(\alpha)$, according to Definition 5, we have $0 > \phi_\delta \geq \phi_\sigma^\alpha(r)$, i.e., $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$.
4. Suppose $r \in \mathcal{R}^-$ and λ is even. Under the guarantee of monotonicity, for any $\mathcal{S} \subseteq \mathcal{R} \setminus \{r\}$ if $\delta < \tau(\beta)$, then $\sigma_{\mathcal{S} \cup \{r\}}^\delta(\alpha) \geq \sigma_{\mathcal{S} \cup \{r\}}(\alpha)$. Since $\sigma_S^\delta(\alpha) = \sigma_S(\alpha)$, according to Definition 5, we have $0 > \phi_\delta \geq \phi_\sigma^\alpha(r)$, i.e., $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$.

Therefore, if $\delta < \tau(\beta)$, then $|\phi_\delta| \leq |\phi_\sigma^\alpha(r)|$. Analogously, if $\delta > \tau(\beta)$, then $|\phi_\delta| \geq |\phi_\sigma^\alpha(r)|$. \square

Proposition 14. *Let r be a multifold edge w.r.t. α . $\phi_\sigma^\alpha(r)$ may violate quantitative variability even if σ satisfies monotonicity.*

Proof. A counterexample is shown in Figure 10. σ is given by DF-QuAD since it satisfies monotonicity. Let us set $\tau(\mu) = 0.8$, $\tau(\alpha) = 0.5$ and $\tau(\beta) = \tau(\gamma) = \tau(\zeta) = \tau(\rho) = 0$. (μ, ρ) is a multifold edge w.r.t. α because it passes 3 (more than 1) paths to reach α . By Definition 5, we have $\phi_{\sigma_{DF}}^{\alpha}((\mu, \rho)) \approx -0.0693$. Let $\delta = 1.0$, then we have $\phi_{\delta} \approx -0.0667$. Since $\tau(\mu)$ increases from 0.8 to 1.0 while $|\phi_{\sigma_{DF}}^{\alpha}((\mu, \rho))|$ decreases from 0.0693 to 0.0667, thus quantitative variability is violated.

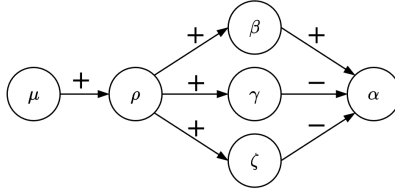


Figure 10: A counterexample.

□

D Proofs for Section 6

Proposition 15 (Computing RAEs exactly). *RAEs can be computed in time $\mathcal{O}(n \cdot 2^n \cdot T(m, n))$.*

Proof. The strength can be computed in time $T(m, n)$ by assumption. The exact RAE for one edge can thus be computed in time $\mathcal{O}(2^n \cdot T(m, n))$. Therefore, the total runtime is $\mathcal{O}(n \cdot 2^n \cdot T(m, n))$. □

Proposition 17 (Approximating RAEs). *If the number of samples for each edge is N , then approximate RAEs can be generated in time $\mathcal{O}(n \cdot N \cdot T(m, n))$.*

Proof. The strength can be computed in time $T(m, n)$ by assumption. The approximate RAE for each edge can thus be computed in time $\mathcal{O}(N \cdot T(m, n))$. Hence, the total runtime is $\mathcal{O}(n \cdot N \cdot T(m, n))$. □

E Additional Details for Section 1

We show the concrete values of the RAEs in Figure 1.

In Figure 1, the base scores are set as $\tau(\alpha) = 0.5$, $\tau(\beta) = 0.9$, $\tau(\gamma) = 0.1$, $\tau(\delta) = 0.1$, $\tau(\xi) = 0.1$, $\tau(\eta) = 0.1$. Applying DF-QuAD semantics according to Section A.1, we have $\sigma^{DF}(\alpha) = 0.95095$, $\sigma^{DF}(\beta) = 0.9$, $\sigma^{DF}(\gamma) = 0.91$, $\sigma^{DF}(\delta) = 0.91$, $\sigma^{DF}(\xi) = 0.009$, $\sigma^{DF}(\eta) = 0.0099$. We apply RAEs, we have $\phi_{\sigma_{DF}}^{\alpha}(r_1) = 0.1981125$, $\phi_{\sigma_{DF}}^{\alpha}(r_2) = 0.2471125$, $\phi_{\sigma_{DF}}^{\alpha}(r_3) = 0.0135$, $\phi_{\sigma_{DF}}^{\alpha}(r_4) = 0.016$, $\phi_{\sigma_{DF}}^{\alpha}(r_5) = -0.034$, $\phi_{\sigma_{DF}}^{\alpha}(r_6) = -0.011875$, $\phi_{\sigma_{DF}}^{\alpha}(r_7) = 0.0221125$.

F Additional Details for Section 3

Example 1. *Consider the QBAF in Figure 11. The base scores of all arguments are set to 0.5, and the strengths of arguments are evaluated by DF-QuAD gradual semantics denoted by σ^{DF} .*

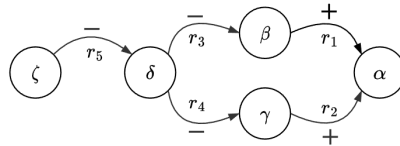


Figure 11: An example of QBAF.

According to the definition of DF-QuAD in Section A.1:

For ζ , $v_{\zeta-att} = 0$ and $v_{\zeta-sup} = 0$, hence $\sigma^{DF}(\zeta) = \tau(\zeta) = 0.5$.

For δ , $v_{\delta-att} = 0.5$ and $v_{\delta-sup} = 0$, hence $\sigma^{DF}(\delta) = \tau(\delta) - \tau(\delta) \cdot (v_{\delta-att} - v_{\delta-sup}) = 0.25$.

For γ , $v_{\gamma-att} = 0.25$ and $v_{\gamma-sup} = 0$, hence $\sigma^{DF}(\gamma) = \tau(\gamma) - \tau(\gamma) \cdot (v_{\gamma-att} - v_{\gamma-sup}) = 0.375$.

For β , $v_{\beta-att} = 0.25$ and $v_{\beta-sup} = 0$, hence $\sigma^{DF}(\beta) = \tau(\beta) - \tau(\beta) \cdot (v_{\beta-att} - v_{\beta-sup}) = 0.375$.

For α , $v_{\alpha-att} = 0$ and $v_{\alpha-sup} = 1 - (1 - 0.375)(1 - 0.375) = 0.609375$, hence $\sigma^{DF}(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot (v_{\alpha-sup} - v_{\alpha-att}) = 0.5 + 0.5 \times 0.609375 = 0.8046875$.

To sum up, we have $\sigma^{DF}(\alpha) = 0.8046875$, $\sigma^{DF}(\beta) = 0.375$, $\sigma^{DF}(\gamma) = 0.375$, $\sigma^{DF}(\delta) = 0.25$, $\sigma^{DF}(\zeta) = 0.5$.

G Additional Experimental Results for Section 6

G.1 Hardware Specifications

We ran all experiments on a Windows PC (OS: Windows 10 Enterprise, Version: 22H2, 64-bit operating system; Processor: AMD Ryzen 7 PRO 3700 8-Core Processor, 3.59 GHz; Memory: 16.0 GB).

G.2 Additional Experimental Results

Figure 12 and Figure 13 show additional experimental results of convergence for larger acyclic and cyclic QBAFs, respectively. Figure 12 is based on DF-QuAD gradual semantics, while Figure 13 is based on QE semantics.

Tables 1 - 4 show the runtime in milliseconds of one trial for different sizes of acyclic/cyclic QBAFs.

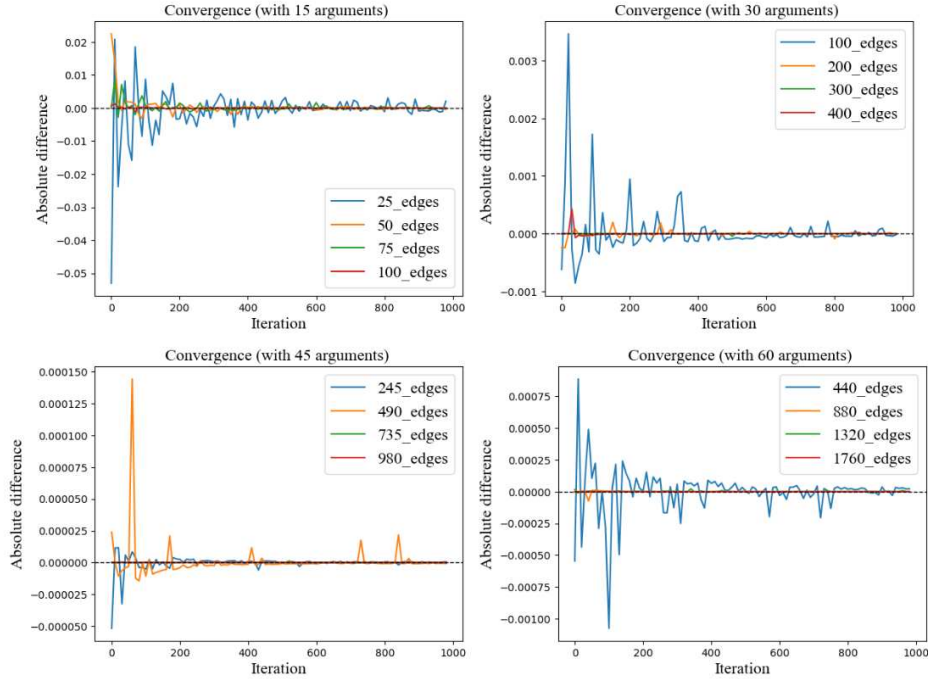


Figure 12: The results are based on random acyclic QBAFs with different number of arguments and edges.

Table 1: Runtime(ms) for random QBAFs with 15 arguments and 25/50/75/100 edges.

	25	50	75	100
15 (acyclic)	0.9	1.0	1.1	1.3
15 (cyclic)	14	18	17	23

Table 2: Runtime(ms) for random QBAFs with 30 arguments and 100/200/300/400 edges.

	100	200	300	400
30 (acyclic)	1.5	2.2	3.2	5.6
30 (cyclic)	36	37	47	62

Table 3: Runtime(ms) for random QBAFs with 45 arguments and 245/490/735/980 edges.

	245	490	735	980
45 (acyclic)	2.5	5.7	10.4	15.3
45 (cyclic)	64	101	127	150

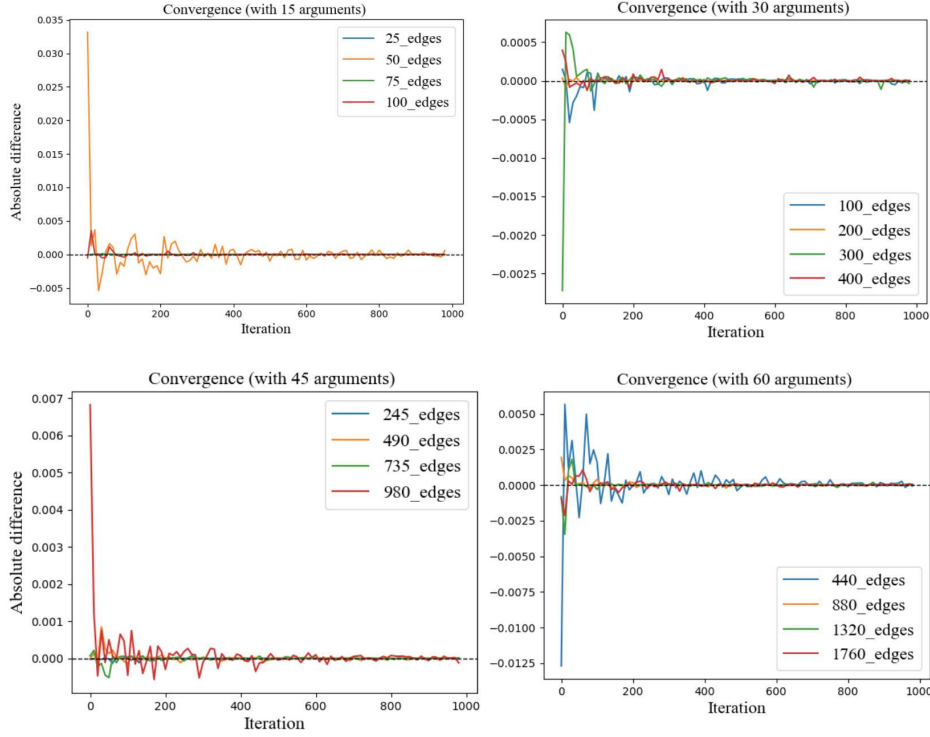


Figure 13: The results are based on random cyclic QBAFs with different number of arguments and edges.

Table 4: Runtime(ms) for random QBAFs with 60 arguments and 440/880/1320/1760 edges.

	440	880	1320	1760
60 (acyclic)	5	14	27	43
60 (cyclic)	106	142	216	296

H Additional Details for Section 7

H.1 Case Study 1: Fraud Detection

We show the contents of the arguments in Figure 5 in Table 5. The specific RAEs of edges are shown in Table 6.

H.2 Case Study 2: Large Language Models

According to the definition of the QE semantics in Section A.2:

For β , $E_\beta = 0$, hence $\sigma(\beta) = \tau(\beta) = 0.6$.

For γ , $E_\gamma = 0$, hence $\sigma(\gamma) = \tau(\gamma) = 0.9$.

For δ , $E_\delta = \sigma(\gamma) - \sigma(\beta) = 0.3$, hence $\sigma(\delta) = \tau(\delta) + (1 - \tau(\delta)) \cdot \frac{E_\delta^2}{1 + E_\delta^2} = 0.72$.

For α , $E_\alpha = \sigma(\gamma) + \sigma(\delta) - \sigma(\beta) = 1.02$, hence $\sigma(\alpha) = \tau(\alpha) + (1 - \tau(\alpha)) \cdot \frac{E_\alpha^2}{1 + E_\alpha^2} = 0.90$.

We next show the exact prompt we used in the case study. However, the response from ChatGPT may vary because the outputs are generated by a sampling algorithm⁷, etc.

Prompt:

Given an argument α : ‘It is easy for children to learn a foreign language well.’

Could you please provide 3 arguments β , γ , and δ that satisfy the following requirements? (each argument should be around 20 words)

⁷The transformer in ChatGPT can only predict token probabilities and output sequences have to be generated by repeatedly sampling tokens and feeding them back into the input. Since the number of tokens is in the order of ten thousands, it is impossible to look at all possible token sequences and token sequences are generated by randomized heuristic search algorithms.

requirement 1: Provide an argument β attacking α ;
requirement 2: Provide an argument γ supporting α ;
requirement 3: Provide an argument δ supporting α ;
requirement 4: β directly attacks δ ;
requirement 5: γ directly supports δ ;
requirement 6: Provide confidence score for α , β , γ , and δ as their initial weights ranging from 0 to 1.
In the end, give some explanations why the provided arguments satisfy the requirements.

Response:

Sure, here are the three arguments as per the requirements:

α : "It is easy for children to learn a foreign language well." Confidence score for α : 0.8

β : "Learning a foreign language requires cognitive maturity, which children lack. Hence, it's difficult for them to excel." Confidence score for β : 0.6

γ : "Studies show that young children possess higher neuroplasticity, making language learning more effective." Confidence score for γ : 0.9

δ : "Children immersed in a foreign language environment from an early age have better language acquisition." Confidence score for δ : 0.7

Explanations:

1. Argument β attacks argument α by emphasizing that children lack the cognitive maturity required to learn a foreign language well. It highlights the challenges faced by children in excelling in language learning, countering the notion that it is easy for them.

2. Argument γ supports α by citing studies that demonstrate young children's higher neuroplasticity, which refers to the brain's ability to adapt and learn. This evidence suggests that children may be more adept at learning a foreign language than adults.

3. Argument δ also supports α by stating that children exposed to a foreign language environment from an early age have better language acquisition. This implies that children have an advantage in learning a foreign language when immersed in the language consistently.

4. Argument β directly attacks argument δ by disregarding the benefits of early exposure to a foreign language and arguing that cognitive maturity is the primary factor affecting language learning.

5. Argument γ directly supports argument δ by reinforcing the idea that young children have a higher capacity for language learning, aligning with the notion that early exposure to a foreign language is beneficial.

The initial confidence scores assigned to each argument are subjective and reflect the confidence in the validity and strength of the arguments based on available evidence and reasoning.

Overall, the provided arguments meet the specified requirements, with one argument attacking α and two arguments supporting it. Additionally, β directly challenges δ , while γ directly aligns with δ , showcasing a clear connection between the supporting arguments. The confidence scores are arbitrary and represent a subjective assessment of the persuasiveness of each argument.

Arg	Content	Arg	Content
2	It is a fraud case.	25	The basic facts of fabrication stand.
3	It is not a fraud case.	26	The victim had abnormal behaviours.
4	The fraud elements are satisfied.	27	The victim delivered the property voluntarily.
5	The fraud elements are not satisfied.	28	The victim delivered the property consciously.
6	The suspect is with limited capacity.	29	The suspect refused to pay refund.
7	The suspect is with incapacity.	30	The suspect betrayed the victim's transaction purpose.
8	The suspect is under the age of criminal responsibility.	31	The property was transferred to the suspect.
9	The suspect has no capacity of criminal responsibility.	32	The collection account is a black account.
10	The suspect had indirect intention.	33	The suspects terminated contact with the victim.
11	The suspect fell into negligence.	34	The victim's complaint is not credible.
12	The behavior elements of fraud cases are satisfied.	35	The transaction is illegal.
13	The result elements of fraud cases are satisfied.	36	The transaction occurred between acquaintances.
14	The suspect had direct intention for fraud.	37	The suspect faked official materials.
15	The auxiliary facts elements are satisfied.	38	The transaction was at variance with business logic.
16	The victim has clear cognition.	39	The suspect faked official identity.
17	The suspect fabricates facts.	40	The victim belongs to a gullible group.
18	The victim fell into cognitive error.	41	The transaction contained high-risk operations.
19	The victim's property has been delivered.	42	The transaction was abnormal.
20	The suspect has obtained the victim's property.	43	The collection account is a mature account.
21	The suspect had an illegal possession purpose.	44	The collection account is an account for daily use.
22	The case is different from common fraud.	45	The collection behavior is suspicious.
23	The victim's complaints are malicious.	46	The collection account has a high-risk record.
24	The victim had speculative motivations.	47	The suspect removed the victim from the communication list.
		48	The suspect had no response.

Table 5: Contents of arguments in Figure 5.

Table 6: RAEs in ascending order for the QBAF in Figure 6.

Relation	RAE	Relation	RAE	Relation	RAE	Relation	RAE
att (3,1)	-4.55E-01	sup (23,16)	-1.60E-03	sup (39,25)	3.49E-05	sup (33,21)	4.95E-04
att (5,2)	-9.13E-02	sup (34,23)	-4.74E-04	sup (41,26)	3.79E-05	sup (18,12)	1.27E-03
sup (16,5)	-1.75E-02	sup (35,24)	-3.62E-04	sup (46,32)	6.24E-05	sup (19,13)	1.38E-03
sup (8,3)	-1.64E-02	sup (36,24)	-3.52E-04	sup (48,33)	6.71E-05	sup (20,13)	1.57E-03
sup (7,3)	-1.59E-02	att (44,32)	-5.82E-05	sup (30,20)	1.40E-04	sup (17,12)	1.73E-03
sup (11,3)	-1.51E-02	att (43,32)	-4.06E-05	sup (31,20)	1.47E-04	sup (21,14)	1.78E-03
sup (6,3)	-1.51E-02	sup (37,25)	1.60E-05	sup (28,19)	1.49E-04	sup (14,4)	8.41E-03
sup (10,3)	-1.46E-02	sup (42,26)	2.31E-05	sup (29,20)	1.49E-04	sup (13,4)	9.36E-03
sup (9,3)	-1.32E-02	sup (38,25)	2.56E-05	sup (27,19)	2.12E-04	sup (12,4)	9.59E-03
sup (15,5)	-1.23E-02	sup (40,26)	2.78E-05	sup (32,21)	2.17E-04	sup (4,2)	1.15E-01
sup (24,16)	-2.89E-03	sup (45,32)	2.79E-05	sup (25,17)	2.98E-04	sup (2,1)	2.60E-01
sup (22,15)	-2.31E-03	sup (47,33)	3.09E-05	sup (26,18)	3.74E-04		